SMA 6304 Factory Planning and Scheduling Lecture 18: Single-stage, multiple part type, MTS

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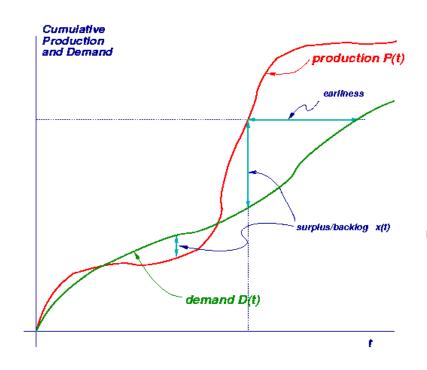
Make to Stock

- Ideal: The product is available (in stock) when the order arrives.
- Service rate: the fraction of orders that are fulfilled with no waiting.
- Costs: Inventory, lost sales.

Make to Stock

- Trade-off:
 - ★ The higher the service rate, the happier the customers (so sales go up).
 - ★ The lower the service rate, the less inventory is required (so costs go down).
- The Quantification Dilemma: Costs are easy to quantify. Benefits are not.

Objective



Objective is to keep the cumulative production line close to the cumulative demand line.

- Setup: It costs less to make a Type i part after making a Type i part than after making a Type j part.
- Examples:
 - * Tool change
 - ★ Paint color change

Costs

Setups

Setup costs can include

- Money costs, especially in labor. Also materials.
- Time, in loss of capacity and delay.
- Setup motivates lots or batches: a set of parts that are processed without interruption by setups.

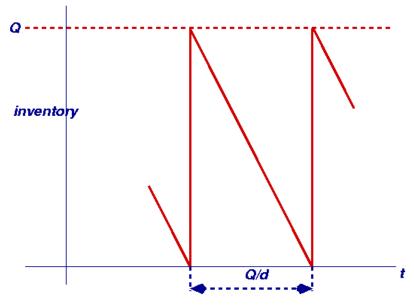
Setups

(EOQ)

- Money cost only.
- ullet Each time the firm obtains a lot of Q items, it must pay A+cQ dollars for that lot.
 - \star The unit cost is c, and the cost of ordering is A.
- The firm orders Q items when its inventory level is 0, and it receives them instantly.

Setups

(EOQ)



- Inventory is depleted at a constant rate d until it is 0, at time Q/d.
- The average inventory level is Q/2.

Setups

EOQ Costs

Holding cost

- If h is the dollar cost per time unit of holding an item in inventory, then hQ/2 is the average inventory holding cost per time unit.
- Over a long time interval of length T, the total holding cost is hQT/2.

Setups

EOQ Costs

Acquisition cost

- Over that interval, Td units of material is acquired, in lot sizes of Q. The *number* of times it is ordered is Td/Q.
- ullet The total cost of acquiring material is (Td/Q)(A+cQ)=TdA/Q+Tdc

Setups

Cost Minimization

The total cost over the interval is

$$rac{TdA}{Q} + Tdc + rac{ThQ}{2} = T\left\{rac{dA}{Q} + dc + rac{hQ}{2}
ight\}.$$

• The minimizing lot size is therefore

$$Q^\star = \left(rac{2dA}{h}
ight)^{rac{1}{2}}.$$

Setups

The Order-up-to policy

In a random environment, EOQ can be converted into a real-time policy:

• When the inventory goes to 0, order Q^* units.

When delivery is not instantaneous, a variation is

• When the inventory goes below some level Q_{min} , order enough to bring it up to Q^* units.

Setups

Loss of Capacity

Assume

- there is one setup for every Q parts (Q=lot size),
- \bullet the setup time is S,
- \bullet the time to process a part is τ .

Then the time to process Q parts is $S+Q\tau$. The average time to process one part is $\tau+S/Q$.

Loss of Capacity

If the demand rate is d parts per time unit, then the demand is feasible only if

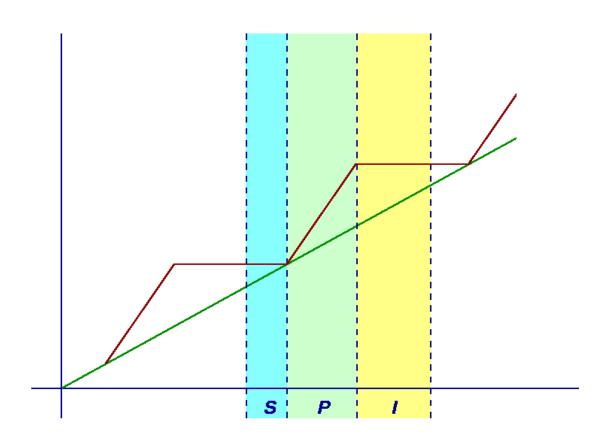
$$au + S/Q < 1/d$$
 or $d < rac{Q}{S+Q au} < rac{1}{ au}$

This is not satisfied if S is too large or Q is too small.

- Single part type (simplification!)
- Constant demand
- Deterministic setup and operation times.
- Setup/production/(idleness) cycles.
- Policy: Produce at maximum rate until the inventory is enough to last through the next setup time.

Setups

Cycle:



Cycle:

- Setup period. Duration: S. Production: 0. Demand: dS. Net change of surplus, ie of P-D is $\Delta_S=-dS$.
- Production period. Duration: $t=Q\tau$. Production: Q. Demand: dt. Net change of P-D is $\Delta_P=Q-dt=Q\tau(1/\tau-d)=Q(1-d\tau)$.

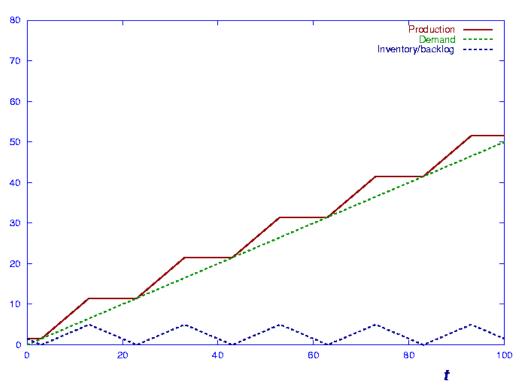
- *Idleness period.* Duration: *I*. Production: 0. Demand: dI. Net change of P D is $\Delta_I = -dI$.
- Total (desired) net change over a cycle: 0.
- ullet Therefore, net change of P-D over whole cycle is $\Delta_S + \Delta_P + \Delta_I = Q(1-d au) dS dI = 0.$

- Since $I \geq 0$, $Q(1 d\tau) dS \geq 0$.
- If I = 0, $Q(1 d\tau) = dS$.
- If d au>1, net change in P-D will be negative.

Deterministic Example

Production & inventory history

$$S=3,\,Q=10,\, au=1,\,d=.5$$

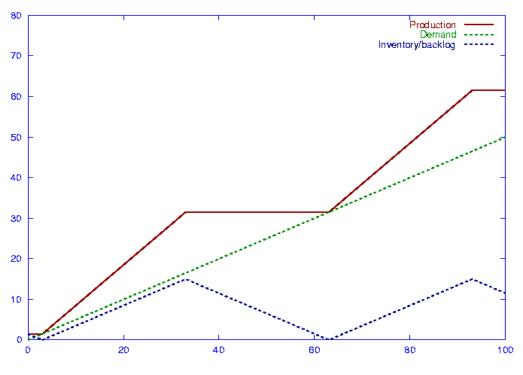


- Production period duration = $Q\tau = 10$.
- Idle period duration = 7.
- Total cycle duration = 20.
- Maximum inventory is Q(1- au d)=5.

Setups

Not frequent enough

$$S=3,\,Q=30,\, au=1,\,d=.5$$

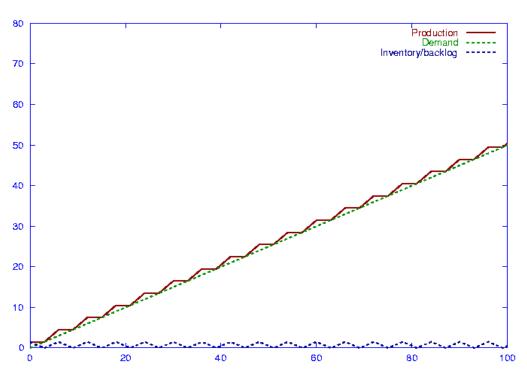


- Large batches big inventories.
- Production period duration $= Q\tau = 30$.
- Idle period duration = 27.
- Total cycle duration = 60.
- Maximum inventory is Q(1- au d)=15.

Setups

Just right!

$$S=3,\,Q=3,\, au=1,\,d=.5$$

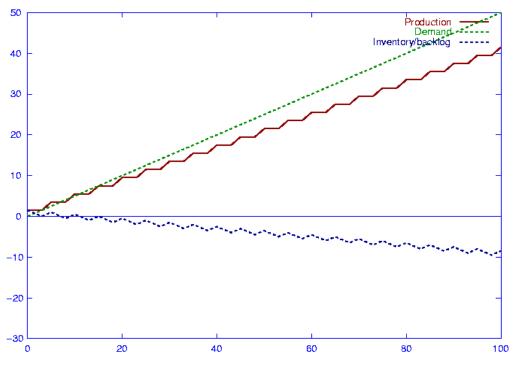


- Small batches small inventories.
- Maximum inventory is $Q(1-\tau d)=1.5$.

Setups

Too frequent

$$S=3,\,Q=2,\, au=1,\,d=.5$$

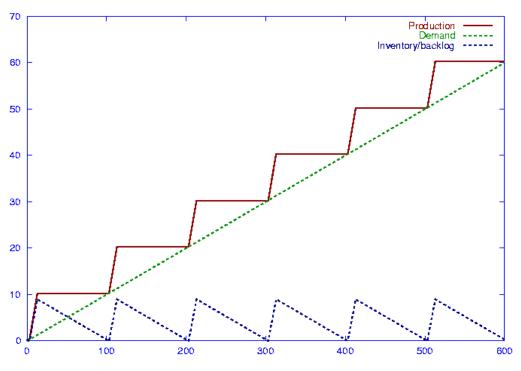


- Batches too small demand not met.
- ullet Q(1-d au)-dS= -0.5
- Backlog grows.
- Too much capacity spent on setups.

Setups

Other parameters

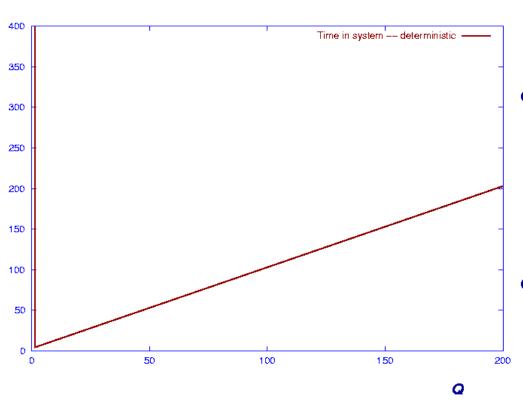
$$S = 3, Q = 10, \tau = 1, d = .1$$



 Not always symmetrical.

Deterministic Example

Time in the system



 Each batch spends Q au + S time units in the system if $Q(1-d\tau)-dS\geq 0.$

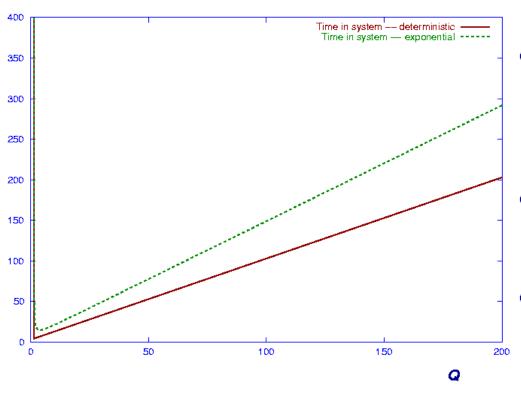
$$Q(1-d au)-dS\geq 0$$

Optimal batch size:

$$Q = dS/(1 - d\tau)$$

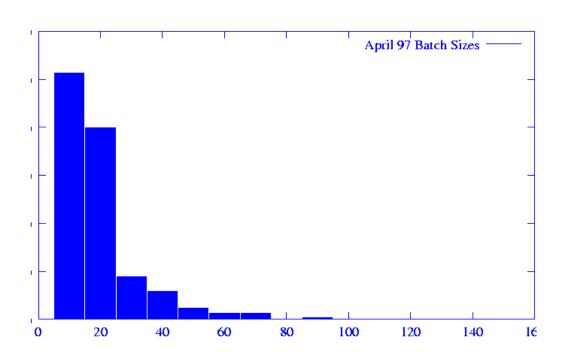
- Batch sizes equal (Q); processing times random.
 - \star Average time to process a batch is $Q\tau+S=1/\mu$.
- Random arrival times (exponential inter-arrival times)
 - \star Average time between arrivals of batches is $Q/d=1/\lambda.$
- Infinite buffer for waiting batches

Stochastic Example

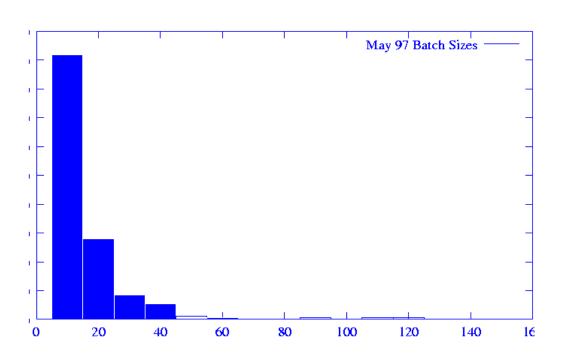


- Treat system as an M/M/1 queue in batches.
- Average delay for a batch is $1/(\mu \lambda)$.
- Variability increases delay

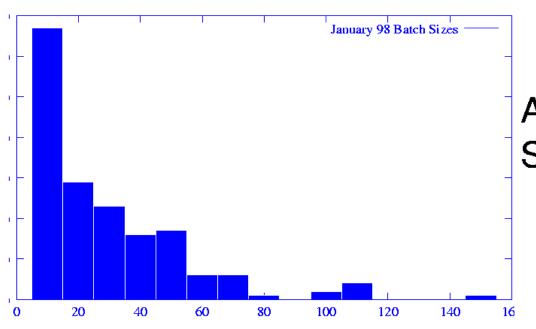
from a factory



from a factory



from a factory



Avg Lot Size=25 Std Dev=27

- Wagner-Whitin (1958) problem
- Assumptions:
 - \star Discrete time periods (weeks, months, etc.); t=1,2,...,T.
 - \star Known, but non-constant demand $D_1, D_2, ..., D_T$.
 - * Production, setup, and holding cost.
 - ★ Infinite capacity.

Dynamic Lot Sizing

Other notation

- c_t = production cost (dollars per unit) in period t
- A_t = setup or order cost (dollars) in period t
- $ullet h_t$ = holding cost; cost to hold one item in inventory from period t to period t+1
- I_t = inventory at the end of period t the state variable
- Q_t = lot size in period t the decision variable

Dynamic Lot Sizing

Problem

minimize
$$\sum_{t=1}^T (A_t \delta(Q_t) + c_t Q_t + h_t I_t)$$
 (where $\delta(Q) = 1$ if $Q \geq 0$; $\delta(Q) = 0$ if $Q = 0$)

subject to

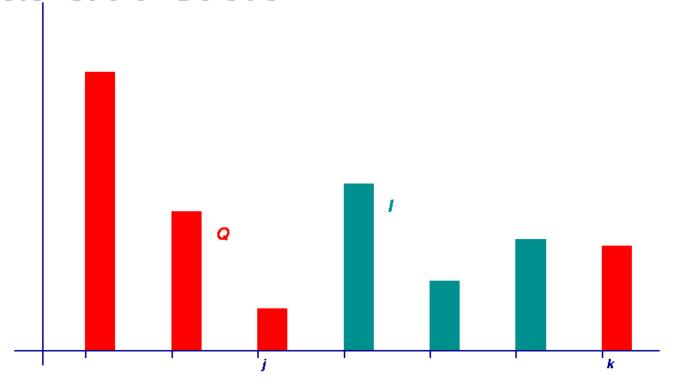
$$ullet I_{t+1} = I_t + Q_t - D_t$$

$$\bullet I_t > 0$$

Dynamic Lot Sizing

Wagner-Whitin Property

Characteristic of Solution:



Dynamic Lot Sizing

Wagner-Whitin Property

Characteristic of Solution:

- Either $I_t = 0$ or $Q_{t+1} = 0$. That is, produce only when inventory is zero. Or,
 - \star If we assume $I_j=0$ and $I_k=0$ (k>j) and $I_t>0, t=j+1,...,k,$
 - \star then $Q_j > 0$, $Q_k > 0$, and

$$Q_t = 0, t = j + 1, ..., k.$$

Dynamic Lot Sizing

Wagner-Whitin Property

Then

$$\bullet \, I_{j+1} = Q_j - D_j,$$

$$ullet I_{j+2} = Q_j - D_j - D_{j+1}, ...$$

$$ullet I_k = 0 = Q_j - D_j - D_{j+1} - ... - D_k$$

Or,
$$Q_j = D_j + D_{j+1} + ... + D_k$$

which means produce enough to exactly satisfy demands for some number of periods, starting now.

Dynamic Lot Sizing

Wagner-Whitin Property

- This is not enough to determine the solution, but it means that the search for the optimal is limited.
- It also gives a qualitative insight.

Real-Time Scheduling

Setups

- Problem: How to decide on batch sizes (ie, setup change times) in response to events.
- Issue: Same as before.
 - * Changing too often causes capacity loss; changing too infrequently leads to excess inventory and lead time.

Real-Time Scheduling

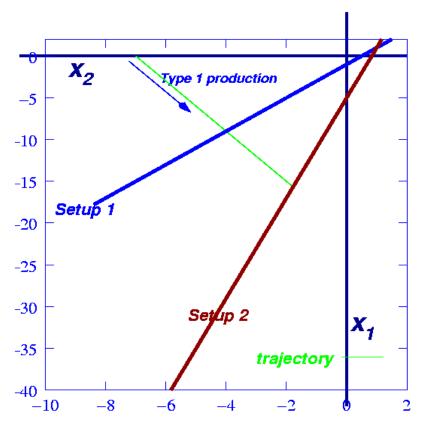
One Machine, Two Part Types

Model:

- d_i = demand rate of type i
- $\mu_i = 1/\tau_i$ = maximum production rate of type i
- S = setup time
- $u_i(t)$ = production rate of type i at time t
- $x_i(t)$ = surplus (inventory or backlog) of type i

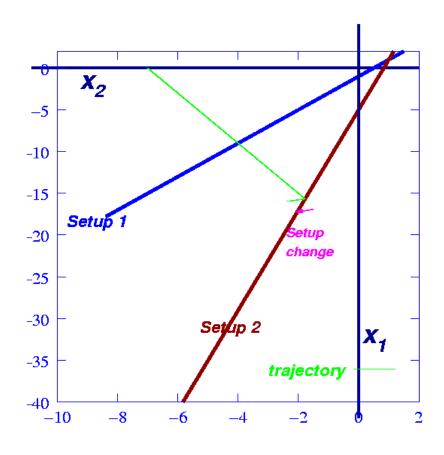
$$ullet rac{dx_i}{dt} = u_i(t) - d_i, i = 1, 2.$$

Real-Time Scheduling

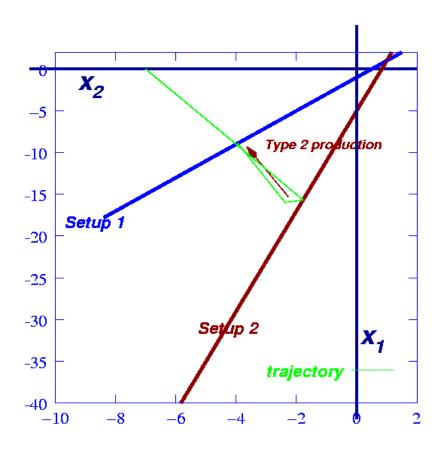


- Draw two lines, labeled
 Setup 1 and Setup 2.
- Keep the system in setup i until x(t) hits the Setup j line.
- Change to setup j.
- Etc.

Real-Time Scheduling

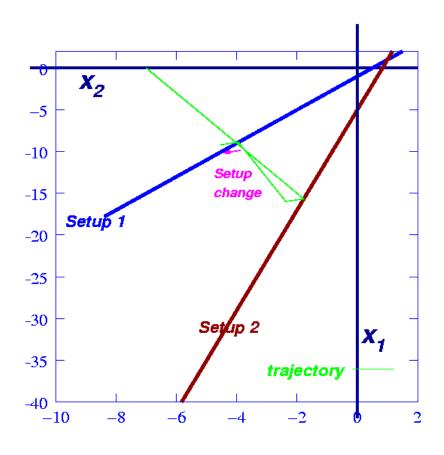


Real-Time Scheduling

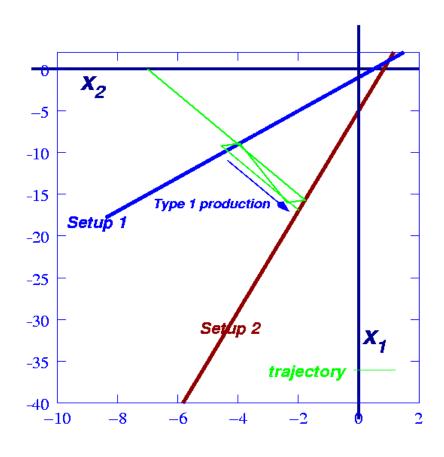


Real-Time Scheduling

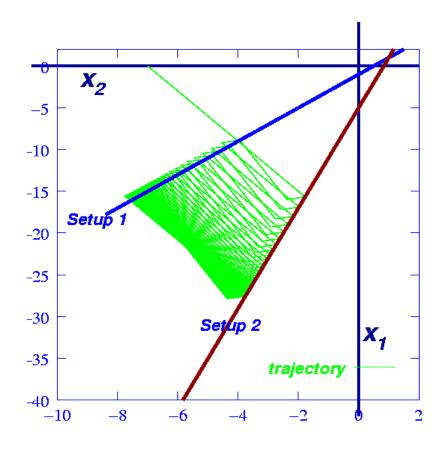
Setups



Real-Time Scheduling



Real-Time Scheduling



Real-Time Scheduling

- In this version, batch size is a function of time.
- Also possible to pick parallel boundaries, with an upper limit. Then batch size is constant until upper limit reached.

Real-Time Scheduling

Corridor Policy

Two possibilities:

- Converges to limit cycle only if demand is within capacity, ie if $\sum_i \tau_i d_i < 1$.
- Diverges if
 - * demand is not within capacity, or
 - * corridor boundaries are poorly chosen.

Real-Time Scheduling

More Than Two Part Types

Three possibilities for the corridor policy:

- Limit cycle only if demand is within capacity,
- Divergence if
 - * demand is not within capacity, or
 - * corridor boundaries are poorly chosen.
- Chaos if demand is within capacity, and corridor boundaries chosen ... not well?