XP
A Common Lisp Pretty Printing System

by

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Abstract

XP provides efficient and flexible support for pretty printing in Common Lisp. Its single greatest advantage is that it allows the full benefits of pretty printing to be obtained when printing data structures, as well as when printing program code.

XP is efficient, because it is based on a linear time algorithm that uses only a small fixed amount of storage. XP is flexible, because users can control the exact form of the output via a set of special format directives. XP can operate on arbitrary data structures, because facilities are provided for specifying pretty printing methods for any type of object.

XP also modifies the way abbreviation based on length, nesting depth, and circularity is supported so that they automatically apply to user-defined functions that perform output—e.g., print functions for structures. In addition, a new abbreviation mechanism is introduced that can be used to limit the total number of lines printed.

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1. Introduction

XP is a portable pretty printer for Common Lisp. As a pretty printer for Lisp code, XP has the advantage that it is fast and allows the user to easily change the way code is formatted. Beyond this, unlike most pretty printers, XP provides comprehensive facilities for pretty printing data structures.

**Four levels of use.** XP can be used at four different levels. At the lowest (weakest) level, you merely have to install XP as described below. This will give you the benefits of more attractive and faster pretty printing. If this is all you want to do, you need not read anything more than this introduction. However, you should take note of the issues raised in the beginning of Section 5.

Moving to a slightly higher level, Section 2 describes several variables that can be used to control XP and some simple additional functionalities provided by XP.

Section 3 exposes the heart of the approach taken by XP and describes a set of new *format* directives that allows users to control the layout of their output dynamically in response to the space available for printing it.

The highest level of using XP is discussed in Section 4. Facilities are provided for specifying how particular types of objects should be pretty printed. These facilities can be used to define (or redefine) the way program code and data structures are displayed.

Section 5 describes a number of shortcomings of XP that stem from the fact that it is supported as a portable package of functions and macros rather than being a standard part of Common Lisp. A number of things cannot be supported quite as they should be, because Common Lisp does not provide all the necessary hooks. Section 5 also describes how proper integration with Common Lisp can be achieved.

**Setting up XP.** To use XP, the file containing it has to be loaded. At the MIT AI Laboratory, XP resides in the file "$b:1mlib>xp.lisp". Compiled versions exist for Symbolics and Lucid (Sun) Common Lisp.

If a compiled version of this file does not exist at your site, one needs to be created. Information about how to get this file via the ARPANET can be obtained by sending a message to the author “dick@ai.mit.edu”. No other method of distribution is available.

All of the functions and forms discussed below are defined in the package "xp". To make these names easily accessible, you must use the package "xp". The most convenient way to do this is to call the function *xp:* :install, which also sets up some additional features of XP. The examples in this manual assume that the form (*xp:* :install) has been evaluated.

```
* xp:* :install &key (:package *package*) (:macro T) (:shadow T) (:remove nil) ⇒ T
```

Calling this function sets up XP for use in the package :package. The argument :package can either be a package, a package name, or a symbol whose name is the name of a package. It defaults to the current package.

to-string, xp::print-to-string, xp::finish-output, xp::force-output, xp::clear-output, and xp::defstruct are shadowing imported into :package. This introduces these functions into :package in place of the standard printing functions. If :shadow is nil you will have to refer to these functions with explicit package prefixes whenever you wish to print something using XP.

If :remove is not nil, the effects of having previously installed XP are undone. In particular, the package is unused and any shadowing is undone. However, any changes to the readtable are left in place.

XP son of PP [7] son of Gprint [5, 6] son of #print is the latest in a line of pretty printers that goes back 12 years. All of these printers use essentially the same basic algorithm and conceptual interface. Further, except for #print, which was not used outside of the MIT AI Laboratory, each of these printers has had extensive use. XP has been in experimental use as the pretty printer in CMU Common Lisp for 8 months. PP has been the pretty printer in DEC Common Lisp for the past 3 years. Prior to three years ago, Gprint was used for 2 years as the pretty printer in DEC Common Lisp. In addition, Gprint has been the pretty printer in various generations of Symbolics Lisp for upwards of 5 years. Both the algorithms and the approach have survived the test of time. A chapter describing the XP interface will appear in the next edition of [4] and this interface is being considered as a formal extension to Common Lisp by the Common Lisp standardization committee.

The basic pretty printing algorithm used by #print, Gprint, PP, and XP has been independently developed by a number of people [2, 3] in addition to the author. This paper does not go into a detailed discussion of the algorithm since it is fully discussed elsewhere (see [3, 5]). However, it should be noted that the algorithm is an inherently fast linear algorithm that uses very little storage. As a result, pretty printing need not be significantly slower than ordinary printing.

The relevance of XP is not restricted to Common Lisp. Many of the ideas discussed in this paper have a wider area of applicability. For instance, the mechanisms for allowing the user to exercise control over the dynamic arrangement of output could be incorporated into almost any programming language (e.g., into the formatted output statements of Fortran or PL/I). Similarly, the mechanisms for linking special pretty printing functions with various types of objects could be introduced into almost any programming environment, given the same kind of type information from the compiler that a good debugger requires.
2. Pretty Printing

Additional Printer Control Variables

XP supports all the standard printer control variables (see [4]). In addition, it supports several new control variables that are relevant to its printing method. None of these new variables has any effect when *print-pretty* is nil.

- **print-dispatch** default value causes standard pretty printing
  This variable holds a printing dispatch table (see Section 4) that controls the way pretty printing is performed. It is initialized to a value that causes traditional Lisp pretty printing. Altering this table, or setting *print-dispatch* to a different table alters the style of pretty printing.

- **print-lines** default value nil
  When not nil, this variable controls the number of lines that will be printed by a single pretty printing event. If an attempt is made to pretty print more than *print-lines* lines, “---” is printed at the end of the last line and pretty printing prematurely halts.

  ```lisp
  (pprint '(setq a b c d))
  With *print-lines* 3 and line width 20 prints:
  (SETQ A 2
       B 3
       C 3 ---
  ```

  Experience has shown that abbreviation based on the number of lines printed can be much more useful than the traditional depth and length abbreviation mechanisms. This is particularly true when the user wants to limit output to a small space. To do this with depth and length abbreviation, the length and depth have to be limited to very small values such as 3 or 4. This often has the unfortunate effect of producing output that consists almost totally of “#”s and “...”s grouped in parentheses. In contrast, limiting the total number of lines printed to 3 or 2, or even 1 produces legible output. Seeing the first few lines of output is usually more informative than seeing only the top level skeletal structure of the output.

- **print-miser-width** default value 40
  A fundamental problem with pretty printing arises when printing deeply nested structures. As line breaks have to be inserted at deeper and deeper levels of nesting, the indentation gets greater and greater. This causes the line width available for printing to get smaller and smaller until it is no longer possible to print substructures in the space available.

  An approach to dealing with this problem, which has been used at least since the original Goldstein pretty printer, [1] is to introduce a special compact kind of layout (called ‘miser’ style) and to use this style once the line width begins to get small. The
key idea behind miser style is that by trading off readability for compactness, it reduces
the width necessary for printing a structure and limits the increase in indentation when
further descending into the structure.

XP switches to an extra compact style of output whenever the width available for
printing a substructure is less than or equal to *print-miser-width*. If *print-miser-
width* is given the value nil, miser style is never used.

A constant default value of 40 for *print-miser-width*, irrespective of the right
margin, is used because the point at which miser style should be triggered does not depend
on the total line width, but rather on the length of the minimal unbreakable
units in the output being printed. When pretty printing programs, the size of these
units depends on the lengths of the symbols in the program. Experience suggests that
*print-miser-width* should be set at from two to four times the length of the typical
symbol.

\*print-right-margin* default value nil

When not nil, this variable specifies the right margin to be used when pretty printing.
By introducing line breaks, XP attempts to prevent any line of output from going beyond
*print-right-margin*.

The left margin for printing is set to the column position where the output begins (or
0 if this position cannot be determined). Except for unconditional newlines (see page 12)
and negative indentation (see page 13), XP will not allow anything but white space to
appear to the left of the left margin.

\*default-right-margin* default value 70

When *print-right-margin* is nil, XP queries the output stream, to determine the
width available for printing. However, there are some situations (e.g., printing to a string)
where the stream has no inherent width limits. In this situation, the right margin is set
to the value of *default-right-margin*, which must be an integer.

Extensions to Output Functions

XP alters the standard printing functions (see the description of xp:install) so that
they operate via XP when (and only when) *print-pretty* is not nil. In addition, XP
provides some additional functionalities

• write object &key :dispatch :right-margin :lines :miser-width ... ⇒ object

The list of keywords accepted by the standard output function write is extended by
adding four more keywords corresponding to the first four control variables above.

A neutral format directive. XP provides a new format directive *W, which bares the
same basic relationship to write that *S and *A bare to print and princ. In particular,
*W prints the corresponding argument obeying every printer control variable, rather than
overriding the value of *print-escape*. In addition, *W interacts correctly with depth
abbreviation, rather than resetting the depth counter to zero. To get the best results
when using XP you should use *W whenever possible instead of *S or *A.
Unlike "S and "A, "W does not accept parameters. If given the colon modifier, "W binds *print-pretty* to T. If given the atsign modifier, "W binds *print-level* and *print-length* to nil.

**Efficient support for formatted output.** The control strings used by format are essentially programs that perform printing. Given that almost all of these strings are constants, there is no reason why they cannot be compiled into highly efficient code. In fact, most languages other than Lisp have always compiled their format-like constructs. However, every implementation of Common Lisp known to the author operates on format control strings in a purely interpretive manner.

XP supports compiled format strings through the combination of two mechanisms. The standard function format (and the directives "? and "{" with no body) are extended so that they can accept functional arguments in addition to standard format control strings. (This recovers the functionality of the original MacLisp format directive "Q.) The functions passed to format must accept the pattern of arguments shown below.

\[
\text{formatting-function stream \&rest args} \Rightarrow \text{unused-args}
\]

When format (or "? or "{" with no body) is called with a functional argument instead of a control string, it applies the function to an appropriate output stream and the appropriate list of arguments. The function should perform whatever output is necessary. In addition, it should return a list of any of the input arguments that it does not use when producing this output. (This is needed for the proper operation of the directives "Q? and "{" with no body.)

A function to be passed to format can be defined in any way desired. However, XP supports a new readmacro character ""..."" that makes it possible to obtain the efficiency of using a function as the second argument to format without loosing the compactness of format control strings. The string following ""#"" in this notation is identical in every respect to a format control string. The readmacro merely translates it into an equivalent function that follows the conventions discussed above. (As in any string, any instances of """" in the delimited region must be preceded by ""T".) Note that ""..."" is not limited to appearing only as the second argument to format. It can appear anywhere—e.g., passed to some other function that will eventually pass it to format or used for some totally different purpose.

\[
\text{(format T "\%Prices: S", "\{" 1.50 3.23 4.50) \equiv}
\]
\[
\text{(format T "'(lambda (stream \&rest args)
\terpri stream)
\write-string "Prices: " stream)
\loop (prini (pop args) stream)
\(\text{if (null args) (return nil))
\write-string "", " stream))
\args)
\1.50 3.23 4.50)}
\]

Using compiled format control strings instead of interpreted ones approximately triples the speed of XP when running on the Symbolics Lisp Machine.
3. Dynamic Control of the Arrangement of Output

Through the introduction of several additional format directives, XP allows the user to exercise precise control over what should be done when a piece of output is too large to fit in the line width available for displaying it. The discussion below assumes that the reader has a basic understanding of the function format and the standard format directives, as described in [4].

Basic concepts. Three concepts underlie the way XP supports the dynamic arrangement of output—logical blocks, conditional newlines, and sections. The positions of logical blocks and conditional newlines are specified by means of format directives. Together, these directives divide the output as a whole into a hierarchy of sections within sections.

The first line of Figure 3.1 shows a schematic piece of output. The characters in the output are represented by “-”s. The positions of conditional newlines are indicated by digits. The beginnings and ends of logical blocks are indicated by “<” and “>” respectively.

The output as a whole is always treated as a logical block and forms one (outermost) section. This section is indicated by the 0’s on the second line of Figure 3.1. Each conditional newline defines two additional sections (one before it and one after it) and is associated with a third (the section immediately containing it).

The section after a conditional newline consists of: all the output up to, but not including, (a) the next conditional newline immediately contained in the same logical block; or if (a) is not applicable, (b) the next newline that is at a lesser level of nesting in logical blocks; or if (b) is not applicable, (c) the end of the output as a whole.

The section before a conditional newline consists of: all the output back to, but not including, (a) the previous conditional newline that is immediately contained in the same logical block; or if (a) is not applicable, (b) the beginning of the immediately containing logical block. The last four lines in Figure 3.1 indicate the sections before and after the four conditional newlines.

The section immediately containing a conditional newline is the shortest section that contains the conditional newline in question. In Figure 3.1, the first conditional newline is immediately contained in the section marked with 0’s, the second and third conditional newlines are immediately contained in the section before the fourth conditional newline, and the fourth conditional newline is immediately contained in the section after the first conditional newline.

```
<-1--<-><-2--3--4-->  
000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000000
It also makes sense to talk about the section immediately containing a given logical block—i.e., the shortest section containing the logical block. Note that this section immediately contains every conditional newline that is immediately contained in the block. In Figure 3.1, the outermost logical block is immediately contained in the section marked with 0's, the second logical block is immediately contained in the section before the fourth conditional newline, and the innermost logical block is immediately contained in the section after the first conditional newline.

Whenever possible, XP prints the entire contents of a given section on the same line. However, if a section is too long to fit in the line width available, XP inserts line breaks at one or more conditional newline positions within the section—printing the section on more than one line. The pretty printing algorithm uses internal buffering of the output so that it can determine which way to print a section. The algorithm is fast, because the amount of lookahead required is strictly limited by the maximum line width available for printing.

**Logical Blocks**

If `"::"` is used to terminate a `"<...>"`, the directive delimits a logical block. In addition, the directive descends into the corresponding format argument (which should be a list) in the same way as the standard directive `"1{...}"` (iterate once over list).

```lisp
(format T #"+ "<Roads "<"W, "::"W:"> ":: Town "<"W:">":" +"
'((elm cottonwood) (boston)))
```

With line width 50 prints:

+ Roads ELM, COTTONWOOD Town BOSTON +

With line width 25 prints:

+ Roads ELM, COTTONWOOD Town BOSTON +

With line width 21 prints:

+ Roads ELM,
  COTTONWOOD
  Town BOSTON +

(As discussed in the next section, the directive `"::_` indicates a conditional newline. An instance of `"::_` is replaced by a line break when the following section cannot fit on the end of the current line. Whenever line breaks are introduced, indentation is also introduced so that each line in a logical block begins in the same column as the logical block as a whole.)

If the atsian modifier is used with `"<...>:"`, the directive operates on the remaining format arguments in the same way as the standard directive `"10{...}"` (iterate once over remaining arguments), except that all of the remaining arguments are always consumed by the `"<...>:"`, whether or not they are actually used by the format string nested in the directive. The directive `"-` (termination test) can be used to exit from `"<...>:"` just as it can be used to exit from `"{...}"`. 
Dynamic Control of the Arrangement of Output

```lisp
(format T #"+ "'@<Roads '"W" ';"'W":> "': Town '"W":>""> +"
'(elm) '(boston))
```

With line width 21 prints:
```
+ Roads ELM
  Town BOSTON +
```

The portion of a format control string enclosed in a "<...>:" directive can be divided into segments "<prefix";body";suffix"> by ";" directives. It is an error for the enclosed portion to be divided into more than three segments. If the portion is only divided into two segments, the suffix defaults to the null string. If the portion consists of only a single segment, both the prefix and the suffix default to the null string. The prefix and suffix must both be constant strings. They cannot contain format directives. The body can be any arbitrary format control string.

When a "<prefix";body";suffix"> directive is processed, the prefix is printed out just before the logical block begins and the suffix is printed out just after the logical block ends. This behavior is the same as if the characters in the prefix and suffix simply appeared before and after the "<...>:" directive, except for the way error situations are handled and the way "<...>:" interacts with depth and circularity abbreviation (see page 17).

```
(format T #"+ "'<Roads '";'W Town '<;'W--1-.->-:> +11
'(('elm cottonwood) boston))
```

With line width 21 prints:
```
+ Roads [ELM
  COTTONWOOD]
  Town BOSTON +
```

An interesting additional feature of "<...>:" is illustrated by the example above. When a "<...>:" directive is applied to an argument that is not a list, the directive is ignored and the offending argument is printed using "W. Among other things, this means that while the argument is printed, the prefix and suffix are not. The soft failure of "<...>:" when presented with non-lists makes it possible to write robust format strings that produce reasonable output for a wide range of possible arguments. This is particularly useful in debugging situations.

During the processing of the format string nested in "<...>:", arguments are taken one by one from the list passed to "<...>:". If an attempt is made to access an argument at a time when the remaining portion of this argument list is not a cons, then "." is inserted in the output, "W is used to print out the remaining argument list, and the processing of the logical block is terminated, except for printing the suffix (if any). This makes it easier to write format strings that are robust in the face of malformed argument lists. (Note that "" exits only when the remaining argument list is nil.)

```
(format T #"+ "'<Roads '";'W "':"'W":> "': Town '";'W":>""> +"
'(('elm . cottonwood) boston))
```

With line width 21 prints:
```
+ Roads [ELM
  . COTTONWOOD]
  Town BOSTON +
```
If the colon modifier is used with "<...:">, the prefix and suffix default to "(" and ")" (respectively) instead of to the null string. Note that the prefix and suffix are printed out even when the argument corresponding to "<...:" is an empty list.

```
(format T "+ " "Roads ":"W ":"W":> ": Town ":"W":>" +" 
'((elm cottonwood) ()))
```

With line width 21 prints:
+ Roads (ELM
     COTTONWOOD)
   Town () +

If the directive "@;" is used to terminate the prefix in a "<...:" directive, the prefix is treated as a per-line prefix. A per-line prefix is printed at the beginning of every line in the logical block, rather than just before the start of the block as a whole. This is done in such a way that the prefixes on subsequent lines are directly below the occurrence of the prefix on the first line.

```
(format T "" "@;Roads " "(elm cottonwood) (boston))"
```

With line width 50 prints:
@; Roads = ELM, COTTONWOOD Town BOSTON

With line width 25 prints:
@; Roads = ELM,
@; = COTTONWOOD
@; Town BOSTON

If a "<...:" directive is terminated with "::", then a ":" is automatically inserted after each group of blanks immediately contained in the body (except for blanks after a "<newline> directive). This makes it easy to achieve the equivalent of paragraph filling.

```
(format T "" "<:("W") street goes to " (:"W") ::" )=" (main boston)" 
(format T "" "<:("W") ":_street ":_goes ":_to ":_(:"W") ::" )=" 
 ' (main boston))
```

With line width 12 prints:
Main street
goesto
Boston.

To a considerable extent, the basic form of the directive "<...:" is incompatible with the dynamic control of the arrangement of output by "W", "_", "<...:">, "I", and ":T. As a result, it is an error for any of these directives to be nested within "<...:". (Note that in standard format, it makes little sense to have anything that can cause a line break within a "<...:" and [4] does not define what would happen if this were the case.)

Beyond this, it is also an error for the "<...:";...:" form of "<...:" to be used at all in conjunction with any of these directives. This is a functionality that is rendered obsolete by per-line prefixes in "<...:". 

 Logical Blocks
Conditional Newlines

The format directive _- is used to specify conditional newlines. By means of the colon and atsign modifiers, the directive can be used to specify four different criteria for the dynamic insertion of line breaks: linear-style, fill-style, miser-style, and mandatory-style.

By default, whenever a line break is inserted by a conditional newline, indentation is also introduced so that the following line begins in the same column as the first character in the immediately containing logical block. This default can be changed by using the directive "I (p. 13).

Linear-style conditional newlines. Without any modifiers, _- specifies linear-style insertion of line breaks. This style calls for the subsections of a logical block to be printed either all on one line or each on a separate line.

Linear-style conditional newlines are replaced by line breaks if and only if the immediately containing section cannot be printed on one line. As soon as the line width available becomes less than the length of a given section, every linear-style conditional newline in it is replaced by a line break.

\[
\text{(format T \\
\"\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\\)}
\[
\text{With line width} \geq 25 \text{ prints:} \n\text{(LIST FIRST SECOND THIRD)}
\]

With line width < 25 prints:
\[
\text{(LIST FIRST SECOND THIRD)}
\]

Miser-style conditional newlines. If the atsign modifier is used with _, the directive specifies miser-style insertion of line breaks. Miser-style conditional newlines are replaced by line breaks if and only if miser style is in effect in the immediately enclosing logical block and the immediately containing section cannot be printed on one line. Miser style is in effect for a given logical block if and only if the starting column of the logical block is less than or equal to *print-miser-width* columns from the end of the line. (Note that the \("C\) in the example below is not in the logical block, but rather before it.)

\[
\text{(format T \\
\"\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\:\\)}
\[
\text{With line width} 10 \text{ and } *\text{print-miser-width}* < 9 \text{ prints:} \n\text{(LIST FIRST SECOND THIRD)}
\]

With line width 10 and \(*\text{print-miser-width}* \geq 9 \text{ prints:}
\[
\text{(LIST FIRST SECOND THIRD)}
\]

Even in miser style, the pretty printing algorithm is not guaranteed to succeed in keeping the output within the line width available. In particular, line breaks are never
inserted except at conditional newline positions. As a result, a given output requires a
certain minimum amount of line width to print it. If the amount of line width available
is less than this amount, characters are printed beyond the end of the line.

(GPRINT [5] supported an additional mechanism for dealing with deeply nested struc-
tures. When indentation reduced the line width to a small percentage of its initial value,
major program structures (such as prog and let) were radically shifted to the left by re-
ducing the indentation to nearly zero. This violated standard Lisp pretty printing style,
but significantly increased the line width available for printing. Unfortunately, experi-
ence showed that, though this was very useful in some situations, it was, in general, more
confusing than helpful.)

**Fill-style conditional newlines.** If the colon modifier is used with `:`, the directive
specifies fill-style insertion of line breaks. This style calls for the subsections of a section
to be printed with as many as possible on each line.

Fill-style conditional newlines are replaced by line breaks if and only if either (a)
the following section cannot be printed on the end of the current line, (b) the preceding
section was not printed on a single line, or (c) the immediately containing section cannot
be printed on one line and miser style is in effect in the immediately containing logical
block. If a logical block is broken up into a number of subsections by fill-style conditional
newlines, the basic effect is that the logical block is printed with as many subsections as
possible on each line. However, if miser style is in effect, fill-style conditional newlines
act like linear-style conditional newlines.

For instance, consider the example below. The format control string shown uses the
standard directives `@{(...)}` (iterate over arguments) and `~` (terminate iteration) to
decompose the list argument into pairs. Each pair in the list is itself decomposed into
two parts using a `<(...)~>` directive. A space and a `:_` are placed after each pair except
the last.

```
(format T #"("LET `:~"@{:"~"W:"~"W~":"~":"~"}~":"~"~"~:"~"~"~:"~"~"")"
  '((x 4) (*print-length* nil) (z 2) (list nil)))
```

With line width 35 prints:

```
(LET ((x 4) (*PRINT-PRETTY* NIL)
       (z 2) (LIST NIL))
    ...
)
```

With line width 22 prints:

```
(LET ((x 4)
       (*PRINT-LENGTH*
        NIL)
       (z 2)
       (LIST NIL))
    ...
)
```

Note that when the line width is 35, only one line break (before "(Z 2 )") has to be
introduced. Once this is done, "(Z 2 )" and "(LIST NIL)"
both fit on the next line. Also note that when the line width is 25, there is a line break after "NIL )", even though
"(Z 2 )" would fit on the end of the previous line. The line break is introduced, due to
criteria (b) in the definition above. This criteria is used, because many people judge that
output of the following form violates the basic aesthetics of Lisp pretty printing.
(LET ((X 4)
      (*PRINT-LENGTH*
        NIL) (Z 2)
      (LIST NIL))
  ...
)

It is often useful to mix different kinds of conditional newlines together in a single logical block. In general, this works well without any conflicts arising between the ways the various conditional newlines work. However, it is not a good idea to have a miser-style conditional newline immediately after a fill-style conditional newline. The problem is that the miser-style conditional newline will terminate the section following the fill-style conditional newline. As a result, no account will be taken of what follows the miser-style conditional newline when deciding whether or not to insert a line break at the fill-style conditional newline. This can cause the output after the miser-style conditional newline to extend beyond the end of the line.

**Mandatory-style conditional newlines.** If the colon and atsign modifiers are both used with `-`, the directive specifies mandatory insertion of a line break. Among other things, this implies that none of the containing sections can be printed on a single line. This will trigger the insertion of line breaks at linear-style conditional newlines in these sections. With regard to indentation, mandatory-style conditional newlines are treated just like any other kind of conditional newline. This makes them different from unconditional newlines.

**Unconditional newlines.** There are at least six ways to introduce a newline into the output without using `-`. You can use the directives `-%` and `-&. You can include a newline character in a format string. You can call `terpri` or `print`. You can print a string that contains a newline. Each of these methods produces an unconditional newline that must always appear in the output. Like a mandatory conditional newline, this prevents any of the containing sections from being printed on one line.

It is not completely obvious how unconditional newlines should be handled. There are two very suggestive cases, which unfortunately contradict each other. First, suppose a program containing a string constant containing a newline character is printed out and then read back in. To ensure that the result of the read will be equal to the original program, it is important that indentation not be inserted after the newline character in the string. On the other hand, suppose that the same program is being printed into a file by a format string specifying a per-line prefix of `";;;"`, with the intention that the program appear in the output as a comment. To ensure that this comment will not interfere with subsequent reading from the file, it is important that the prefix be printed when the newline in the string is printed.

In an attempt to satisfy the spirit of both of the cases above, XP applies the following heuristic. Indentation is used only if a newline is created with `-`.

```
(format T ""-\:LIST "\%\"\-\:"\="" (first "string on
two lines")
)
```

With line width 100 prints:
```
(LIST FIRST
   "string on
    two lines")
```

However, any per-line prefixes (and any indentation preceding them) are always printed no matter how a newline originates.

\[(\text{format } T \#""Q<;; ;""Q;:"<\text{LIST } ""Q<"W "¬"W":>:">:">" (first "string on two lines")))\]

With line width 100 prints:

```
;;; (\text{LIST FIRST}
;;; "string on
;;; two lines")
```

**Discarding trailing spaces.** Conditional newline directives are typically preceded by some amount of blank space. This is done so that the subsections of a section will be visually separated when they are printed on a single line. However, without anything more being said, this would lead to the printing of unnecessary blank spaces at the end of most lines when line breaks are inserted. In the interest of efficiency, XP suppresses the printing of blanks at the end of a line if (and only if) the line break was caused by a "_" directive. For instance, there are no blanks at the end of any of the lines in the examples above, except that, if there are blanks after the word "on" in the string in the two unconditional newline examples, these blanks appear in the output.

**Indentation**

By default, the second and subsequent output lines corresponding to a logical block are indented so that they line up vertically under the first character in the block. The `format` directive "I makes it possible to specify a different indentation.

The directive "nI specifies that the indentation within the immediately containing logical block should be set to the column position of the first character in the block plus n. If omitted, n defaults to zero. The parameter can be negative, which will have the effect of moving the indentation point to the left of the first character in the block. However, the total indentation cannot be moved left of the beginning of the line or left of the end of the rightmost per-line prefix.

The directive "n:I is exactly the same as "nI except that it operates relative to the position in the output of the directive itself, rather than relative to the position of the first character in the block.

As an example of using "I, consider the following: If the line width available is 25, the "_" directive is replaced by a line break. The "I directive specifies that the statement in the body of the `defun` should be printed at a relative indentation of 1 in the logical block. If the line width is 15, the "_" is also replaced by a line break. The ":I directive before the "W printing the function name causes the argument list to be lined up under the function name. The column position corresponding to the ":I is determined dynamically as the output is printed.

\[(\text{format } T ":"<"W "Q_""I""W ":"<"W "I""W":>:">" 
  '(\text{defun prod (x y) (* x y))}))\]

With line width 50 prints:

```
(\text{DEFUN PROD } (X Y) (* X Y))
```
Dynamic Control of the Arrangement of Output

Pretty Printing as Selection

Stepping back a moment, it is useful to reflect on how the format directives above interact to support pretty printing. The "<...:" and "_" directives in a format control string divide the output up into a hierarchy of sections within sections. The "_" and "I" directives simultaneously specify three ways (on one line, on multiple lines, and in miser style) for printing each section. The job of the pretty printer boils down to selecting (based on the length of the section, the line width available, and the value of *print-miser-width*) which of the three ways is appropriate for printing each section.

The format directives have been designed so that it is relatively easy to specify three different ways to print a logical block in a single format control string. In particular, except for line breaks, white space, (and per-line prefixes), the characters to be printed are exactly the same in each of the three styles. The "_" and "I" directives specify how the logical block sections are to be arranged when printed on multiple lines and in miser style.

Many other kinds of pretty printing directives could have been supported—for example, arbitrary sections of text that are output only when a section is printed on multiple lines. XP supports only the limited set of directives above, because experience has shown them to be a good compromise between the requirements of expressive power, easy understandability, and efficiency.

Tabbing Within a Section

The standard format directive "T is extended so that it supports the colon modifier in addition to the atsign modifier. If the colon modifier is specified, the tabbing computation
is done relative to the beginning of the immediately containing section, rather than with respect to the beginning of the line. (When this computation is performed, any unconditional newlines in the section are ignored.) As an example of using \( \ast :T \), consider the following. Each street name is followed by a \( \ast 8:T \), which ensures that the total width taken up will be a multiple of 8. Fill-style conditional newlines are used to put as many streets as possible on each line.

\[
(\text{format} \ T \ #\"-\>:\text{Roads} \ -I-\@\{\"W\}\-8:T\-_:\}\>:\"
\'(\text{elm main maple center}))
\]

With line width 25 prints:

Roads ELM MAIN
MAPLE CENTER

The fact that \( \ast :T \) operates solely within the immediately containing section means that the number of spaces to insert is independent of whatever indentation is in effect. In the example above, a column spacing of 8 is used, but the entire table is shifted over 6 columns. (Note the way the \( \ast \@ \) delimits the beginning of the section containing the first road.)

(The fact that \( \ast :T \) operates solely within the immediately containing section and ignores unconditional newlines means that the amount of space to insert can be determined before deciding which (if any) conditional newlines have to be replaced with line breaks. This is essential for the efficient operation of the pretty printing algorithm.)

When the normal directive \( \ T \) (without a colon) is processed, tabbing is computed relative to the beginning of the line and all conditional newlines are ignored. (Again, this is important so that the number of spaces to insert can be determined before making any decisions about conditional newlines.)

As a practical matter, one should not use \( \ast T \) after a conditional newline nor \( \ast :T \) after an unconditional newline.

### User-Defined Format Directives

XP provides a mechanism for allowing the user to define new format directives. In a somewhat simplified form, this revives a feature of format that was left out when Common Lisp was initially developed.

XP allows a format string to contain a directive of the form \( \ast /name/ \). When this is the case, it is assumed that a function named name has been defined. This function must accept the pattern of arguments shown below. (Name cannot contain any instances of "/". In addition, the "/"'s must be used even if the name only has one character. Among other things, this means that this mechanism cannot be used to redefine the standard format directives.)

\[
\text{name stream arg colon? atsign? &rest parameters } \Rightarrow \text{ignored}
\]

The colon modifier, the atsign modifier, and arbitrarily many parameters can be specified with the \( \ast /name/ \) directive. This information, along with the output stream and one argument from the current argument list, are passed to the function name. The input colon? is \( \ T \) if and only if the colon modifier was specified. The input atsign? is \( \ T \) if
and only if the atsign modifier was specified. The function \texttt{name} should perform whatever operations are required to print \texttt{arg} into \texttt{stream}. Any value returned by the function is ignored.

**Packages.** A key problem with \texttt{/name/} directives derives from the fact that, as written by the user, \texttt{name} is a string, not a symbol. This string has to be converted into a symbol to identify the function. The question is, what package should this symbol be in? You can write a directive of the form \texttt{/package:name/}, in which case the package is explicitly specified and there is no problem. However, if no package name is specified, a default package has to be chosen.

In Common Lisp, symbols without explicit package prefixes are placed in the package that is contained in the variable \texttt{*package*} at the moment when the symbol is first read. To continue this policy, \texttt{name} should be placed in the package that is contained in the variable \texttt{*package*} at the moment when the format string is first read. If a \texttt{/name/} directive appears in a format string specified using \texttt{#"..."}, then this correct behavior is obtained.

Unfortunately, if \texttt{/name/} appears in a format string specified using simply \texttt{"..."}, then \texttt{name} is placed in the package that is contained in the variable \texttt{*package*} at the moment when the format string is first evaluated. It is very possible that this package will not be the same as the one in effect when the string was read. As a result, it is advisable to use \texttt{#"..."} whenever using \texttt{/name/} directives.

**Special format directives for lists.** XP provides three special format directives for printing lists. These are defined and accessed using the mechanisms for user-defined format directives described above.

The directive \texttt{/linear-style/} prints out the elements of a list either all on one line or each on a separate line. Parentheses are printed around the list if the colon modifier is specified.

\begin{verbatim}
(defun linear-style (stream list &optional (colon? T) atsign?)
  (declare (ignore atsign?))
  (if colon?
      (format stream #"\(\text{list}\))
      (format stream #"\<\(\text{list}\))
  )
  (format T #"/:linear-style/" '(one two three))
\end{verbatim}

With line width 15 prints:

\texttt{(ONE TWO THREE)}

With line width 14 prints:

\texttt{(ONE TWO THREE)}

The directive \texttt{/fill-style/} prints out the elements of a list with as many elements as possible on each line. Except for the fact that it uses \texttt{\_} instead of \texttt{\texttt{\_}}, it is identical to \texttt{/linear-style/}. 


Abbreviation

\[(\text{format T \#"/fill-style/" '\{one two three four five\}})\]

With line width 25 prints:
ONE TWO THREE FOUR FIVE

With line width 15 prints:
ONE TWO THREE FOUR FIVE

The directive \"/tabular-style/\ is similar to \"/fill-style/\, except that it prints the elements of the list so that they line up in a table. In addition to the colon modifier, \"/tabular-style/\ takes a parameter (default 16) that specifies the width of columns in the table.

\[(\text{format T \#"8:/tabular-style/" '\{one two three four five\}})\]

With line width 20 prints:

\begin{verbatim}
(ONE TWO THREE FOUR FIVE)
\end{verbatim}

Abbreviation

XP fully supports abbreviation controlled by \texttt{*print-level*}, \texttt{*print-length*}, and \texttt{*print-circle*}. In addition, (see Section 2) XP supports a new abbreviation mechanism that limits the total number of lines printed. All four mechanisms are supported in such a way that they automatically apply to user-defined functions that perform output.

**Depth abbreviation.** XP obeys \texttt{*print-level*} in its internal operation. In addition, it makes it very easy to write format control strings that obey \texttt{*print-level*}. This is done by basing depth abbreviation on the concept of logical blocks. Whenever a \"<...:\> directive is encountered at a dynamic nesting depth in logical blocks greater than \texttt{*print-level*}, \"#\ is printed instead of the block. In addition, the argument (or for \"@<...:\> arguments) that would have been consumed by the directive are skipped.

The following example illustrates how \"<...:\> supports depth abbreviation. The most important feature of the example is that it shows that depth abbreviation is controlled by the \textit{dynamic} nesting of \"<...:\> directives, not their \textit{static} nesting. In the second output shown, the statically outermost instance if \"<...:\> in \"/linear-style/\ (p. 22) is at a dynamic nesting depth of 3. (Note that since there is an implicit logical block dynamically wrapped around the entire output, the dynamically outermost instance of \"<...:\> is at a dynamic nesting depth of 1.)

\[(\text{format T \#"+:<<W ":<<W ":/linear-style/":>>":> +" '\{1 (2 (3))\}})\]

With \texttt{*print-level* nil} prints:
+ (1 (2 (3))) +

With \texttt{*print-level* 2} prints:
+ (1 (2 #)) +
Length abbreviation. XP obeys *print-length* in its internal operation. In addition, it makes it very easy to write format control strings that obey *print-length*. This is done by basing length abbreviation on the concept of logical blocks.

"<..." *> provides automatic support for length abbreviation. If *print-length* is not nil, a count is kept of the number of arguments used within the "<...":>. If this count ever reaches *print-length*, "..." is inserted in the output and the processing of the logical block is terminated, except for printing the suffix (if any). As with depth abbreviation, the processing depends on dynamic relationships, not static ones.

(format T #"+ "#:<@"*"-">:";}+" '(1 2 3 4 5))

With *print-length* nil prints:
+ (1 2 3 4 5) +

With *print-length* 2 prints:
+ (1 2 ...) +

Circularity abbreviation. XP obeys *print-circle* in its internal operation. In addition, it makes it very easy to write format control strings that obey *print-circle*. This is done by supporting circularity abbreviation through the combined actions of "W and "<..." ".

In situations where *print-circle* is not nil, the following extra processing is performed. When "W is applied to a non-list, a check is made to see whether the argument has previously been encountered. If so, an appropriate #n# marker is printed out instead of printing the argument. Similarly, when a "<..." :> is applied to a list, a check is made to see whether the list has previously been encountered.

In addition, if an attempt is made to access an argument from the list passed to "<..." :>, at a time when the remaining portion of this list has already been encountered during the printing process, " . #n#" is inserted in the output and the processing of the logical block is terminated, except for printing the suffix (if any). This catches instances of cdr circularity in lists.

(format T #"+ "#:<@"*"-">:";}+" '#1=(1 2 #1# 3 . #1#))

With *print-circle* T prints:
+ #1=(1 2 #1# 3 . #1#) +

With *print-circle* T and *print-length* 2 prints:
+ (1 2 ...) +

Circularity detection is an inherently slow process. In particular, two entire passes have to be made over the output: one to determine what #n= markers should be printed and another to perform the actual printing. All and all, setting *print-circle* to T more than doubles the time required for printing using XP and should be avoided unless strictly necessary. In the interest of efficiency, XP does not print circularity abbreviation markers in situations where other abbreviation methods hide the circularity. This is illustrated in the last part of the example above.

For a format string to correctly support circularity abbreviation, every part of the object being printed must be seen by an occurrence of "W or "<..." :>. (If some part is
skipped (e.g., printed with "s), XP will fail to detect circularities involving that part.) (The above criteria are also required for depth and length abbreviation to be handled in a completely correct way.)

Reprinting an abbreviated object. XP keeps track of the last pretty printing event that lead to abbreviation due to *print-level*, *print-length*, or *print-lines*. A hook, is provided for obtaining this information. Using this hook, mechanisms can easily be implemented for reprinting abbreviated objects in full (see page 32).

• *last-abbreviated-printing*

   This variable records the last printing event where abbreviation occurred. Functioning its value (e.g., after turning off abbreviation) causes the printing to happen a second time.

Functional Interface

The primary interface to XP’s operations for dynamically determining the arrangement of output is provided through format. This is done, because format strings are typically the most convenient way of interacting with XP. However, XP’s operations have nothing inherently to do with format per se. In particular, they can also be accessed via the six functions and macros below.

• within-logical-block (stream-symbol list &key :prefix :per-line-prefix :suffix)
  &body body ⇒ nil

   In the manner of "<...>" this macro causes printing to be grouped into a logical block. The value nil is always returned.

   The argument stream-symbol must be a symbol. If it is nil, it is treated the same as if it were *standard-output*. If it is T, it is treated the same as if it were *terminal-io*. The run-time value of stream-symbol must be a stream (or nil meaning *standard-output* or t meaning *terminal-io*). The logical block is printed into this destination stream.

   Within the body, stream-symbol is bound to a special kind of stream that supports dynamic decisions about the arrangement of output and then forwards the output to the destination stream. All and only the output sent to stream-symbol is treated as being in the logical block. (It is an error to send any output directly to the underlying destination stream.)

   The :suffix, :prefix, and :per-line-prefix must all be expressions that (at run time) evaluate to strings. The argument :suffix (which defaults to the null string) specifies a suffix that is printed just after the logical block. The argument :prefix specifies a prefix to be printed before the beginning of the logical block. If the argument :per-line-prefix is supplied, it specifies a prefix that is printed before the block and at the beginning of each new line in the block. It is an error for :prefix and :pre-line-prefix to both be supplied. If neither is supplied, a :prefix of the null string is assumed.

   The argument list is interpreted as being a list that the body is responsible for printing. If list is not a list, it is printed using write on stream-symbol and the body is
skipped along with the printing of the prefix and suffix. If *print-circle* is not nil and list is a cons that has already been printed by or within a dynamically containing logical block, then an appropriate #n# marker is printed on stream-symbol and the body is skipped along with the printing of the prefix and suffix. (If the body is not responsible for printing a list, then the behavior above can be turned off by supplying nil for the list argument.)

If *print-level* is not nil and the logical block is at a dynamic nesting depth of greater than *print-level* in logical blocks, "#" is printed on stream-symbol and the body is skipped along with the printing of the prefix and suffix.

The body can contain any arbitrary Lisp forms. All the standard printing functions (e.g., write, princ, terpri) can be used to print output into stream-symbol. Within a logical block, these functions interact correctly with *print-circle* and *print-depth*.

From the above, it can be seen that within-logical-block supports all of the functionality of "<...>" except for the automatic introduction of fill-style conditional newlines supported by "<...>:0>. This feature is omitted, because it is a transformation on format strings rather than a printing operation.

- **conditional-newline** kind &optional (stream *standard-output*) => nil
  The function conditional-newline supports the functionality of _. The stream argument (which defaults to *standard-output*) follows the standard conventions for stream arguments to printing functions (i.e., nil can be used to mean *standard-output* and T can be used to mean *terminal-io*). If stream is a special stream bound by within-logical-block, a conditional newline is sent to stream. Otherwise, conditional-newline has no effect. The value nil is always returned.
  
  The kind argument specifies the style of conditional newline. It must be one of :linear (_), :fill (:_), :miser (_), or :mandatory (:0).

- **logical-block-indent** relative-to n &optional (stream *standard-output*) => nil
  
  The function logical-block-indent supports the functionality of -I. The stream argument (which defaults to *standard-output*) follows the standard conventions for stream arguments to printing functions. If stream is a special stream bound by within-logical-block, logical-block-indent specifies the indentation within the innermost enclosing logical block. Otherwise, logical-block-indent has no effect. The value nil is always returned.
  
  The argument n specifies the amount of indentation. If relative-to is :block, this indentation is relative to the start of the enclosing block (as for -I). Alternatively, if relative-to is :current, the indentation is relative to the current output position in the immediately containing section (as for -:I). It is an error for relative-to to take on any other value.

- **logical-block-tab** kind colnum colinc &optional (stream *standard-output*) => nil
  
  The function logical-block-tab supports the functionality of -T. The stream argument (which defaults to *standard-output*) follows the standard conventions for stream arguments to printing functions. If stream is a special stream bound by within-logical-
block, tabbing is performed. Otherwise, \texttt{logical-block-tab} has no effect. The value \texttt{nil} is always returned.

The arguments \texttt{colnum} and \texttt{colinc} correspond to the two numeric parameters to \texttt{^T}. The \texttt{kind} argument specifies the style of tabbing. It must be one of \texttt{:line (^T)}, \texttt{:block (^T)}, \texttt{:line-relative (^T)}, or \texttt{:block-relative (^T)}.

\begin{itemize}
  \item \texttt{logical-block-pop \ args \&optional (stream \ *standard-output*) \Rightarrow \ item}
  
  The macro \texttt{logical-block-pop} is identical to \texttt{pop} except that, when used in conjunction with \texttt{within-logical-block}, it supports \texttt{*print-length*} and \texttt{*print-circle*}. It is an error to use \texttt{logical-block-pop} anywhere other than syntactically nested within a call on \texttt{within-logical-block}.

  The argument \texttt{args} must be a symbol or expression acceptable to \texttt{pop}. The \texttt{stream} argument (which defaults to \texttt{*standard-output*}) follows the standard conventions for stream arguments to printing functions. If \texttt{stream} is a special stream bound by \texttt{within-logical-block}, then \texttt{logical-block-pop} performs the special operations described below. Otherwise, \texttt{logical-block-pop} is identical to \texttt{pop}.

  Each time \texttt{logical-block-pop} is called, it performs three tests. First, it checks to see whether \texttt{args} is a cons. If not, \texttt{"."} is printed on \texttt{stream}, \texttt{args} is printed on \texttt{stream} using \texttt{write}, and the execution of the immediately containing \texttt{within-logical-block} is terminated except for the printing of the suffix. Second, if \texttt{*print-length*} is \texttt{nil} and \texttt{logical-block-pop} has already been called \texttt{*print-length*} times within the immediately containing logical block, \texttt{"..."} is printed on \texttt{stream} and the execution of the immediately containing \texttt{within-logical-block} is terminated except for the printing of the suffix. Third, if \texttt{*print-circle*} is not \texttt{nil}, \texttt{args} is checked to see if it is a circular reference. If it is, \texttt{"."} followed by an appropriate \#n\# marker is printed on \texttt{stream} and the execution of the immediately containing \texttt{within-logical-block} is terminated except for the printing of the suffix. If all three of the tests above fail, \texttt{logical-block-pop} pops the top value off of \texttt{args} and returns this value.

  \item \texttt{logical-block-count \&optional (stream \ *standard-output*)} \Rightarrow \texttt{nil}
  
  This macro is identical to \texttt{logical-block-pop} except that it does not take an \texttt{args} argument, always returns \texttt{nil}, and only performs the second test discussed above. It is useful when the components of a non-list are being printed.

As an example of using the functions above, consider that \texttt{tabular-style} is defined as follows. Using \texttt{logical-block-tab} in the definition makes it easy to communicate the parameter \texttt{tabsize} to the algorithm controlling the dynamic arrangement of output. By means of the \texttt{list} argument of \texttt{within-logical-block} and the macro \texttt{logical-block-pop}, the definition is robust in the face of malformed lists and supports \texttt{*print-length*}, \texttt{*print-level*}, and \texttt{*print-circle*}. 
Dynamic Control of the Arrangement of Output

(defun tabular-style (s list &optional (colon? T) atsign? (tabsize nil))
  (declare (ignore atsign?))
  (if (null tabsize) (setq tabsize 16))
  (within-logical-block (s list :prefix (if colon? "(" "")
          :suffix (if colon? ")" ""))
    (when list
      (loop (write (logical-block-pop list s) :stream s)
          (if (null list) (return nil))
          (write-char #\space s)
          (logical-block-tab :block-relative 0 tabsize s)
          (conditional-newline :fill WM)))))

The function below prints a vector using #(...) notation. A dummy list argument of nil for within-logical-block is used along with the macro logical-block-count, because the structure being printed is not a list. Here the functional interface to XP is appropriate, because format control strings do not provide any way to traverse a vector.

(defun print-vector (v *standard-output*)
  (within-logical-block (nil nil :prefix "#(" :suffix ")")
    (let ((end (length v)) (i 0))
      (when (plusp end)
        (loop (logical-block-count)
          (write (aref v i))
          (if (= (incf i) end) (return nil))
          (write-char #\space)
          (conditional-newline :fill)))))))

* fill-style stream list &optional (colon? T) atsign? ⇒ nil
* linear-style stream list &optional (colon? T) atsign? ⇒ nil
* tabular-style stream list &optional (colon? T) atsign? (tabsize 16) ⇒ nil

The directives */fill-style/, */linear-style/, and */tabular-style/ (see page 16) are supported by the three functions above. These functions can also be called directly by the user. Each function prints parentheses around the output if and only if colon? (default T) is not nil. Each function ignores its atsign? argument and returns nil. Each function handles abbreviation and circularity detection correctly, and uses write to print list when given a non-list argument.

The function linear-style prints a list either all on one line, or with each element on a separate line. The function fill-style prints a list with as many elements as possible on each line. The function tabular-style is the same as fill-style except that it prints the elements so that they line up in columns. This function takes an additional argument tabsize (default 16) that specifies the column spacing.
4. Pretty Printing Types of Objects

As discussed in Section 2, the pretty printing performed by XP is directed by the value of *print-dispatch*. The value of this variable is a print dispatch table. This table is initialized with a number of entries that specify how to pretty print all the built-in Common Lisp macros and special forms. You can tailor the pretty printer to your own needs by adding new entries into the table and/or replacing existing entries. Multiple styles of pretty printing can be supported by constructing several tables and switching between them. The primitives supported for operating on print dispatch tables are designed in analogy with the operations associated with read tables.

- **copy-print-dispatch** &optional (table *print-dispatch*) ⇒ copy
  A copy is made of table, which defaults to the current print dispatch table. If table is nil, a copy is made of the standard print dispatch table initially defined by XP.

- **define-print-dispatch** type-specifier options &body function ⇒ T
  This puts an entry into a print dispatch table. The type-specifier is implicitly quoted and is a standard Common Lisp type specifier as defined in [4]. It specifies what type of objects the entry is applicable to. The function specifies how to pretty print that type of object. When appropriate, the function will be called with two arguments: an output stream and the object to print. The options are a list of pairs of a keyword and a value. Three different keywords are possible:

  (:table table)
  This option specifies where to put the print dispatch entry being defined. If this option is not present, the entry is placed in the table stored in *print-dispatch*.

  (:priority number)
  This option specifies a priority that is used to control the order in which entries in the print dispatch table are compared against an object to be printed. If this option is not present, the priority defaults to 0.

  (:name name)
  If present, this option specifies a name to be given to function. This makes it possible to reuse the function—e.g., in another call on define-print-dispatch.

Before creating a new entry in the table, define-print-dispatch removes any existing entry with the same (equal) type specifier no matter what its priority. This guarantees that there will never be two entries that have the same (equal) type specifier. However, given a particular object it is likely that it will match several entries. The entry to use for printing is selected by taking the matching entry with the highest priority.

Before discussing the handling of the function and type-specifier in detail, it is useful to consider a brief example. The definition below specifies a new way to print ratios. Once entered into the print dispatch table, it alters the way every ratio is pretty printed.
(The use of \texttt{\&rest x} in the argument list below makes it possible to use \texttt{/ratio-print/} in a \texttt{format} string.)

\begin{verbatim}
(define-print-dispatch ratio ((:name ratio-print)) (stream obj \&rest x)
  (declare (ignore x))
  (format stream "#.(/ -,OF -,OF)" (numerator obj) (denominator obj)))

(pprint '(2/3 250 -4/5))
\end{verbatim}

Prints:

\begin{verbatim}
(#.(/ 2. 3.) 250 #.(/ -4. 5.))
\end{verbatim}

\section*{Pretty Printing Functions}

The function in a \texttt{define-print-dispatch} call can be specified in one of five ways. First, as shown in the example above, it can be an argument list followed by a body consisting of one or more statements. The argument list must be consistent with the fact that the function will be called with a stream and an object. The function can assume that the object satisfies the associated type specifier.

Second, the function can consist solely of an instance of \texttt{#'name}. If so, the indicated function will be used as the printing function. (Note that if \texttt{ratio-print} used \texttt{"W} instead of \texttt{"OF} to print the numerator or denominator, infinite recursion would occur, because these parts are themselves integers.)

\begin{verbatim}
(define-print-dispatch integer ((:priority 1)) #'ratio-print)

(pprint '(2/3 250 -4/5))
\end{verbatim}

Prints:

\begin{verbatim}
(#.(/ 2. 3.) #.(/ 250. 1.) #.(/ -4. 5.))
\end{verbatim}

Third, the function can be an instance of \texttt{"..."}.

\begin{verbatim}
(define-print-dispatch (and ratio (satisfies plusp)) ((:priority 2))
  #"(+ /ratio-print/)"

(pprint '(2/3 250 -4/5))
\end{verbatim}

Prints:

\begin{verbatim}
((+ #.(/ 2. 3.) #.(/ 250. 1.) #.(/ -4. 5.))
\end{verbatim}

Fourth, the function can be \texttt{nil}. In this case, any currently existing entry for the type specifier is removed without replacing it by anything. Pretty printing for objects that match the indicated type specifier will be controlled by the other entries they match (if any).

\begin{verbatim}
(define-print-dispatch (and ratio (satisfies plusp)) () nil)

(pprint '(2/3 250 -4/5))
\end{verbatim}

Prints:

\begin{verbatim}
(#.(/ 2 3) #.(/ 250 1) #.(/ -4 5))
\end{verbatim}

Fifth, the function can be totally omitted. In this case, any currently existing entry for the type specifier is removed and the standard pretty printing function (if any) corresponding to the type specifier is reentered into the table at the newly specified priority.
Pretty Printing Type Specifiers

When an object is to be pretty printed, the print dispatch table stored in the variable *print-dispatch* is consulted to find out how to print it. This is done by looking at the entries in the table in the order of their priorities and selecting the first entry for which (typep object type-specifier) is not nil. The type specifiers can take any of the forms described in the Common Lisp book [4]. In addition, the type specifier cons is extended to make it more useful.

It is expected that the table may contain entries whose type specifiers partially overlap in various ways. For example, the standard print dispatch table contains a catchall entry for printing lists in general and a number of entries for printing specific kinds of lists. As a result, you must be careful with your choice of priorities. If an object matches two different entries that have the same priority, there are no guarantees as to which entry will be used.

Pretty printing lisp code. The definition below shows the default method XP uses for printing lists that represent data rather than programs. (The functions linear-style, fill-style, and tabular-style are all defined with their colon? and atsign? arguments optional so that they can be used as define-print-dispatch functions.) It can be very useful in some situations to use tabular-style instead of fill-style to print data lists.

\[(\text{define-print-dispatch cons ((:priority -10)) #'fill-style})\]

However, it should be noted that, in Lisp there is no completely reliable way to distinguish between lists that represent program code and lists that merely represent data. Nevertheless, the following type specifier is useful for specifying tests that do a good job most of the time.

- cons &optional (car-type T) (cdr-type T)

When used simply as the symbol cons, this type specifier matches any cons cell. When used in the form above, it matches a cons cell only if the car of the cell matches the type specifier car-type and the cdr of the cell matches the type specifier cdr-type.

The examples below show some of the predefined pretty printing functions for Lisp code. By default, function calls are printed in the standard way—i.e, either all on one line or with the arguments one to a line indented after the function name. Lists beginning with cond are printed the same way as function calls except that the clauses are always printed in linear style, rather than in the format suggested by their cars. Lists beginning with setq are printed with two arguments on each line. Lists beginning with quote are printed using the standard "'" syntax. Note the care taken to ensure that data lists that happen to begin with quote will be printed legibly.

\[(\text{define-print-dispatch (cons (and symbol (satisfies fboundpM ((:priority -5)) W-- ---- :>I$))})\]
\[(\text{define-print-dispatch (cons (member (and symbol (satisfies fboundpM ((:priority -5)) W-- ---- :>I$)) cond))})\]
Pretty Printing Types of Objects

Pretty printing structures. An important use of XP is to print data structures. In fact, typical Lisp interactions call for much more printing of data than printing of programs. Pretty printing can do just as much to enhance the readability of this output as it can to enhance the readability of code. As shown below, pretty printing functions for structures that have been defined without the :type option can be specified with reference to their types.

\[
\text{(define-print-dispatch (cons (member setq)) () #":::":::W":::I"::I":::Q::W":::I"::W":::I"::":::":::})}
\]

\[
\text{(define-print-dispatch (cons (member quote)) () (s list)}
\]

\[(\text{if (and (consp (cdr list)) (null (cddr list))) (format s #":"W" s (cadr list)) (fill-style s list))})
\]

With *print-pretty* T, line width 23, and *print-lines* 3 prints:

(principle-family
 "Lucy" and
 ("Mark" "Bob" ---

With *print-pretty* T, *print-level* 3, and *print-length* 3 prints:

(while (primary-family #<"Lucy" and ("Mark" "Bob" # . . . >))

With *print-pretty* T, *print-escape* nil, and *print-circle* T prints:

(while (primary-family #<Lucy and #1=(Mark Bob #1# Bill . Dan))>

A key thing to notice about the pretty printing function above is that without the programmer having to take any explicit action, it tolerates a malformed kids list and correctly follows the printer control variables *print-lines*, *print-level*, *print-length*, *print-escape* and *print-circle*. This should be contrasted with Common Lisp's current support for structure print self functions, where it is difficult to handle *print-level* and *print-length* correctly and impossible to handle *print-circle* correctly.

There is clearly a close relationship between XP's pretty printing functions for structures and the standard concept of a print function for a structure. However, there is a fundamental difference in approach. XP stores the function in a print dispatch table rather than directly with the structure. This makes it possible to simultaneously support several different styles of printing by maintaining several different dispatch tables and to switch rapidly between them. However, it has the disadvantage that pretty printing functions are only used when *print-pretty* is not nil. This could have the effect of forcing you to define a pretty printing function and a print function for the same structure merely to ensure that the structure is always printed the same way. To avoid this
problem, XP uses the print function for a structure when no pretty printing function is available.

**Efficiency.** Given only what is said above, the process of determining the printing functions to be used for the various parts of an object to be printed would be horrendously inefficient, because every part of the object would have to be compared against every entry in the print dispatch table. XP avoids this problem by speeding up the selection process in two ways.

A hash table is used to very rapidly compare an object against every entry with a type specifier in the print dispatch table that has the form \((\text{cons (member symbol)})\). A second hash table is used to rapidly compare objects with type specifiers that are the types of structures defined without the :type option. It is advisable for you to make as many print dispatching entries as possible fit into these two categories.

**Predefined pretty printing functions.** To support traditional Lisp pretty printing style, XP provides pretty printing functions for all of the Common Lisp macros and standard forms. The user can change the way any given kind of list is printed by defining a new list pretty printing function for it. To facilitate the correct utilization of priorities, Figure 41 summarizes the contents of the standard print dispatch table initially defined by XP.

<table>
<thead>
<tr>
<th>Priority</th>
<th>Type Specifier</th>
<th>Pretty Printing Action</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>((\text{cons (member symbol)}))</td>
<td>(~60) printers for Lisp code.</td>
</tr>
<tr>
<td>-5</td>
<td>((\text{cons (and symbol (satisfies fboundp))}))</td>
<td>Print as function call.</td>
</tr>
<tr>
<td>-10</td>
<td>(\text{cons})</td>
<td>Print using fill-style.</td>
</tr>
</tbody>
</table>

Figure 4.1: Contents of the initial print dispatch table.

If an attempt is made to pretty print an object that does not match any entry in the current print dispatch table, one of the following default actions is taken. Arrays are printed appropriately following the value of \(*\text{print-array}*. Structures are printed using their print functions (if any). Otherwise the object is printed using the standard printer, with \(*\text{print-pretty}* bound to \(\text{nil}\).
5. Ugly details

XP is implemented in fully portable Common Lisp. However, a number of compromises had to be made for this to be true. The discussion above deliberately glosses over these problems on the theory that there is no fundamental need for them to exist and they would not exist if XP were implemented as part of Common Lisp, rather than as a separate package. This section discusses these problems in detail and explains how they have been dealt with in Symbolics Common Lisp [8]. It is hoped that they can be dealt with as easily in other implementations of Common Lisp.

Insufficient integration with non-pretty printing. XP never comes into play unless *print-pretty* is not nil, #"..." is encountered, or a format string is evaluated that contains one of xp's special format directives. This is done as a matter of safety and so that XP will operate purely as an add-on system. However, it has drawbacks. For example, the variable *print-right-margin* only has an effect when XP is in operation. Similarly, *last-abbreviated-printing* only gets set when XP is in operation. If XP were combined into a Common Lisp implementation, it would be natural to combine it directly into the standard output routines, and support variables like *print-right-margin* and *last-abbreviated-printing* all of the time.

Getting XP to take effect. By far the biggest problem is that Common Lisp has no standard mechanism for allowing a new pretty printer to be specified. The function xp::install uses shadowing to redefine the standard Common Lisp printing functions. However, this is of somewhat limited utility for several reasons.

First, shadowing fundamental functions like print and defstruct is a dangerous practice. In particular, while it can work when it is done by one subsystem, it is almost never going to work if two subsystems try to do it.

Second, shadowing only effects programs that are read into the package where XP is installed after XP has been installed. Among other things this means that it will not change the printing that is initiated by the Lisp system itself. For example, it will not change the printing done by the top level read-eval-print loop. You can change this easily enough, but that leaves a host of other places where the system initiates output such as various things printed by the debugger.

You could try to install XP more firmly by altering the function cells of the standard printing functions. However, this is an exceptionally dangerous thing to do and is quite likely to break the system. (To start with, it will break XP.)

In any event, clobbering these function cells would not fix the problem, because many Lisp implementations do output by calling primitive output functions that are not part of the standard set of Common Lisp output functions. As a result, clobbering the standard functions still would not fix all output.

A better answer is to have a hook in the Lisp system that is prepared to accept a new pretty printer. Symbolics Common Lisp has such a hook in the form of the variable scl:*print-pretty-printer*. The symbolics Common Lisp version of the function xp::install sets this variable to a value that causes XP to be used for all pretty printing.

An interesting aspect of the function installed on scl:*print-pretty-printer* is that it traps any errors that occur when printing is done. This is very useful when such errors
are happening while you are trying to debug something else. However, trapping such errors can be very annoying when it is a `define-print-dispatch` function or something like that that is being debugged. You can turn off the error trapping feature by setting the variable `xp:*allow-errors*` to T.

**Obtaining information from output streams.** To operate as intended, XP needs to be able to get two pieces of information from an output stream before starting to print into it. This information is obtained by calling the following two functions.

- `xp::output-width &optional (output-stream *standard-output*) ⇒ width`
  
  Returns the maximum number of characters that can be printed on a single line without causing truncation or wraparound when printing to `output-stream`, or `nil` if this cannot be computed.

- `xp::output-position &optional (output-stream *standard-output*) ⇒ position`
  
  Returns the number of characters printed so far on the current output line in `output-stream`, or `nil` if this cannot be computed.

Unfortunately, although every implementation of Common Lisp probably supports internal functions providing this information, there are currently no standard Common Lisp functions yielding this information. XP contains appropriate definitions of the functions above for several different implementations of Common Lisp; however, in other implementations it is reduced to using default (useless) definitions of these two functions that always return `nil`. If you are operating in one of these other implementations (you can tell by looking at the beginning of the XP file) you should provide better implementations for these functions.

**Imperfect integration with structures.** To operate as intended, XP needs to be able to determine which types are structure types. This is done by calling the following function.

- `xp::structure-type-p type ⇒ boolean`
  
  Returns non-nil if and only if `type` is a structure type defined by `defstruct` without the `:type` option, and `nil` otherwise.

Unfortunately, although every implementation of Common Lisp probably supports an internal function providing this information, there is currently no standard Common Lisp function yielding this information. XP contains appropriate definitions of the functions above for several different implementations of Common Lisp; however, in other implementations it is reduced to pessimistically assuming that the only structures are ones defined using `xp::defstruct` (which is used to shadow `lisp:defstruct` if `xp::install` is called with `:shadow T`). If you are operating in one of these other implementations (you can tell by looking at the beginning of the XP file) you should provide an implementation for `xp::structure-type-p`.

Another potential problem is that XP assumes that if a structure is defined using XP’s shadowed version of `defstruct`, then the structure’s print function (if any) is defined using the XP’s shadowed versions of the various printing functions. As a result, XP does
not hesitate to call such a print function with one of its special pretty printing streams as an argument. Since it is possible for the assumption to be false, this can lead to problems.

**Limitations on the definition of new type specifiers.** Due to the extreme restrictions Common Lisp places on the ways complex type specifiers can be constructed, there is no implementation independent way to support the extended definition of the type specifier `cons` as a first class type specifier, even though it does not violate the 'spirit' of what can and cannot be a type specifier. As a consequence of this limitation, the extend form of `cons` can only be used in conjunction with `define-print-dispatch`. This could easily be remedied if XP were incorporated directly into Common Lisp.

**Imperfect integration with format.** XP supports 99% of the functionality of `format`, but not all of it. In particular, XP takes pains to fully support `format` as described in [4]. However, there is one place where XP falls short of this goal.

As discussed above, the standard `format` directive `<...>` is more or less incompatible with `<...>` and the other pretty printing directives. However, it is permissible to have a garden variety instance of `<...>` nested in a `format` string that also contains some pretty printing directives. In this situation, XP uses the standard function `format` to process the part of the `format` string containing the `<...>`. Unfortunately, this only works when it can be determined exactly how many arguments will be used by the `<...>`. As a result, XP is forced to require that `<...>` cannot contain complex directives like `@{...}`, `~`, and `*` or anything similar. This problem could be straightforwardly fixed if XP duplicated all of the code in the standard function `format` that supports `<...>` instead of merely using the standard function `format`.

Another area of difficulty concerns the fact that XP is oriented around supporting formatting functions (e.g., created by `#...`) rather than `format` control strings. Nevertheless, in the interest of upward compatibility, XP allows `format` strings to be used. However, there are three complications with this.

First, to avoid having to implement an interpreter for `format` strings as well as a compiler for `#...`, XP converts each `format` string that contains any of XP's special directives into a function the first time it is encountered. This works well as long as `format` strings are not modified by side-effect. The caching of converted `format` control strings can be turned off by setting the variable `xp:*format-string-cache*` to `nil`.

Second, some implementations of Common Lisp support `format` directives beyond the ones defined in [4] or support additional features of the standard `format` directives. No attempt is made to support this functionality in conjunction with the special directives supported by XP. However, in order to make sure that merely installing XP will not break any currently running code, XP converts `format` strings to functions only if they contain one or more of XP's special directives. If a `format` string does not contain any of XP's special directives, it is left as a string and the standard function `format` is used to process it.

Third, the dual approach of using XP for some `format` strings and standard `format` for others has some implications with regard to the directive `~` and the usage `~{...}` with no body. If these forms exist in a `format` string that does not contain any of XP's special directives, then the control arguments they receive must be `format` strings rather than
functions. On the other hand, if they exist in a format string that contains any of XP's special directives then the control arguments they receive must either by functions or format strings that can be successfully converted to functions by XP.

Beyond the problems above, there are several points where the documentation in [4] is not entirely clear, and about which different implementations of Common Lisp seem to disagree. XP may not be doing the right thing in these situations. In particular:

How exactly does "@" act in a "{...}" and "@{...}"? Is it relative to the arguments being processed by the whole loop, or relative to the arguments being process by the current step of the loop? XP assumes the former.

There is no detailed grammar given for how a directive can be specified. In particular, can a colon or atsign modifier be specified before all the parameters have been specified? XP assumes not.

What is supposed to happen to the argument list when a cycle of "{...}" is prematurely terminated by a "@" directive? In particular, are the arguments that have been processed supposed to have been removed or not? In the interest of simplicity, XP assumes they should be removed.

Is "@" supposed to operate identically when accessed via "{" with no body and a "@?" directive? XP assumes that it is. (It would be quite difficult for XP to support things any other way.)

Assumptions about the read table. It was possible to more than double the speed of XP by assuming that the characters "a-z", "A-Z", "+", "*", "<", ">", and "-" always have the same syntax as defined in the initial read table. This assumption would not be necessary if Common Lisp provided any quick way to determine what the syntax of a character is.

The delay caused by buffering. As part of its operation, XP buffers up output characters before actually printing them into the appropriate stream. The fundamental source of the efficiency of the pretty printing algorithm is that things are designed so that the buffer never has to contain more than one line width worth of output. The algorithm sends output to the underlying stream one line width at a time. The buffer is not guaranteed to be completely empty until the printing is completed. Thus there is typically a delay between the time characters are put in the buffer (e.g., by a call on some printing function in the pretty printing function for some type of object) and the time they appear in the output stream. This can be confusing if a process which is performing pretty printing is interrupted (e.g., during the debugging of a pretty printing function).

The functions finish-output and force-output can be used to force the internal buffer to empty out. However, to maintain internal consistency in the pretty printing algorithm, all of the logical blocks that have been started but not yet completed are printed as if they will not be able to fit on a single line. As a result, the output may not look the same as it would if the buffer were not prematurely forced to empty out.

Taking full advantage of information about formatting special forms. Symbolics Common Lisp contains a large number of special forms that have to be pretty printed in special ways in addition to the standard Common Lisp special forms. In Symbolics Common Lisp, XP takes advantage of the fact that the ZWIE editor maintains information about these forms in order to determine how to pretty print them. A
similar sharing of information between XP and the Lisp editor might be useful in other environments as well.

Reprinting an abbreviated object. XP supports a special function that facilitates the reprinting of the last abbreviated object. In Symbolics Common Lisp, `xp::install` sets up the the key sequence `<function> <resume>` so that it triggers the reprinting in full of the last abbreviated printing. This turns out to be very convenient. A similar mechanism might be useful in other environments as well.

No support for font variations. The pretty printing algorithm depends on extensive calculations about how much space strings of characters will take up when displayed. These calculations are greatly simplified by assuming that every character will have the same fixed width when displayed. Only newlines are treated specially.

It should be noted that (except for `T`) the standard `format` directives all make the same simplifying assumption. However, this assumption can lead to problems in some situations. For example, it is inadvisable to use literal tab characters when pretty printing and the output produced by XP looks quite strange when it is displayed using a variable width font.

The above notwithstanding, the fundamental algorithms used by XP could be extended to handle characters of variable width and characters whose width depends on the position where they are displayed. In addition, the interface has been designed as much as possible to be independent of this issue.

The only user-visible things that refer to actual lengths are the variables `*print-right-margin*`, `*default-right-margin*`, `*print-miser-width*`, and the numeric arguments to `T`, `I`, and `/tabular-style/` and their functional counterparts. All of these measurements must be in the same units, but it does not matter a great deal what the units are. A good choice would be something like the length of an "m" in the current font. This will work out right for fixed width fonts and pretty well for variable width ones. Programmers should be advised to avoid explicit lengths—i.e., they should rely on streams knowing how wide they are and use `*0:I` whenever possible to indicate indentations.

Doing Things Right

XP is the kind of program that cannot really be supported in a totally portable way in Common Lisp. This is true both due to the various problems outlined above and because there are a number of things where portability has only been achieved at the significant sacrifice of efficiency. The right thing to do when incorporating XP into a Common Lisp is not to merely load the system and use it, but rather to totally integrate it with the way printing is done.

Places where XP needs to be more tightly integrated with the primitive printing facilities. There are a number of places where XP falls back on using the standard printing facilities. The standard function `write` is used to print objects for which there is no special printing function in `*print-dispatch*`. The standard function `format` is used to support complex `format` directives like `R`, `C`, and `F`. In both cases this is done by having the standard functions output into a string and then copying the string into XP's internal buffer.
This is effective, but quite slow. As a result, printing with XP is noticeably slower than printing with \texttt{*print-pretty* nil}. This is unfortunate, because as demonstrated by PP [7], the algorithms underlying XP are sufficiently efficient that it is possible for pretty printing to be virtually as fast as non-pretty printing. The only thing that is missing is proper integration with the printing subprimitives.

To a certain extent, superior integration could have been achieved by duplicating more of the basic printing code as part of XP. However, it would not be possible to achieve perfect integration in a portable way, because Common Lisp does not provide any way to get information out of the read table. As a result, \texttt{write} must be used to print symbols. (As discussed above, XP gets around this problem to some extent by making a few conservative assumptions about the read table.)

The right thing to do when incorporating XP into a Common Lisp is to modify XP so that it directly calls the appropriate printing sub-primitives and modify the sub-primitives so that they put their output directly into XP's internal buffer.

**Places where the primitive printing facilities need to be more tightly integrated with XP.** To get XP to really take over for all pretty printing, it needs to be installed in such a way that it is always used. The right way to do this is to insert a call to it deep in the standard printing code at the point where the variable \texttt{*print-pretty*} is tested.

In addition, all of the functions that make use of \texttt{format} strings (e.g., \texttt{error}) should be extended so that they can make use of \texttt{"\ldots"} and the special pretty printing functions.

Beyond this, there is a more subtle problem. Internally, XP operates in two stages. The first stage supports dispatching through \texttt{*print-dispatch*} along with various kinds of abbreviation. This dispatching is accessed via the directive \texttt{\textbackslash W} and the function \texttt{write}. The second stage performs the actual dynamic formatting decisions.

The second stage essentially operates as a special kind of output stream. This stream receives output characters and commands related to logical blocks and conditional new-lines. After deciding where to insert line breaks, the output is sent on to the ordinary stream that is the eventual destination of the output. This organization is largely hidden. However, it becomes apparent in one key situation.

When writing special printing functions (i.e., to be used with \texttt{define-print-dispatch}, \texttt{defstruct}, or \texttt{\textbackslash /\ldots\textbackslash /}) it is permissible to use any kind of printing function. However, these functions are called with special XP streams as arguments rather than ordinary streams. (This is essential, because XP must be able to catch all output before it gets to the real output stream.) As a result, all of the standard printing functions (e.g., \texttt{print}, \texttt{terpri}, \texttt{force-output}) have to be modified so that they will operate correctly when passed a special XP stream.

Alternatively, the fundamental concept of what an output stream is can be altered so that every stream is capable of supporting the operations of the second stage of XP. This approach was taken by PP, and worked very well. A particular advantage of this is that it allows proper integration of XP with functions like \texttt{with-open-file} and \texttt{with-output-to-string}. Unfortunately, it is impossible to create a new kind of stream in a portable manner, because Common Lisp does not provide any appropriate primitives.

**Using XP to the full.** Because the capabilities of XP go way beyond typical pretty
printers, XP can be used in many ways that typical pretty printers cannot. As a result. It is useful to extend a Common Lisp so that it takes better advantage of pretty printing. To start with, since (when properly integrated) XP is just as fast as a non-pretty printer their is no reason not to have the default value of *print-pretty* be T.

Beyond this, many kinds of output done by the system itself should be upgraded to take advantage of XP. As an example, in Symbolics Common Lisp, the trace facilities can be used to print out information about the arguments a function is called with whenever it is called. This output is produced using standard format control strings and always prints all the arguments on one line. If the arguments are large, this output ends up being more or less unreadable. The Symbolics version of xp:install replaces the trace printer with a new function that takes full advantage of xp. There are dozens of other places where such changes could profitably be made.
Bibliography

(Revised version MIT/AIM-611a, September 1982.)
Historical Note

The original #print system was written in MacLisp in 1977. The primary motivation behind #print was producing a pretty printer significantly faster than the Goldstein pretty printer [1] then in use. By means of the same basic algorithms that are still in use in XP, #print succeeded in being almost as fast as ordinary printing. In addition to using fundamentally the same algorithms as XP, #print followed the same basic approach of having pretty printing control strings for specifying how to control the dynamic layout of output and mechanisms for associating pretty printing functions with types of objects. However, the interface was markedly different in two respects.

First, the pretty printing control strings used by #print were developed before format came into wide use. Although fundamentally similar to format control strings, they looked very different, because they treated unmarked characters as directives instead of literal characters to be printed out. Literal output had to be specified by enclosing it in apostrophes. The pretty printing control strings were also described in a confusing way that exposed unnecessarily much of the underlying algorithm. Second, the mechanisms used by #print to associate pretty printing functions with types of objects were significantly more cumbersome and less powerful than those supported by XP.

#print was released for general use in the MIT AI Laboratory in January of 1978. However, probably because satisfactory documentation was never produced, #print was not extensively used by anyone other than the author.

In early 1980, #print was cleaned up and re-released under the name Gprint. The primary change was that, Gprint extended the power of (and further complicated) the mechanisms for associating pretty printing functions with types of objects. In late 1981, full documentation was prepared [5] and Gprint began to reach a wide audience.

In the spring and summer of 1982, Gprint was converted to run on the Symbolics Lisp Machine. After a delay of a year or so, Gprint was adopted as the standard pretty printer on the Symbolics Lisp Machine, in which role it is still being used today. However, Symbolics decided not to publicize the interface that can be used to define new ways of pretty printing things. In the summer of 1983, DEC converted Gprint into their Common Lisp and adopted it and its interface as official parts of their Common Lisp.

In 1984, Gprint was totally rewritten in Symbolics Lisp Machine Lisp and re-emerged as pp [7]. The key advance was that PP unified the concepts of format control strings and pretty printing control strings, recasting everything in format's syntax. From the point of view of people who understand format, this simplified things tremendously. PP also somewhat simplified the mechanisms for associating pretty printing functions with types of objects by eliminating the least used features. In the fall of 1985, DEC upgraded their Common Lisp to include PP and its interface instead of Gprint.

In 1988, PP was totally rewritten in completely portable Common Lisp and re-emerged as XP. XP's major contribution is that, by taking an entirely different approach, it greatly simplifies the mechanisms for associating pretty printing functions with types of objects and makes them even more powerful than the mechanisms supported by Gprint. Since September 1988, XP has been in experimental use as the pretty printer in CMU Common Lisp. Currently, XP and its interface are being considered by the Common Lisp standardization committee for adoption as a formal part of Common Lisp.
Functional Summary and Index

The entries below describe the various functions, variables, and macros supported by XP, showing their inputs and outputs, the pages where documentation can be found, and one line descriptions. The next page summarizes the extensions to format.

conditional-newline *optional (stream *standard-output*) ⇒ nil
  p. 20  Functional interface to "_."

copy-print-dispatch *optional (table *print-dispatch*) ⇒ copy
  p. 23  Copies a print dispatch table.

*default-right-margin* default value 70.
  p. 4  Default line width to use when pretty printing.

define-print-dispatch type-specifier options &body function ⇒ T
  p. 23  Defines a new print dispatch table entry.

cons *optional (car-type T) (cdr-type T)
  p. 25  Type specifier that matches a cons if its car and cdr are of the indicated types.

fill-style stream list *optional (colon? T) atsign? ⇒ nil
  p. 22  Function underlying "/fill-style/.

xp::install &key (:package *package*) (:macro T) (:shadow T) (:remove nil) ⇒ T
  p. 1  Makes XP ready for use.

*last-abbreviated-printing*
  p. 19  Variable recording last printing event that was abbreviated.

linear-style stream list *optional (colon? T) atsign? ⇒ nil
  p. 22  Function underlying "/linear-style/.

logical-block-count *optional (stream *standard-output*) ⇒ nil
  p. 21  Supports length abbreviation.

logical-block-indent relative-to n *optional (stream *standard-output*) ⇒ nil
  p. 20  Functional interface to "I.

logical-block-pop args *optional (stream *standard-output*) ⇒ item
  p. 21  Supports length and circularity abbreviation.

logical-block-tab kind colnum colinc *optional (stream *standard-output*) ⇒ nil
  p. 20  Functional interface to "T.

*print-lines* default value nil
  p. 3  Variable limiting the total number of lines pretty printed.

*print-miser-width* default value 40.
  p. 3  Variable specifying when pretty printing should switch to space saving mode.

*print-right-margin* default value nil
  p. 4  Variable specifying the line width to use when pretty printing.

*print-dispatch* default value causes standard pretty printing
  p. 3  Variable containing the current print dispatch table controlling pretty printing.

*print-dispatch* default value causes standard pretty printing
  p. 3  Variable containing the current print dispatch table controlling pretty printing.

tabular-style stream list *optional (colon? T) atsign? (tabsize 16) ⇒ nil
  p. 22  Function underlying "/tabular-style/.

within-logical-block (stream-symbol list &key :prefix :per-line-prefix :suffix)
  &body body ⇒ nil
  p. 19  Functional interface to "<...>".
Functional format control string.

The directive °W (write object p. 4) uses the function write to output the corresponding format argument without forcing the setting of any output control variables.

°W (p. 4) Prints an argument following all output control variables.
°:W (p. 4) Forces pretty printing.
°QW (p. 4) Suppresses length and depth abbreviation.

There are three special directives for printing lists. Each of them prints parentheses around the output when used with the colon modifier.

°/fill-style/ (p. 16) Prints as many elements as possible on each line.
°/linear-style/ (p. 16) Prints elements all on one line or one to a line.
°c/tabular-style/ (p. 17) Prints elements in a table with column spacing c.

The directive °<prefix>; body°; suffix°> (logical block, p. 7) iterates over a list argument using body to print the elements of the list in a logical block. The prefix and suffix are printed before and after the block respectively.

°<...°: (p. 7) Denotes a logical block and descends into a list argument.
°@<...°: (p. 7) Operates on all the remaining arguments.
°: <body°: (p. 9) Prefix and suffix default to "(" and ")" respectively.
°<body°:°> (p. 9) Body printed to fill the line width.
°<prefix°;...» (p. 9) Prefix printed on each line.

The indentation in a logical block is initially set to the column position of the first character in the block. The directive °I (set indentation, p. 13) is used to alter the indentation within a logical block. If omitted, the parameter defaults to zero. When a logical block is printed in miser style, all instances of °I are ignored.

°nI (p. 13) Indentation set to position of first character in block plus n.
°n: I (p. 13) Indentation set to position of directive plus n.

The directive °_ (conditional newline, p. 10) specifies a place where a newline can be inserted in a logical block. For a discussion of line breaks inserted by other means than °_, see page 12.

°_ (p. 10) Linear-style conditional newline.
°: _ (p. 11) Fill-style conditional newline.
°0 _ (p. 10) Miser-style conditional newline.
°: 0 _ (p. 12) Mandatory-style conditional newline.

The directive °T has augmented capabilities.

°: T (p. 14) Tab relative to containing section.