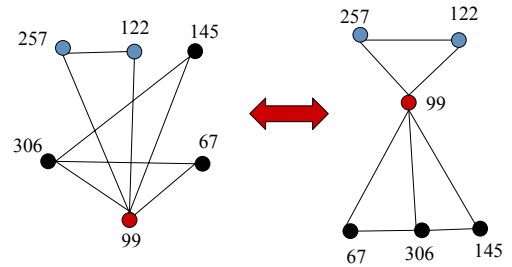




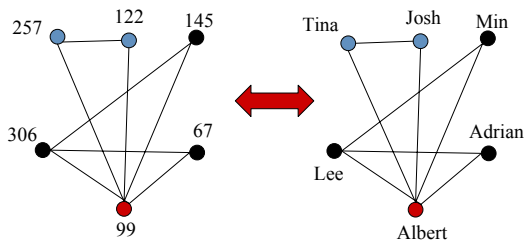
# More Graphs



## Topology, not Geometry



## Equivalent (Isomorphic) Graphs



## Graph Isomorphism

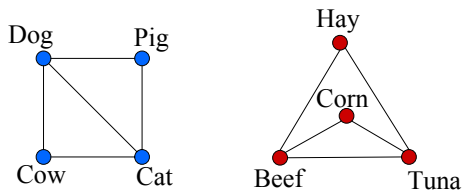
Graphs  $G_1$  and  $G_2$  are **isomorphic** if there exists a **bijection**  $f: V_1 \rightarrow V_2$  such that for all  $u, v$  in  $V_1$

- the **edge**  $(u, v)$  is in  $G_1$
- iff the **edge**  $(f(u), f(v))$  is in  $G_2$

- If there is a one-to-one correspondence between the nodes of  $G_1$  and  $G_2$  that preserves all edge connections.

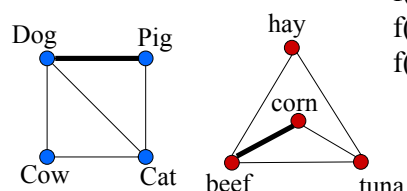


## Are these Isomorphic?



## Find a Mapping

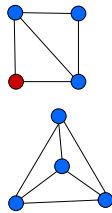
**Function**  
 $f(\text{Dog}) = \text{beef}$   
 $f(\text{Cat}) = \text{tuna}$   
 $f(\text{Cow}) = \text{hay}$   
 $f(\text{Pig}) = \text{corn}$



6	9	13	7
12	10	5	
3	1	4	11
15	8	14	2

## Finding the Mapping

- Not easy, can try all possible mappings
  - Roughly  $n!$  possibilities
- Can test for Invariants
  - Same number of nodes, edges
  - Same degree distributions
  - Preserves cycles, longest path, etc



Copyright © Radhika Nagpal, 2002.

L4-2.7

6	9	13	7
12	10	5	
3	1	4	11
15	8	14	2

## In-class Problem 1

Copyright © Radhika Nagpal, 2002.

L4-2.8

6	9	13	7
12	10	5	
3	1	4	11
15	8	14	2

## Problem with False Proof 1

Proof (silently) assumes any 2-ended  $G_{n+1}$  can be built from a 2-ended  $G_n$ . This isn't true!

Consider the counter example, it is two ended but I cannot construct it by adding an edge to another two ended graph.



Copyright © Radhika Nagpal, 2002.

L4-2.9

6	9	13	7
12	10	5	
3	1	4	11
15	8	14	2

## Problem with False Proof 2

After removing a vertex from  $G_{n+1}$ , the claim that  $G_n$  still has 2 vertices of degree 1 and rest degree 2. This is not true!

The same counter example works:



Copyright © Radhika Nagpal, 2002.

L4-2.10

6	9	13	7
12	10	5	
3	1	4	11
15	8	14	2

## Revisit: Coloring with $d_{\max}$ colors

- Induction Hypothesis
  - $P(n)$  = a graph with  $n$  vertices and maximum degree  $d_{\max}$  can be colored with  $d_{\max} + 1$  colors
- Inductive Step
  - Do you justify why your proof doesn't have the same pitfalls?

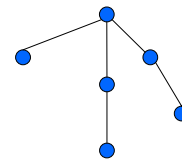
Copyright © Radhika Nagpal, 2002.

L4-2.11

6	9	13	7
12	10	5	
3	1	4	11
15	8	14	2

## Trees

- *Definition:* A tree is simple connected graph with no cycles.



Copyright © Radhika Nagpal, 2002.

L4-2.12

6	9	13	7
12	10	5	
3	1	4	14
15	8	11	2

## Applications of Trees

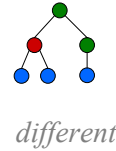
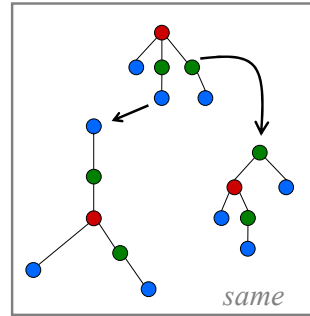
- Data structures for sorting, searching
- Spanning Trees
- Game Trees (alpha-beta trees)
- Prefix codes (Huffman encoding)
- Many algorithms based on trees (6.046)

Copyright © Radhika Nagpal, 2002.

L4-2.13

6	9	13	7
12	10	5	
3	1	4	14
15	8	11	2

## Tree Isomorphisms



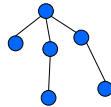
Copyright © Radhika Nagpal, 2002.

L4-2.14

6	9	13	7
12	10	5	
3	1	4	14
15	8	11	2

## Trees

- *Definition:* A tree is a simple connected graph with no cycles.



- **Exercise:** Draw a tree with 5 vertices
- **Question:** How many edges does your tree have? 3, 4 or 5?

Copyright © Radhika Nagpal, 2002.

L4-2.15

6	9	13	7
12	10	5	
3	1	4	14
15	8	11	2

## Another Tree Definition

- No matter how you draw it, you get 4 edges.
- *Definition 2:* A tree is a connected graph with  $n$  vertices and  $n - 1$  edges.
- In fact, *a tree is the smallest connected graph on  $n$  vertices!*

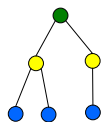
Copyright © Radhika Nagpal, 2002.

L4-2.16

6	9	13	7
12	10	5	
3	1	4	14
15	8	11	2

## Equivalent Definitions of Trees

- A connected graph with no cycles
- A connected graph where  $|E| = |V| - 1$
- A connected graph where removing any edge leaves a disconnected graph
- A graph such that there exists a unique simple path between any two vertices



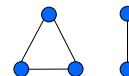
Copyright © Radhika Nagpal, 2002.

L4-2.17

6	9	13	7
12	10	5	
3	1	4	14
15	8	11	2

## Be careful with these definitions

- What is wrong with this definition?
  - A tree is a graph with  $n$  vertices and  $n - 1$  edges.
- Counter-example



Copyright © Radhika Nagpal, 2002.

L4-2.18



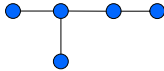
## Different Trees with 5 vertices



Vertex degrees



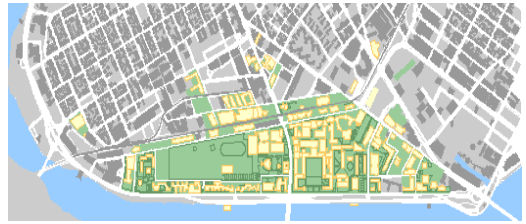
1 2 2 2 1  
4 1 1 1 1  
1 3 1 2 1



Sum is always 8  
( $2 \times$  edges)



## MIT Building Connections



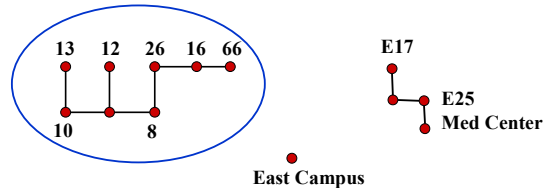
## Connectivity and Paths

- Can you get from building 10 to 36 without crossing more than 5 other buildings
  - Is there a path of length  $k$  from  $u$  to  $v$ ?
- How many different ways are there to get from building 10 to building 36?
  - How many different paths are there from  $u$  to  $v$ ?



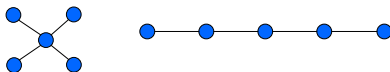
## Connected Components

- Can we get from building 10 to building E17?
  - Is there a path between  $u$  and  $v$ ?
  - Are  $u$  and  $v$  connected?



## Smallest Connected Graph

- MIT administration wants the number of physical connections between buildings to be minimum but still have everything connected.
  - What is the smallest connected graph I can construct? **ANY TREE**



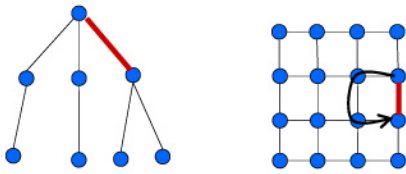
## Cut Edge

- Definition:** An edge is a **cut edge** if removing it from the graph disconnects two connected components

*Problem with our smallest connected graph*  
 – any edge disruption disconnects the graph!

6	9	13	7
12	10	8	
3	1	14	5
15	4	11	2

## Fault Tolerant Design



Robustness versus Cost  
(Networks, Highways, etc)

Copyright © Radhika Nagpal, 2002.

L4-2.25

6	9	13	7
12	10	8	
3	1	14	5
15	4	11	2

## Newark Airport Shutdown

January 9, 1995

... Among the hundreds receiving citations so far are a contractor and two sub-contractors who were working on a parking garage at Newark airport when **three 26,000-volt power cables** were cut. This shut down the airport and affected travel throughout the world. Incoming domestic and international flights were waved off to other airports for nearly 24 hours, and outbound planes couldn't get off the ground, leaving gaping holes in the schedules of virtually every major airline.

Copyright © Radhika Nagpal, 2002.

L4-2.26

6	9	13	7
12	10	8	
3	1	14	5
15	4	11	2

## In-class Problems

Copyright © Radhika Nagpal, 2002.

L4-2.27

6	9	13	7
12	10	8	
3	1	14	5
15	4	11	2

## Extra slides

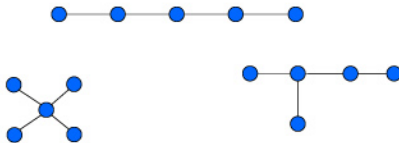
Copyright © Radhika Nagpal, 2002.

L4-2.28

6	9	13	7
12	10	8	
3	1	14	5
15	4	11	2

## In-class Problem 0

*Prove that all trees with five vertices are isomorphic to these three:*



Copyright © Radhika Nagpal, 2002.

L4-2.29

6	9	13	7
12	10	8	
3	1	14	5
15	4	11	2

## Solution to Problem 0

Use the following three facts:

- sum of degrees =  $2 \cdot \text{edges}$
- tree has  $n - 1$  edges, 5 vertex tree has 4 edges
  - therefore sum of degrees is 8
- tree is connected
  - therefore each vertex must have degree at least 1.
  - 3 degrees left over to distribute
- The only possibilities are:
  - 4 1 1 1 1
  - 3 2 1 1 1
  - 2 2 2 1 1

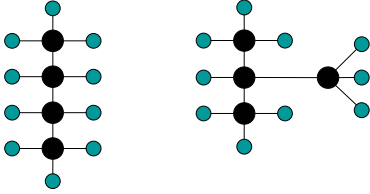
Copyright © Radhika Nagpal, 2002.

L4-2.30

6	9	13	7
12	10	5	
3	4	8	11
15	2	14	1

## Graphs with Same Degrees

- Example: isomers of butane
  - Same degree distribution, but not isomorphic



Copyright © Radhika Nagpal, 2002.

1.4-2.31

6	9	13	7
12	10	5	
3	4	8	11
15	2	14	1

## Why are some problems easier than others?

- **2 colorable** (no odd cycles)
- **3 colorable** (NP- complete)
- **Euler circuits** (connected and all even degree)
- **Hamiltonian circuits** (NP complete)

Copyright © Radhika Nagpal, 2002.

1.4-2.32