Robustness of Interdependent Random Geometric Networks

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Abstract—We propose an interdependent random geometric graph (RGG) model for interdependent networks. Based on this model, we study the robustness of two interdependent spatially embedded networks where interdependence exists between geographically nearby nodes in the two networks. We study the emergence of the giant mutual component in two interdependent RGGs as node densities increase, and define the percolation threshold as a pair of node densities above which the mutual giant component first appears. In contrast to the case for a single RGG, where the percolation threshold is a unique scalar for a given connection distance, for two interdependent RGGs, multiple pairs of percolation thresholds may exist, given that a smaller node density in one RGG may increase the minimum node density in the other RGG in order for a giant mutual component to exist. We derive analytical upper bounds on the percolation thresholds of two interdependent RGGs by discretization, and obtain 99% confidence intervals for the percolation thresholds by simulation. Based on these results, we derive conditions for the interdependent RGGs to be robust under random failures and geographical attacks.

I. INTRODUCTION

Cyber-physical systems such as smart power grids and smart transportation networks are being deployed towards the design of smart cities. The integration of communication networks and physical networks facilitates network operation and control. In these integrated networks, one network depends on another for information, power, or other supplies in order to properly operate, leading to interdependence. For example, in smart grids, communication networks rely on the electric power from power grids, and simultaneously control power generators [1]. Failures in one network may cascade to another network, which potentially makes the interdependent networks vulnerable to failures and attacks.

Cascading failures in interdependent networks have been extensively studied in the statistical physics literature since the seminal work in [2], where each of the two interdependent networks is modeled as a random graph. Nodes in the two random graphs are interdependent, and a node is functioning if both itself and its interdependent node are in the giant components of the respective random graphs. After initial node failures in the first graph, their interdependent nodes in the second graph fail. Thus, a connected component in the second graph may become disconnected, and the failures of the disconnected nodes cascade back to (their interdependent) nodes in the first graph. As a result of the cascading failures, removing a small fraction of nodes in the first random graph destroys the giant components of both graphs.

To model spatially embedded networks, an interdependent lattice model was studied in [3]. Under this model, geographical attacks may cause significantly more severe cascading failures than random attacks. Removing nodes in a finite region (i.e., a zero fraction of nodes) may destroy the infinite clusters in both lattices [4].

If every node in one network is interdependent with multiple nodes in the other network, and a node is content to have at least one supply node, failures are less likely to cascade [5]. Although the one-to-multiple interdependence exists in real-world spatially embedded interdependent networks (e.g., a control center can be supported by the electric power generated by more than one power generator), it has not been previously studied using spatial graph models.

We use a random geometric graph (RGG) to model each of the two interdependent networks. The two RGGs are allowed to have different connection distances and densities, which can represent two networks that have different average link lengths and scales. These differences between the two networks were not captured in the lattice model studied in the previous literature. Moreover, the interdependent RGG model is able to capture the one-to-multiple interdependence in spatially embedded networks, and provides a more general framework for studying interdependent networks.

Robustness is a key design objective for interdependent networks. We study the conditions under which a positive fraction of nodes are functioning in interdependent RGGs as the number of nodes approaches infinity. In this case, the interdependent RGGs percolate. Consistent with previous research [2], [3], [5], the robustness of interdependent RGGs under random failures and geographical attacks is measured by whether percolation exists after the failures and attacks. To the best of our knowledge, our paper is the first to study the percolation of interdependent spatial network models using a mathematically rigorous approach.

The main contributions of this paper are as follows.

1) We propose an interdependent RGG model for two interdependent networks, which captures the differences in the scales of the two networks as well as the one-to-multiple interdependence in spatially embedded networks.
2) We derive the first analytical upper bounds on the percolation thresholds of the interdependent RGGs, above which a positive fraction of nodes are functioning.

3) We obtain 99% confidence intervals for the percolation thresholds, by mapping the percolation of interdependent RGGs to the percolation of a square lattice where the probability that a bond in the lattice is open is evaluated by simulation.

4) We characterize sufficient conditions for the interdependent RGGs to percolate under random failures and geographical attacks. In particular, if the node densities are above any upper bound on the percolation threshold obtained in this paper, the interdependent RGGs remain percolated after a geographical attack.

The rest of the paper is organized as follows. We state the model and preliminaries in Section II. We derive analytical upper bounds on percolation thresholds in Section III and obtain confidence intervals for percolation thresholds in Section IV. In Section V, we study the robustness of interdependent RGGs under random failures and geographical attacks. Section VI concludes the paper.

II. MODEL

A. Preliminaries on RGG and percolation

An RGG in a two-dimensional square consists of nodes generated by a Poisson point process and links connecting nodes within a given connection distance. Let \( G(\lambda, d, a^2) \) denote an RGG with node density \( \lambda \) and connection distance \( d \) in an \( a \times a \) square. The studies on RGG focus on the regime where the expected number of nodes \( n = \lambda a^2 \) is large.

We first present some preliminaries useful for developing our model. The giant component of an RGG is a connected component that contains \( \Theta(n) \) nodes. A node belongs to the giant component with a positive probability \( \Theta(n)/n \) if the giant component exists. For a given connection distance, the percolation threshold is a node density above which a node belongs to the giant component with a positive probability (i.e., a giant component exists) and below which the probability is zero (i.e., no giant component exists). By scaling, if the percolation threshold is \( \lambda^* \) under connection distance \( d \), then the percolation threshold is \( \lambda^* c^2 \) under connection distance \( d/c \). Therefore, without loss of generality, in this paper, we study the percolation thresholds represented by node densities, for given connection distances.

The RGG is closely related to the Poisson boolean model, where nodes are generated by a Poisson point process on an infinite plane. Let \( G(\lambda, d) \) denote a Poisson boolean model with node density \( \lambda \) and connection distance \( d \). The difference between \( G(\lambda, d) \) and \( G(\lambda, d, a^2) \) is the number of nodes in \( G(\lambda, d) \) is infinite while the expected number of nodes in \( G(\lambda, d, a^2) \) is large but finite. The Poisson boolean model can be viewed as a limit of the RGG as the number of nodes approaches infinity. The percolation threshold of \( G(\lambda, d) \) under a given \( d \) is defined as the density above which a node belongs to the infinite cluster with positive probability and below which the probability is zero. It has been shown that a node belongs to the infinite cluster with positive probability if and only if an infinite cluster exists, and thus the percolation of \( G(\lambda, d) \) can be equivalently defined as the existence of the infinite cluster. Moreover, the percolation threshold of \( G(\lambda, d) \) is identical with the percolation threshold of \( G(\lambda, d, a^2) \).

B. Interdependent RGGs

Two interdependent networks are modeled by two RGGs \( G_1(\lambda_1, d_1, a_1^2) \) and \( G_2(\lambda_2, d_2, a_2^2) \) on the same \( a \times a \) square. A node in one graph is interdependent with all the nodes in the other graph within the interdependent distance \( d_{dep} \). See Fig. 1 for an illustration. Nodes in one graph are supply nodes for nodes in the other graph within \( d_{dep} \). The physical interpretation of supply can be either electric power or information that is essential for proper operation.

![Fig. 1. Two interdependent RGGs with interdependent distance \( d_{dep} \).](image)

We define mutual component and giant mutual component in interdependent RGGs, in the same way as one defines the connected component and giant component in a single RGG.

**Definition 1.** Let \( V_i^0 \) denote nodes in a connected component in \( G_i(\lambda_i, d_i, a_i^2), \forall i \in \{1, 2\} \). If each node in \( V_i \subseteq V_i^0 \) has at least one supply node in \( V_{\overline{i}} \subseteq V_{\overline{i}}^0 \) within \( d_{dep} \), \( \forall i \neq j \), then nodes \( V_1 \) and \( V_2 \) form a mutual component of the interdependent RGGs.

If, in addition, \( V_i \) contains \( \Theta(n_i) \) nodes, where \( n_i = \lambda_i a_i^2 \), \( \forall i \in \{1, 2\} \), then \( V_1 \) and \( V_2 \) form a giant mutual component.

A mutual component can be viewed as an autonomous system in the sense that nodes in a mutual component have supply nodes in the same mutual component, and in each graph, nodes that belong to a mutual component are connected regardless of the existence of nodes outside the mutual component. Note that a node can receive supply from any of its supply nodes in the same mutual component. Nodes in a giant mutual component are functioning, since they are connected to a large number of nodes in the network. This definition of functioning is consistent with previous research on interdependent networks based on random graph models.

For a fixed \( d_{dep} \), if a giant mutual component exists in interdependent RGGs \( G_1(\lambda_1, d_1, a_1^2) \) and \( G_2(\lambda_2, d_2, a_2^2) \), then a giant mutual component exists in interdependent RGGs \( G_1'(\lambda_1', d_1, a_1^2) \) and \( G_2(\lambda_2, d_2, a_2^2) \), where \( \lambda_1' > \lambda_1 \). This can be explained by coupling \( G_1' \) with \( G_1 \) as follows. By randomly removing each node in \( G_1' \) independently with probability
there exist percolation thresholds $\lambda_1^*$ and $\lambda_2^*$, the remaining nodes in $G'_1$ has density $\lambda_1$, and a giant mutual component exists in the interdependent RGGs that consist of $G_2$ and the RGG formed by the remaining nodes in $G'_1$. Since adding nodes to a graph does not disconnect any mutual component, a giant mutual component exists in the interdependent RGGs $G'_1$ and $G_2$. By the same analysis, a giant mutual component also exists in interdependent RGGs $G_1(\lambda_1, d_1, a^2)$ and $G_2'(\lambda_2', d_2, a^2)$ for a fixed $d_{dep}$, if $\lambda_2' > \lambda_2$.

We define a percolation threshold of two interdependent RGGs as follows.

**Definition 2.** A pair of node densities $(\lambda_1^*, \lambda_2^*)$ is a percolation threshold of two interdependent RGGs, given connection distances $d_1$, $d_2$ and the interdependent distance $d_{dep}$, if a giant mutual component exists in $G_1(\lambda_1, d_1, a^2)$ and $G_2(\lambda_2, d_2, a^2)$ for $\lambda_1 > \lambda_1^*$ and $\lambda_2 > \lambda_2^*$, and no giant mutual component exists otherwise.

For fixed $d_1$, $d_2$ and $d_{dep}$, there may exist multiple percolation thresholds. Intuitively, the larger the node density is in one graph, the smaller the required node density is in the other graph in order for the giant mutual component to exist. This is in contrast with the situation for a single RGG where there is a unique percolation threshold for a fixed $d$.

There is a non-trivial phase transition in the interdependent RGGs. If $\lambda_i$ is smaller than the percolation threshold of a single RGG $G_i(\lambda_i, d_i, a^2)$, clearly there does not exist a giant mutual component in the interdependent RGGs. Therefore, $\lambda_i^* > 0$, $\forall i \in \{1, 2\}$. As we will see in the next section, there exist percolation thresholds $\lambda_i^* < \infty$, $\forall i \in \{1, 2\}$, which concludes the non-trivial phase transition.

Given that the conditions for the percolation of $G_i(\lambda_i, d_i, a^2)$ and $G_i(\lambda_i, d_i)$ are the same, in most parts of the paper we study the percolation of two interdependent Poisson boolean models on the same infinite plane, $G_{IntDep} = (G_1(\lambda_1, d_1), G_2(\lambda_2, d_2), d_{dep})$, by applying techniques in continuum percolation. The percolation of $G_{IntDep}$ is defined as the existence of an infinite mutual cluster, which consists of an infinite number of connected nodes in both $G_1(\lambda_1, d_1)$ and $G_2(\lambda_2, d_2)$ where every node has at least one supply node in the same infinite mutual cluster. In the rest of the paper we sometimes use $G_i$ to denote both $G_i(\lambda_i, d_i, a^2)$ and $G_i(\lambda_i, d_i)$. The model that it refers to will be clear from the context.

### C. Related work

The model which is closest to ours is the interdependent lattice model, first proposed in [10] and further studied in [3], [4]. In this model, nodes in a network are represented by the open sites (nodes) of a square lattice, where every site is open independently with probability $p$. Network links are represented by the bonds (edges) between adjacent open sites. Every node in one lattice is interdependent with one randomly chosen node within distance $r_d$ in the other lattice. The percolation threshold of the interdependent lattice model is characterized as a function of $r_d$, assuming the same $p$ in both lattices [10]. Percolation of the model where some nodes do not need to have supply nodes was studied in [3]. The analysis relies on quantities estimated by simulation and extrapolation, such as the fraction of nodes in the infinite cluster of a lattice for any fixed $p$, which cannot be computed rigorously. In contrast, we study the percolation of the interdependent RGG model using a mathematically rigorous approach.

### III. Analytical upper bounds on percolation thresholds

In interdependent RGGs, nodes in the giant mutual component are viewed as functioning while all the other nodes are not. Thus, a node is functioning only if it is in the giant component of its own graph, and it depends on at least one node in the giant component of the other graph. For any node $b_1$ in $G_1$, although the number of nodes in $G_2$ within distance $d_{dep}$ from $b_1$ follows a Poisson distribution, the number of functioning nodes is hard to calculate, since the fraction of nodes in the giant component of $G_2$ is unknown. Moreover, the nodes in the giant component of $G_2$ are clustered, and thus the thinning of the nodes in $G_1$ due to a lack of supply nodes in $G_2$ is difficult to characterize. To overcome these difficulties, we consider the percolation of two RGGs jointly, instead of studying the percolation of one RGG with reduced node density due to a lack of functioning supply nodes.

We now give an overview of our approach. We develop mapping techniques (discretizations) to characterize the percolation of $G_{IntDep}$ by the percolation of a discrete model. Mappings from a model whose percolation threshold is unknown to a model with known percolation threshold are commonly employed in the study of continuum percolation. For example, one can study the percolation threshold of the Poisson boolean model $G(\lambda, d)$ by mapping it to a triangle lattice and relating the state of a site in the triangle lattice to the point process of $G(\lambda, d)$. By the mapping, the percolation of the triangle lattice implies the percolation of $G(\lambda, d)$. Consequently, an upper bound on the percolation threshold of $G(\lambda, d)$ is given by $\lambda$ for which the triangle lattice percolates, a known quantity [11], [8]. In general, more than one mapping can be applied, and the key is to search for a mapping that gives a good (smaller) upper bound. Following this idea, we propose different mappings that fit different conditions to obtain upper bounds on the percolation thresholds of $G_{IntDep}$.

#### A. Ratio $d_2/d_1$ is small

Given $G_{IntDep} = (G_1(\lambda_1, d_1), G_2(\lambda_2, d_2), d_{dep})$, without loss of generality we assume that $d_1 \leq d_2$. Moreover, we assume that $d_{dep} \geq \max(d_1/2, d_2/2)$ (see the remark at the end of the section for comments on this assumption). Let $c = \lfloor d_2/d_1 \rfloor = \max\{c : d_2/d_1 \geq c, c \in \mathbb{N}\}$. For small $c$, we study the percolation of $G_{IntDep}$ by mapping it to an independent bond percolation of a square lattice, and prove the following result.

**Theorem 1.** If $(\lambda_1, \lambda_2)$ satisfies

\[
(1 - e^{-\lambda_1 d_1^2/8})^c (1 - e^{-\lambda_2 e^2 d_2^2/8}) > 1/2, 
\]

then...
then \( G_{\text{IntDep}} = (G_1(\lambda_1, d_1), G_2(\lambda_2, d_2), d_{\text{dep}}) \) \( \text{percolates} \), where \( c = [d_2/d_1], \ d_1 \leq d_2, \text{ and } d_{\text{dep}} \geq d_2/2. \)

**Proof.** We first construct a square lattice as follows. Partition the plane into small squares of side length \( s = d_1/2\sqrt{2} \). A large square consists of \( c \times c \) small squares and has side length \( cs \). The **diagonals** of the large squares form the bonds of a square lattice \( L \), illustrated by the thick line segments in Fig. 2.

The state of a bond in \( L \) is determined by the point process \( \lambda \). The event that a bond is open depends on the point process \( \lambda \). The probability that the second condition is satisfied is \( p_2 = 1 - e^{-\lambda_2 cs^2 d_1^2/8} \).

To obtain a closed-form formula, instead of computing the exact probability, we compute a lower bound on the probability that the first condition is satisfied. The probability is lower bounded by the probability that the \( c \) small squares that intersect the bond each contain at least one node from \( G_1 \), given by

\[
p_1 = 1 - e^{-\lambda_2 cs^2 d_1^2/8}c.
\]

The probability that the second condition is satisfied is

\[
p_2 = 1 - e^{-\lambda_2 cs^2 d_1^2/8}.
\]

Given that the two Poisson point processes in \( G_1 \) and \( G_2 \) are independent, the probability that a bond is open is \( p_1p_2 \).

It remains to prove that the percolation of \( L \) implies the percolation of \( G_{\text{IntDep}} \). Consider two adjacent open bonds \((v_1, v_2), (v_2, v_3)\) in \( L \). Let \( S_1 \) and \( S_2 \) denote the two adjacent large squares of side length \( cs \) that contain the two open bonds.

Let \( S_1' \) and \( S_2' \) denote two adjacent small squares of side length \( s \) that contains \( v_2 \), within \( S_1 \) and \( S_2 \), respectively. See Fig. 2 for an illustration. Since \((v_1, v_2), (v_2, v_3)\) are open, under the second condition, nodes of \( G_2 \) exist in \( S_1 \) and \( S_2 \) and they are connected, because they are within distance \( 2\sqrt{2}cs \leq d_2 \). Under the first condition, nodes of \( G_1 \) form a connected path from the small square (within \( S_1 \), marked as 7 in Fig. 2) containing \( v_1 \) to \( S_1' \), and another path from the small square (within \( S_2 \)) containing \( v_3 \) to \( S_2' \). Moreover, the two paths are joined, because any pair of nodes in \( S_1' \) and \( S_2' \) are within distance \( 2\sqrt{2}s = d_1 \). Given that any pair of nodes within a large square have distance at most \( 2\sqrt{2}cs \leq d_2/2 \leq d_{\text{dep}} \), all the nodes have at least one supply node inside the large square that contains an open bond. To conclude, if the open bonds in \( L \) form an infinite cluster, then the nodes in \( G_{\text{IntDep}} \) form an infinite mutual cluster.

The event that a bond is open depends on the point processes in the large square that contains the bond, and is independent of whether any other bonds are open. As long as the probability that a bond is open, \( p_1p_2 \), is larger than \( 1/2 \), which is the threshold for independent bond percolation in a standard square lattice \([12]\), \( G_{\text{IntDep}} \) percolates.

The bound can be made tighter for any given \( c = [d_2/d_1] \), by computing more precisely the probability that the first condition is satisfied. We provide an example to illustrate the computation of an improved upper bound.

**Example:** Consider an example where \( d_1 = 1, d_2 = 2d_{\text{dep}} = 3 \). The probability that there is at least one node from \( G_2 \) in the large square of side length \( 3/2\sqrt{2} \) is \( p_2 = 1 - e^{-\lambda_2 \sqrt{2}/8} \).

The probability that a small square contains at least one node from \( G_1 \) is \( p_1 = 1 - e^{-\lambda_1 \sqrt{2}/8} \). The probability that the first condition is satisfied is

\[
p_1 = p_1^2 + (1 - p_1)p_1^4 + (1 - p_1)p_1^4 - (1 - p_1)p_1^6, \quad (1)
\]

obtained by considering all the sequences of adjacent small squares. Since \( p_1 \) computed by Eq. (1) is larger than \( p_2 \), the bound on \( \lambda_2 \) is smaller for any fixed \( \lambda_1 \).

**B. Ratio \( d_2/d_1 \) is large**

In the mapping from \( G_{\text{IntDep}} \) to the square lattice \( L \), the condition for a bond to be open becomes overly restrictive as \( d_2/d_1 \) increases. A path joining the two large squares that contain two adjacent bonds do not have to cross the small squares that contain the common end of the two bonds. We obtain another upper bound on the percolation threshold of \( G_{\text{IntDep}} \), given by the following theorem. This upper bound is tighter than the bound in Theorem 1 for larger values of \( d_2/d_1 \).

**Theorem 2.** If \((\lambda_1, \lambda_2)\) satisfies

\[
\left[ 1 - \frac{4}{3} (m + 1)e^{m \log 3(1 - p)} \right] \left[ 1 - \frac{4}{3} (2m + 1)e^{m \log 3(1 - p)} \right] p^m > 0.8639,
\]

then \( G_{\text{IntDep}} = (G_1(\lambda_1, d_1), G_2(\lambda_2, d_2), d_{\text{dep}}) \) percolates, where

\[
p = 1 - e^{-\lambda_1 d_1^2/8}, \quad p' = 1 - e^{-2D \lambda_2}, \quad D = \min\{d_2/\sqrt{10}, d_{\text{dep}}/\sqrt{5}\}, \ m = \lfloor 2D/d_1 \rfloor, \ d_1 \leq d_2, \text{ and } d_{\text{dep}} \geq d_2/2.
\]
This upper bound is obtained by mapping $G_{\text{IntDep}}$ to a dependent bond percolation model $L_D$. The mapping from the Poisson boolean model $G(\lambda, d)$ to $L_D$ was first proposed in [13] to study the percolation threshold of $G(\lambda, d)$, and later applied to the study of the robustness of random geometric networks [14]. We briefly describe the method in the previous literature that uses $L_D$ to study the percolation of $G(\lambda, d)$, and then prove Theorem 3 based on a similar method.

1) 1-dependent bond percolation model $L_D$: In the standard bond percolation model on a square lattice $L$, the event that a bond is open is independent of the event that any other bond is open. If in a square lattice $L_D$, the event that a bond is open may depend on the event that its adjacent bond is open, but is independent of the event that any non-adjacent bond is open, then $L_D$ is a 1-dependent bond percolation model on a square lattice. With the additional restriction that each bond is open with an identical probability, an upper bound on the percolation threshold of $L_D$ is 0.8639 [13].

The 1-dependent bond percolation model $L_D$ can be used to study the percolation of $G'$ where the points are generated by homogeneous Poisson point processes. To construct a mapping $G$ to $L_D$, consider two adjacent $D \times D$ squares $S_1$ and $S_2$ and let $R$ be the rectangle formed by the two squares. A bond $(v_1, v_2)$ that connects the centers of $S_1$ and $S_2$ is associated with $R$. Figure 3 illustrates the square lattice formed by the bonds, represented by thick line segments.

**Lemma 3.** Let the state of bond $(v_1, v_2)$ be determined by the homogeneous Poisson point processes of $G'$ inside $R$, and the conditions for a bond to be open be identical for all bonds. Then the bonds form a 1-dependent bond percolation model $L_D$ with identical bond open probabilities.

The proof of this lemma can be found in the technical report [15]. By properly setting the conditions for a bond to be open, the percolation of $L_D$ can imply the percolation of $G'$. We first look at an example in [12] that studies the percolation of $\Gamma(G, \lambda, d)$, and extend the technique to study $G_{\text{IntDep}}$.

**Example 12:** Let a bond be open if a path in $G(\lambda, d)$ crosses horizontally and another path in $G(\lambda, d)$ crosses vertically, where $R'$ is a $(2D - 2d) \times (D - 2d)$ rectangle that has the same center as $R$, and $S_i$ is a $(D - 2d) \times (D - 2d)$ square that has the same center as $S_1$. The reason for considering $R'$ and $S_i$ is that the existence of the two crossing paths over $R'$ and $S_i$ is entirely determined by the point process within $R$, while the existence of links within distance $d$ from the boundaries (and thus the crossings over $R$) may depend on nodes outside $R$.

If two adjacent bonds are open, the paths in $G(\lambda, d)$ in the two rectangles are joined. To see this, note that in Fig. 4 if the black and blue bonds (same direction) are both open, the crossings 1 and 2 intersect. If the black and red bonds (perpendicular) are both open, the crossings 1 and 3 intersect.

**Fig. 4.** Crossings over rectangles associated with two adjacent open bonds are joined.

If the square lattice $L_D$ percolates, open bonds form an infinite cluster. Paths in $G(\lambda, d)$ across the rectangles associated with the open bonds are connected and form an infinite cluster. Therefore, a node density above which $L_D$ percolates is an upper bound on the percolation threshold of $G(\lambda, d)$.

2) Proof of Theorem 3: We map $G_{\text{IntDep}}$ to $L_D$ by letting a bond in $L_D$ be open if the following three conditions are satisfied in its associated rectangle $R = S_1 \cup S_2$. The size of the rectangle satisfies $D = \min(d_2 / \sqrt{10}, d_{\text{dep}} / \sqrt{5}) \geq d_2 / 2\sqrt{5}$.

1) A path from $G_1$ crosses $R'$ horizontally, where $R'$ is a $(2D - 2d_1) \times (D - 2d_1)$ rectangle that has the same center as $R$.
2) A path from $G_1$ crosses $S_i'$ vertically, where $S_i'$ is a $(D - 2d_1) \times (D - 2d_1)$ square that has the same center as $S_1$.
3) There is at least one node from $G_2$ in $R$.

To bound the percolation thresholds of $G_{\text{IntDep}}$, in the technical report [15], we prove that the percolation of $L_D$ implies the percolation of $G_{\text{IntDep}}$ and compute the probability that the three conditions are satisfied using a method similar to [16].

**Remark:** We have assumed that $d_{\text{dep}} \geq \max(d_1 / 2, d_2 / 2) = d_2 / 2$ in this section. To see that this is a reasonable assumption, note that nodes in $G_1$ that have at least one functioning supply node are restricted in the region $R_{\text{dep}}$, where $R_{\text{dep}}$ is a union of disks with radius $d_{\text{dep}}$ centered at nodes in the giant component of $G_2$. If $R_{\text{dep}}$ is fragmented, it is not likely for disks of radius $d_1 / 2 < d_2 / 2$ centered at random locations within $R_{\text{dep}}$ to overlap, and it is not likely that a functioning giant component will exist in $G_1$, unless the node density in $G_1$ is large. Therefore, the interdependent distance $d_{\text{dep}}$ should be large enough so that $R_{\text{dep}}$ is a connected region, to avoid a large minimum node density in $G_1$. The region $R_{\text{dep}}$ can be made larger by increasing either $\lambda_2$ or $d_2$. Setting
then the bond is open. Otherwise, the bond is closed. The probability that a bond is open is larger than 0.8639, is an
component in \(S\). If the probability that a bond is open is larger than 0.8639, is an
complement of the connection process forms a connected
empty space that do not intersect nodes or links. If
vertical crossing of the complement process is defined analogously.

IV. CONFIDENCE INTERVALS FOR PERCOLATION THRESHOLDS

In the previous section, we discussed a method of mapping the percolation of \(G_{\text{IntDep}}\) to the percolation of the 1-dependent bond percolation model \(L_D\). The previous mapping and the mapping that we consider in this section both satisfy the following: 1) the percolation of \(L_D\) implies the percolation of \(G_{\text{IntDep}}\); 2) the event that determines the state of a bond depends only on the point process within its associated rectangle, thus preserving the 1-dependent property. The event probability can be computed or bounded analytically in the previous section. In contrast, in this section, we consider events whose probabilities are larger under the same point processes, but can only be evaluated by simulation. Since the events that we consider in this section are more likely to occur under the same point processes, the mappings yield tighter bounds.

Our mappings from \(G_{\text{IntDep}}\) to \(L_D\) extend the mappings from \(G(\lambda, d)\) to \(L_D\) proposed in [13]. For completeness, we first briefly summarize the mappings in [13] that determine upper and lower bounds on the percolation threshold of \(G(\lambda, d)\).

Upper bound for \(G(\lambda, d)\) [13]: Recalling Fig. 3 the event that a bond \((v_1, v_2) \in L_D\) is open is determined by the point process of \(G(\lambda, d)\) in the rectangle \(R = S_1 \cup S_2\), where \(S_1\) and \(S_2\) are squares. Let \(V_i\) denote the largest component formed by the points of \(G(\lambda, d)\) in \(S_i\). If \(V_i\) is the unique largest component in \(S_i\) (\(\forall i \in \{1, 2\}\)) and \(V_1\) and \(V_2\) are connected, then the bond is open. Otherwise, the bond is closed.

If \(L_D\) percolates, open bonds form an infinite cluster. As a result, the largest components in the squares that intersect the open bonds are connected in \(G(\lambda, d)\) and they form an infinite cluster. Therefore, a node density \(\lambda\), above which the probability that a bond is open is larger than 0.8639, is an upper bound on the percolation threshold of \(G(\lambda, d)\).

Lower bound for \(G(\lambda, d)\) [13]: Let the connection process of \(G(\lambda, d)\) be the union of nodes and links in \(G(\lambda, d)\). Let the complement of the connection process be the union of the empty space that do not intersect nodes or links. If the complement of the connection process forms a connected infinite region, then all the connected clusters in \(G(\lambda, d)\) have finite sizes and \(G(\lambda, d)\) does not percolate [13, 17]. Consider the complement of the connection process in rectangle \(R\). Let a bond (in \(L_D\)) associated with rectangle \(R\) be open if the complement process forms a horizontal crossing over the rectangle \(R'\) and a vertical crossing over the square \(S_1'\).

If \(L_D\) percolates, the complement process forms an infinite region and \(G(\lambda, d)\) does not percolate. To conclude, a node density, under which the probability that the complement process forms the two crossings is above 0.8639, is a lower bound on the percolation threshold of \(G(\lambda, d)\).

A. Upper bounds for \(G_{\text{IntDep}}\)

Two mutual components \(M = V_1 \cup V_2\) and \(\hat{M} = \hat{V}_1 \cup \hat{V}_2\) form one mutual component if and only if \(V_i\) and \(\hat{V}_i\) are connected in \(G(\forall i \in \{1, 2\})\). The necessity of the condition is obvious. To see that this condition is sufficient, note that every node in the connected component formed by \(V_i\) and \(\hat{V}_i\) has at least one supply node which belongs to the connected component formed by \(V_j\) and \(\hat{V}_j\) (\(\forall i, j \in \{1, 2\}, i \neq j\)). The condition can be generalized naturally for more than two mutual components to form one mutual component.

Algorithm 1 An algorithm that greedily computes a mutual component \(M_{\text{greedy}}(S)\) in region \(S\).

1) Identify the largest connected component \(V_i^0(S)\) in \(G_i(S)\), where \(G_i(S)\) consists of the nodes and links of \(G_i\) in \(S\). If there are multiple largest connected components, apply any deterministic tie-breaking rule (e.g., choose the component that contains a nodes with the smallest x-coordinate).
2) Remove nodes in \(V_i^0(S)\) that do not have supply nodes in \(V_i^0(S)\) (\(\forall i, j \in \{1, 2\}, i \neq j\)). Identify the largest connected component \(V_i^1(S)\) formed by the remaining nodes in \(V_i^0(S)\) (\(\forall i \in \{1, 2\}\)) and apply the same tie-breaking rule.
3) Repeat step 2 until \(V_i^{k+1}(S) = V_i^k(S)\left(\forall i \in \{1, 2\}\right)\). Let \(M_{\text{greedy}}(S) = V_i^k(S) \cup V_2(S)\).

Let a bond \((v_1, v_2)\) in \(L_D\) be open if the two components \(M_{\text{greedy}}(S_1)\) and \(M_{\text{greedy}}(S_2)\) form one mutual component, where \(M_{\text{greedy}}(S)\) is computed by Algorithm 1. See the technical report [15] for the rationale behind this algorithm. Since \(M_{\text{greedy}}(S_i)\) is unique in any square \(S_i\), a connected cluster in \(L_D\) implies that \(\{M_{\text{greedy}}(S_i)\}\) form one mutual component in \(G_{\text{IntDep}}\), where \(S_i\) are the squares that intersect the open bonds in the connected cluster. If the probability that a bond is open is larger than 0.8639, \(L_D\) percolates and \(G_{\text{IntDep}}\) also percolate.

An alternative condition for a bond to be open is that nodes in \(M_{\text{greedy}}(R)\) form a horizontal crossing over rectangle \(R'\) and a vertical crossing over square \(S_1'\) in both graphs (recall Fig. 4 and the condition for two mutual components to form one mutual component). In order for the existence of the two crossings to only depend on the point processes in \(R\), in the definition of \(R'\) and \(S_1'\), \(d = \max(d_1, d_2, d_{\text{dep}})\).

An upper bound on the percolation threshold can be obtained by either approach. The smaller bound obtained by the two approaches is a better upper bound on the percolation threshold of \(G_{\text{IntDep}}\).

B. Lower bounds for \(G_{\text{IntDep}}\)

In \(G_{\text{IntDep}}\), the connection process consists of nodes and links in mutual components. To avoid the heavy computation of mutual components, we study another model in which the
connection process $\tilde{P}_i$ of $G_i$ in the new model dominate the connection process $P_i$ of $G_i$ in $G_{\text{IntDep}}$. As a consequence, the complement of the connection process $\tilde{P}_i^c$ of $G_i$ in the new model is dominated by $P_i^c$ ($\forall i \in \{1, 2\}$). If $\tilde{P}_i^c$ percolates, then $P_i^c$ percolates and $P_i$ does not percolate. If either $P_1$ or $P_2$ does not percolate, then $G_{\text{IntDep}}$ do not percolate. Thus, node densities under which at least one of $\tilde{P}_i^c$ and $P_i^c$ percolates are lower bounds on the percolation thresholds of $G_{\text{IntDep}}$.

The new model can be viewed to have a relaxed supply requirement. In this model, every node (as opposed to nodes in the same mutual component) is viewed as a valid supply node for nodes in the other graph. A node $b_i$ in $G_i$ is removed if and only if there is no node in $G_j$ within the interdependent distance $d_{\text{dep}}$ from $b_i$ ($\forall i, j \in \{1, 2\}, i \neq j$). After all such nodes are removed, the remaining nodes in $G_i$ are connected if their distances are within the connection distance $d_i$. The computation of the connection process $\tilde{P}_i$ is efficient and avoids the computation of mutual components in $G_{\text{IntDep}}$ through multiple iterations.

The connection process $\tilde{P}_i$ in the new model dominates $P_i$ in the original model $G_{\text{IntDep}}$. On the one hand, for any realization, all the links in $P_i$ are present in $\tilde{P}_i$, because all the nodes in a mutual component have supply nodes, and links between these nodes are present in the new model as well. On the other hand, in the new model, nodes in a connected component $V_i$ in $G_j$ may depend on nodes in several connected components in $G_j$. In contrast, in $G_{\text{IntDep}}$, $V_i$ may be divided into several mutual components and links do not exist between two disjoint mutual components.

An algorithm that computes a lower bound on the percolation threshold of $G_{\text{IntDep}}$ is as follows. First obtain the connection process $\tilde{P}_i$ in the new model. Next in the $2D \times D$ rectangle $R$, consider the complement of the connection process $P_i^c$. Let $p_i$ denote the probability that there is a horizontal crossing over $R'$ and a vertical crossing over $S_i^c$ in the process $P_i^c$, where $R'$ and $S_i^c$ are the same as before. A lower bound on the percolation threshold of $G_{\text{IntDep}}$ is given by node densities under which $\max(p_1, p_2) \geq 0.8639$.

**C. Confidence intervals**

The probability that a bond is open can be represented by an integral that depends on the point processes in the rectangle $R$. However, direct calculation of the integral is intractable, so instead the integral is evaluated by simulation. In every trial of the simulation, nodes in $G_1$ and $G_2$ are randomly generated by the Poisson point processes with densities $\lambda_1$ and $\lambda_2$, respectively. The events that a bond is open are independent in different trials. Let the probability that a bond is open be $p$ given $(\lambda_1, \lambda_2)$. The probability that a bond is closed in $k$ out of $N$ trials follows a binomial distribution. The interval $[0.8639, 1]$ is a 99.5% confidence interval for $p$, given that $N = 100$ and $k = 5$. If $k < 5$, $p \in [0.8639, 1]$ with a higher confidence. This suggests that if $k \leq 5$, with 99.5% confidence $p \geq 0.8639$ and the 1-dependent bond percolation model $L_D$ percolates given $(\lambda_1, \lambda_2)$.

Based on this method, with 99.5% confidence an upper bound on the percolation threshold of $G_{\text{IntDep}}$ can be obtained by declaring a bond to be open using the method in Section [V-A] and with 99.5% confidence a lower bound can be obtained by declaring a bond to be open using the method in Section [V-B]. For a fixed $\lambda_2$, a 99% confidence interval for $\lambda_1$ is obtained, given by the interval between the upper and lower bounds. Confidence intervals for different percolation thresholds can be obtained by changing the value of $\lambda_2$ and repeating the computation. We make a similar remark as in [13]. The confidence intervals are rigorous, and the only uncertainty is caused by the stochastic point process in the $2D \times D$ rectangle. This is in contrast with the confidence intervals obtained by estimating whether $G_{\text{IntDep}}$ percolate based on extrapolating the observations of simulations in a finite region (which is usually not very large because of limited computational power).

**D. Numerical results**

The simulation-based confidence intervals are much tighter than the analytical bounds. Given that $d_1 = d_2 = 2d_{\text{dep}} = 1$, and $\lambda_2 = 2$, the upper and lower bounds on $\lambda_1^*$ are 2.25 and 1.80, respectively, both with 99.5% confidence. In contrast, even if $\lambda_2 \to \infty$, the analytical upper bound on $\lambda_1^*$ is no less than 3.372, which is the best available analytical upper bound for a single $G_1$ [11]. Confidence intervals for the percolation thresholds are plotted in Fig. 5.

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\( ^1 \) A connection process dominates another if the nodes and links in the first process form a superset of the nodes and links in the second process, for any realization of $G_i$. 

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![Fig. 5. The intervals between bars are 99% confidence intervals for percolation thresholds. The confidence intervals of two different $G_{\text{IntDep}}$ are plotted.](image)
which graph is able to resist more random node removals, by comparing the gap between the node density $\lambda_i$ and the percolation threshold $\lambda_i^*$ given $\lambda_j (i, j \in \{1, 2\}, i \neq j)$. The graph that can resist a smaller fraction of node removals is the bottleneck for the robustness of the interdependent RGGs. Besides, we can compute the maximum fraction of nodes that can be randomly removed from two graphs while guaranteeing the interdependent RGGs to be percolated.

We next show that the interdependent RGGs still percolate after a geographical attack that removes nodes in a finite connected region, if the node densities of the two graphs before the attack are above any upper bound on the percolation thresholds obtained in this paper (either analytical or simulation-based).

Recall that we obtained upper bounds on the percolation thresholds of $G_{\text{IntDep}}$ by mapping the percolation of $G_{\text{IntDep}}$ to the bond percolation of either a standard square lattice $L$ or the 1-dependent square lattice $L_D$. Moreover, whether a bond $e$ is open is entirely determined by the point processes in a finite region $R_e$ that contains the bond. After removing nodes of $G_{\text{IntDep}}$ in a connected finite geographical region, the state of a bond $e$ may change from open to closed only if $R_e$ intersects the attack region. Let $R_f$ be the union of $R_e$ that intersects the attack region. The region $R_f$ is also a connected finite region. As long as $L$ or $L_D$ still percolates after setting bonds in $R_f$ to be closed, $G_{\text{IntDep}}$ percolate.

Results from the percolation theory indeed indicate that setting all the bonds in a finite region $R_f$ to be closed does not affect the percolation of $L$ or $L_D$. For any percolated $L$, the probability that there exists a horizontal crossing of open bonds over a $kl \times l$ rectangle approaches 1 for any integer $k > 1$, as $l \to \infty$ (Lemma 8 on Page 64 of [12]). The percolation of $L$ (after setting all bonds in $R_f$ to be closed) is justified by the fact that the connected open bonds across rectangles form a square annulus that does not intersect $R_f$ (shown in Fig. 6), which is a standard approach to prove the percolation of $L$. Moreover, the percolation of $L_D$ after all bonds in $R_f$ can be proved in the same approach, by noting that the probability that open bonds of $L_D$ form a horizontal crossing over a rectangle approaches 1 as the rectangle size increases to infinity [13].

If the $kl \times l$ rectangle is large but finite, the probability that a horizontal crossing formed by open bonds exists is close to 1 if $L$ or $L_D$ percolates. Therefore, the same analysis demonstrates the robustness of finite interdependent RGGs under geographical attacks that remove a positive fraction of nodes in a connected region. Simulations for the robustness of interdependent RGGs under geographical attacks can be found in the technical report [15].

VI. CONCLUSION

We developed an interdependent RGG model for interdependent spatially embedded networks. We obtained analytical upper bounds and confidence intervals on the percolation thresholds. The percolation thresholds of two interdependent RGGs form a curve, which shows the tradeoff between the two node densities in order for the interdependent RGGs to percolate. The curve can be used to study the robustness of interdependent RGGs to random failures. Moreover, if the node densities are above any upper bound on the percolation thresholds obtained in this paper, then the interdependent RGGs remain percolated after a geographical attack.

REFERENCES