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Abstract—In this paper, we present a concise and coherent analysis of the constrained $\ell_1$ minimization method for stable recovering of high-dimensional sparse signals both in the noiseless case and noisy case. The analysis is surprisingly simple and elementary, while leads to strong results. In particular, it is shown that the sparse recovery problem can be solved via $\ell_1$ minimization under weaker conditions than what is known in the literature. A key technical tool is an elementary inequality, called Shifting Inequality, which, for a given nonnegative decreasing sequence, bounds the $\ell_2$ norm of a subsequence in terms of the $\ell_1$ norm of another subsequence by shifting the elements to the upper end.

Index Terms—$\ell_1$ minimization, restricted isometry property, shifting inequality, sparse recovery.

I. INTRODUCTION

RECONSTRUCTING a high-dimensional sparse signal based on a small number of measurements, possibly corrupted by noise, is a fundamental problem in signal processing. This and other related problems in compressed sensing have attracted much recent interest in a number of fields including applied mathematics, electrical engineering, and statistics. In signal processing setting, this new sampling theory has many applications. For example, interesting applications of compressed sensing in magnetic resonance imaging are described in [14]. Compressed sensing is also closely connected to coding theory, see, e.g., [1], and [8].

In compressed sensing, one considers the following model:

$$y = \Phi \beta + z$$

where the matrix $\Phi \in \mathbb{R}^{n \times p}$ (with $n \ll p$) and $z \in \mathbb{R}^n$ is a vector of measurement errors. The goal is to reconstruct the unknown signal $\beta \in \mathbb{R}^p$. In this paper, our main interest is the case where the signal $\beta$ is sparse and the noise $z$ is Gaussian, $z \sim \mathcal{N}(0, \sigma^2 I_n)$. We shall approach the problem by considering first the noiseless case and then the bounded noise case, both of significant interest in their own right. The results for the Gaussian case will then follow easily.

It is now well understood that the method of $\ell_1$ minimization provides an effective way for reconstructing a sparse signal in many settings. The $\ell_1$ minimization method in this context is

$$\min_{\gamma} \|\gamma\|_1 \quad \text{subject to } y - \Phi \gamma \in \mathcal{B}$$

where $\mathcal{B}$ is a bounded set determined by the noise structure. For example, $\mathcal{B} = \{0\}$ in the noiseless case and $\mathcal{B}$ is the feasible set of the noise in the case of bounded error.

The sparse recovery problem has now been well studied in the framework of the Restricted Isometry Property (RIP) introduced by [8]. A vector $v = (v_i) \in \mathbb{R}^p$ is $k$-sparse if $|\text{supp}(v)| \leq k$, where $\text{supp}(v) = \{i : v_i \neq 0\}$ is the support of $v$. For an $n \times p$ matrix $\Phi$ and an integer $k$, $1 \leq k \leq p$, the $k$-restricted isometry constant $\delta_k(\Phi)$ is the smallest constant such that

$$\sqrt{1 - \delta_k(\Phi)}\|c\|_2 \leq \|\Phi c\|_2 \leq \sqrt{1 + \delta_k(\Phi)}\|c\|_2$$

for every $k$-sparse vector $c$. If $k + k' \leq p$, the $k, k'$-restricted orthogonality constant $\theta_{k, k'}(\Phi)$, is the smallest number that satisfies

$$\quad |\langle \Phi c, \Phi c' \rangle| \leq \theta_{k, k'}(\Phi)\|c\|_2\|c'\|_2$$

for all $c$ and $c'$ such that $c$ and $c'$ are $k$-sparse and $k'$-sparse, respectively, and have disjoint supports. For notational simplicity, we shall write $\delta_k$ for $\delta_k(\Phi)$ and $\theta_{k, k'}$ for $\theta_{k, k'}(\Phi)$ hereafter.

It has been shown that $\ell_1$ minimization can recover a sparse signal with a small or zero error under various conditions on $\delta_k$ and $\theta_{k, k'}$. See, for example, [8], [9], and [7]. These conditions essentially require that every set of columns of $\Phi$ with certain cardinality approximately behaves like an orthonormal system. For example, the condition $\delta_{m} + \theta_{k, k + \theta_{k, 2k}} < 1$ was used in [8], $\delta_{3k} + 3\delta_{k} < 2$ in [7], and $\delta_{2k} + \theta_{k, 2k} < 1$ in [9]. Simple conditions involving only $\delta$ have also been used in the literature on sparse recovery, for example, $\delta_{2k} < \sqrt{2} - 1$ was used in [6]. In a recent paper, [5] sharpened the previous results by showing that stable recovery can be achieved under the condition $\delta_{1.75k} + \theta_{1.5k, 1.5k} < 1$ (or a stronger but simpler condition $\delta_{1.75k} < \sqrt{2} - 1$).

The present paper, we provide a concise and coherent analysis of the constrained $\ell_1$ minimization method for stable recovery of sparse signals. The analysis, which yields strong results, is surprisingly simple and elementary. At the heart of our simplified analysis of the $\ell_1$ minimization method for stable recovery is an elementary, yet highly useful, inequality. This inequality, called Shifting Inequality, shows that, given a finite decreasing sequence of nonnegative numbers, the $\ell_2$ norm of a subsequence can be bounded in terms of the $\ell_1$ norm of another

$^1$For a positive real number $\alpha$, $\delta_{\alpha k}$, and $\theta_{\alpha k, \alpha k}$ are understood as $\delta_{[\alpha k]}$ and $\theta_{[\alpha k], [\alpha k]}$.
subsequence by “shifting” the terms involved in the $\ell_2$ norm to the upper end.

The main contribution of the present paper is twofold: first, it is shown that the sparse recovery problem can be solved under weaker conditions and second, the analysis of the $\ell_1$ minimization method can be very elementary and much simplified. In particular, we show that stable recovery of $k$-sparse signals can be achieved if

$$\delta_{1,25k} + \theta_{k,1,25k} < 1.$$  

This condition is weaker than the ones known in the literature. In fact, our general treatment of this problem produces a family of sparse recovery conditions. Interesting conditions include

$$\delta_{1,025k} < \sqrt{2} - 1 \quad \text{and} \quad \delta_{3k} < 4 - 2\sqrt{3} \approx 0.535.$$  

In the case of Gaussian noise, one of the main results is the following.

**Theorem 1:** Consider the model (1) with $z \sim N(0, \sigma^2 I_n)$. If $\beta$ is $k$-sparse and

$$\delta_{1,25k} + \theta_{k,1,25k} < 1$$

then, an $\ell_1$ minimizer

$$\hat{\beta}^{DS} = \arg \min \left\{ \|y\|_1 : \|\Phi^T (y - \Phi \gamma)\|_\infty \leq \sigma \sqrt{2 \log p} \right\}$$

satisfies, with probability at least $1 - (1/(2\sqrt{\pi} \log p))$

$$\|\hat{\beta}^{DS} - \beta\|_2 \leq \frac{\sqrt{10}}{1 - \delta_{1,25k} - \theta_{k,1,25k}} \sqrt{k \sigma \sqrt{2\log p}}$$

(5)

and an $\ell_1$ minimizer

$$\hat{\beta}^{2} = \arg \min \left\{ \|y\|_1 : \|y - \Phi \gamma\|_2 \leq \sigma \sqrt{n + 2 \sqrt{n \log n}} \right\}$$

satisfies, with probability at least $1 - (1/n)$

$$\|\hat{\beta}^{2} - \beta\|_2 \leq \frac{2\sqrt{2(1 + \delta_{1,25k})}}{1 - \delta_{1,25k} - \theta_{k,1,25k}} \sigma \sqrt{n + 2 \sqrt{n \log n}}.$$  

(6)

See Section IV for a more general result. In comparison to [9, Th. 1.1], the result given in (5) for $\hat{\beta}^{DS}$ weakens the condition from $\delta_{3k} + \theta_{k,2k} < 1$ to $\delta_{1,25k} + \theta_{k,1,25k} < 1$ and improves the constant in the bound from $4/(1 - \delta_{2k} - \theta_{k,2k})$ to $\sqrt{10}/(1 - \delta_{1,25k} - \theta_{k,1,25k})$. Although our primary interest in this paper is to recover sparse signals, all the main results in the subsequent sections are given for general signals that are not necessarily $k$-sparse.

Weakening the RIP condition also has direct implications on the construction of compressed sensing (CS) matrices. It is important to note that it is computationally difficult to verify the RIP for a given design matrix $\Phi$ when $p$ is large and the sparsity $k$ is not too small. It is required to bound the condition numbers of $\binom{p}{k}$ submatrices. The spectral norm of a matrix is often difficult to compute and the combinatorial complexity makes it infeasible to check the RIP for reasonable values of $p$ and $k$. A general technique for avoiding checking the RIP directly is to generate the entries of the matrix $\Phi$ randomly and to show that the resulting random matrix satisfies the RIP with high probability using the well-known Johnson-Lindenstrauss Lemma. See, for example, Baraniuk, et al. [3]. Weakening the RIP condition makes it easier to prove that the resulting random matrix satisfies the CS properties.

The paper is organized as follows. After Section II, in which basic notations and definitions are reviewed, we introduce in Section III-A the elementary Shifting Inequality, which enables us to make finer analysis of the sparse recovery problem. We then consider the problem of exact recovery in the noiseless case in Section III-B and stable recovery of sparse signals in Section III-C. The Gaussian noise case is treated in Section IV. Section V discusses various conditions on RIP and effects of the improvement of the RIP condition on the construction of CS matrices. The proofs of some technical results are relegated to the Appendix.

**II. Preliminaries**

We begin by introducing basic notations and definitions related to the RIP. We also collect a few elementary results needed for the later sections.

For a vector $v = (v_i) \in \mathbb{R}^p$, we shall denote by $v_{\max}(k)$ the vector $v$ with all but the $k$ largest entries (in absolute value) set to zero and define $v_{\max}(k) = v - v_{\max}(k)'$, the vector $v$ with the $k$ largest entries (in absolute value) set to zero. We use the standard notation $\|v\|_q = (\sum_{i=1}^p |v_i|^q)^{1/q}$ to denote the $\ell_q$ norm of the vector $v$. We shall also treat a vector $v = (v_i)$ as a function $v : \{1, \ldots, p\} \rightarrow \mathbb{R}$ by assigning $v(i) = v_i$.

For a subset $T$ of $\{1, \ldots, p\}$, we use $\Phi_T$ to denote the submatrix obtained by extracting the columns of $\Phi$ according to the indices in $T$. Let $SSV_T = \{\lambda : \lambda$ an eigenvalue of $\Phi_T^T \Phi_T\}$, and $\Lambda_{\min}(k) = \min \{\bigcup_{|T| \leq k} SSV_T\}$, $\Lambda_{\max}(k) = \max \{\bigcup_{|T| \leq k} SSV_T\}$. It can be seen that

$$1 - \delta_k \leq \Lambda_{\min}(k) \leq \Lambda_{\max}(k) \leq 1 + \delta_k.$$  

Hence, (3) can be viewed as a condition on $\Lambda_{\min}(k)$ and $\Lambda_{\max}(k)$.

The following relations can be easily checked.

$$\delta_k \leq \delta_{k_1}, \text{ if } k \leq k_1 \leq p$$

$$\theta_{k,k'} \leq \theta_{k_1,k'_1}, \text{ if } k \leq k_1, k' \leq k_1', \quad k_1 + k'_1 \leq p.$$  

(7)

(8)

 Candès and Tao [8] showed that the constants $\delta_k$ and $\theta_{k,k'}$ are related by the following inequalities:

$$\theta_{k,k'} \leq \delta_{k+k'} \leq \theta_{k,k'} + \max(\delta_k, \delta_{k'}).$$  

(9)

Cai, Xu, and Zhang obtained the following properties for $\delta$ and $\theta$ in [5], which are especially useful in producing simplified recovery conditions:

$$\theta_k \sum_{i=1}^k k_i \leq \sqrt{\sum_{i=1}^k \theta_{k_i}^2 \leq \sqrt{\sum_{i=1}^k \delta_{k_i}^2}.$$  

(10)

It follows from (10) that for any positive integer $a$, we have $\theta_{k,ak'} \leq \sqrt{a} \theta_{k,k'}$. This can be further generalized.
Lemma 1: For any $a \geq 1$ and positive integers $k, k'$ such that $ak'$ is an integer

$$\theta_{k, ak'} \leq \sqrt{a} \theta_{k, k'}.$$  \hspace{1cm} (11)

A proof of this lemma can be found in the Appendix.

Consider the $\ell_1$ minimization problem $(P_B)$. Let $\beta$ be a feasible solution to $(P_B)$, i.e., $y = \Phi \beta \in B$. Without loss of generality we assume that supp($\beta_{\text{max}}(k)$) = $\{1, 2, \ldots, k\}$. Let $\tilde{\beta}$ be a solution to the minimization problem $(P_\tilde{B})$. Then it is clear that $||\beta||_1 \leq ||\tilde{\beta}||_1$. Let $h = \tilde{\beta} - \beta$ and $h_0 = h_{\{1, 2, \ldots, k\}}$ for some positive integer $k \leq p$. Here $\mathbb{I}_A$ denotes the indicator function of a set $A \subseteq \{1, 2, \ldots, p\}$, i.e., $\mathbb{I}_A(j) = 1$ if $j \in A$ and 0 if $j \not\in A$.

The following is a widely used fact. See, for example, [5], [7], [9], and [11].

Lemma 2: $||h - h_0||_1 \leq ||h_0||_1 + 2 ||\beta_{\text{max}}(k)||_1$.

This follows from the fact that $||\beta||_1 \geq ||\tilde{\beta}||_1 = ||\beta_{\text{max}}(k) + h_0||_1 + ||h - h_0 + \beta_{\text{max}}(k)||_1 \geq ||\beta_{\text{max}}(k)||_1 + ||h - h_0||_1 - ||\beta_{\text{max}}(k)||_1$.

Note also that the Cauchy-Schwarz Inequality yields that for $u \in \mathbb{R}^k$

$$||u||_1 \leq \sqrt{k} ||u||_2.$$  \hspace{1cm} (12)

III. SHIFTING INEQUALITY, EXACT AND STABLE RECOVERY

In this section, we consider exact recovery of high-dimensional sparse signals in the noiseless case and stable recovery in the bounded noise case. Recovery of sparse signals with Gaussian noise will be discussed in Section IV. We begin by introducing an elementary inequality which we call the Shifting Inequality. This useful inequality plays a key role in our analysis of the properties of the solution to the $\ell_1$ minimization problem.

A. The Shifting Inequality

The following elementary inequality enables us to perform finer estimation involving $\ell_1$ and $\ell_2$ norms as can be seen from the proofs of Theorem 2 in Section III-B and other main results.

Lemma 3 (Shifting Inequality): Let $q, r$ be positive integers satisfying $q \leq 3r$. Then any nonincreasing sequence of real numbers $a_1 \geq a_2 \geq \cdots \geq a_r \geq b_1 \geq \cdots \geq b_q \geq c_1 \geq \cdots \geq c_r \geq 0$ satisfies

$$\sqrt{\sum_{i=1}^{q} t_i^2 + \sum_{r=1}^{r} t_i^2} \leq \sqrt{\sum_{i=1}^{q} a_i + \sum_{i=1}^{r} b_i} \sqrt{q + r}.$$ \hspace{1cm} (13)

Equivalently, any nonincreasing sequence of real numbers $b_1 \geq \cdots \geq b_q \geq 0$ satisfies

$$\sqrt{\sum_{i=1}^{q} t_i^2 + r^2 t_i^2} \leq \sqrt{\sum_{i=1}^{q} a_i + \sum_{j=1}^{r} b_i} \sqrt{q + r}.$$ \hspace{1cm} (14)

The proof of this lemma is presented in the Appendix.

We will see that the Shifting Inequality, albeit very elementary, not only simplifies the analysis of $\ell_1$ minimization method but also weakens the required condition on the RIP.

B. Exact Recovery of Sparse Signals

We shall start with the simple setting where no noise is present. In this case the goal is to recover the signal $\beta$ exactly when it is sparse. This case is of significant interest in its own right as it is also closely connected to the problem of decoding of linear codes. See, for example, Candès and Tao [8]. The ideas used in treating this special case can be easily extended to treat the general case where noise is present.

Suppose $y = \Phi \gamma$. Based on $(\Phi, y)$, we wish to reconstruct the vector $\beta$ exactly when it is sparse. Equivalently, we wish to find the sparsest representation of the signal $y$ in the dictionary consisting of the columns of the matrix $\Phi$. Let $\tilde{\beta}$ be the minimizer to the problem

$$(\text{Exact}) \quad \min_{\gamma \in \mathcal{B}} ||\gamma||_1 \quad \text{subject to} \quad \Phi \gamma = y.$$ \hspace{1cm} (15)

Note that this is a special case of the $\ell_1$ minimization problem $(P_B)$ with $B = \{0\}$. We have the following result.

Theorem 2: Suppose that $\beta$ is $k$-sparse and that

$$\delta_{k+a} + \frac{1}{b} \sqrt{k} \theta_{k+a,b} < 1$$ \hspace{1cm} (16)

holds for some positive integers $a$ and $b$ satisfying $a < b \leq 4a$. Then the solution $\tilde{\beta}$ to the $\ell_1$ minimization problem (Exact) recovers $\beta$ exactly. In general, if (15) holds, then $\tilde{\beta}$ satisfies

$$||\beta - \tilde{\beta}||_2 \leq \frac{1 - \delta_{k+a} + \theta_{k+a,b} - 2}{1 - \delta_{k+a} - \sqrt{\frac{1}{b} \theta_{k+a,b}}} \frac{1}{\sqrt{k} \beta_{\text{max}}(k)}.$$ \hspace{1cm} (17)

Remark 1: We should note that in this and following main theorems, we use the general condition $\delta_{k+a} + \sqrt{k} \theta_{k+a,b} < 1$, which involves two positive integers $a$ and $b$ in addition to the sparsity parameter $k$. The flexibility in the choice of $a$ and $b$ in the condition allows one to derive interesting conditions for compressed sensing matrices. More discussions on special cases and comparisons with the existing conditions used in the current literature are given in Section V.

As noted earlier, for a real number $\alpha > 0$, $\delta_{\alpha k}$ and $\theta_{\alpha k}$ are meant to be $\delta_{[\alpha k]}$ and $\theta_{[\alpha k]}$. A particularly interesting choice is $b = k$ and $\alpha = \lfloor k/4 \rfloor$. Theorem 2 shows that if $\beta$ is $k$-sparse and

$$\delta_{1.25k} + \theta_{1.25k} < 1$$ \hspace{1cm} (18)

then the $\ell_1$ minimization method recovers $\beta$ exactly. This condition is weaker than other conditions on RIP currently available in the literature. Compare, for example, Candès and Tao [8], [9], Candès, Romberg, and Tao [7], Candès [6], and Cai,
Xu, and Zhang [5]. See more discussions in Section V. For a general signal \( \beta \), under (16), one has
\[
\| \hat{\beta} - \beta \|_2 \leq \frac{1 - \delta_{1,25k} + \theta_{k,1,25k}}{1 - \delta_{1,25k} - \theta_{k,1,25k}} \frac{2}{\sqrt{k}} \| \beta_{\text{max}}(k) \|_1.
\] (17)

**Proof:** The proof of Theorem 2 is elementary. The key to the proof is the Shifting Inequality. Again, set \( h = \hat{\beta} - \beta \). We shall cut the error vector \( h \) into pieces and then apply the Shifting Inequality to subvectors.

Without loss of generality, we assume the first \( k \) coordinates of \( \beta \) are the largest in magnitude. Making rearrangement if necessary, we may also assume that
\[
|h(k + 1)| \geq h(k + 2) \geq \cdots
\]

Set \( T_0 = \{1, 2, \cdots, k\} \), \( T_* = \{k + 1, k + 2, \cdots, k + a\} \) and \( T_i = \{k + a + (i - 1)b + 1, \cdots, k + a + ib\} \), \( i = 1, 2, \cdots \) with the last subset of size less than or equal to \( b \). Let \( h_0 = \Pi(T_0) \), \( h_* = \Pi(T_*) \), and \( h_i = \Pi(T_i) \) for \( i \geq 1 \).

\[
\begin{array}{ccccccc}
& h_0 & h_* & h_1 & h_2 & h_3 & \cdots \\
\text{Support Size} & k & a & b & b & b & \cdots
\end{array}
\]

To apply the Shifting Inequality, we shall first divide each vector \( h_i \) into two pieces. Set
\[
T_{i1} = \{k + a + (i - 1)b + 1, \cdots, k + ib\}
\]
and
\[
T_{i2} = T_i \setminus T_{i1} = \{k + 1 + ib, \cdots, k + a + ib\}.
\]

We note that \( |T_{i1}| = b - a \) and \( |T_{i2}| = a \) for all \( i \geq 1 \). Let \( h_{i1} = h_i|_{T_{i1}} \) and \( h_{i2} = h_i|_{T_{i2}} \).

It then follows from Lemma 2 and (12) that
\[
\sum_{i \geq 1} ||h_i||_2 \leq \frac{||h_*||_1 + \sum_{i \geq 1} ||h_i||_1}{\sqrt{b}} = \frac{||h - h_0||_1}{\sqrt{b}} \leq \frac{||h_0||_1 + 2 \sqrt{b} \| \beta_{\text{max}}(k) \|_1}{\sqrt{b}} \leq \sqrt{\frac{k}{b}} ||h_0||_2 + \frac{2 \sqrt{b} \| \beta_{\text{max}}(k) \|_1}{\sqrt{b}}.
\]

Now the fact that \( \Phi h_i = 0 \) yields
\[
0 = \langle \Phi h_i, \Phi(h_0 + h_*) \rangle = \langle \Phi(h_0 + h_*), \Phi(h_0 + h_*) \rangle + \sum_{i \geq 1} \langle \Phi h_i, \Phi(h_0 + h_*) \rangle \geq (1 - \delta_{k+a}) ||h_0 + h_*||_2^2 - \sum_{i \geq 1} \theta_{k+a,b} ||h_0 + h_*||_2 ||h_i||_2
\]
\[
\geq ||h_0 + h_*||_2 \left( 1 - \delta_{k+a} - \sqrt{\frac{k}{b}} \theta_{k+a,b} \right) ||h_0 + h_*||_2
\]
\[
- ||h_0 + h_*||_2 \theta_{k+a,b} \frac{2 \| \beta_{\text{max}}(k) \|_1}{\sqrt{b}}.
\]

This implies
\[
||h_0 + h_*||_2 \leq \frac{\theta_{k+a,b}}{1 - \delta_{k+a} - \sqrt{\frac{k}{b}} \theta_{k+a,b}} \frac{2 \| \beta_{\text{max}}(k) \|_1}{\sqrt{b}}.
\]

Therefore
\[
||h||_2 \leq ||h_0 + h_*||_2 + \sum_{i \geq 1} ||h_i||_2
\]
\[
\leq \left( 1 + \sqrt{\frac{k}{b}} \right) ||h_0 + h_*||_2 + \frac{2 \| \beta_{\text{max}}(k) \|_1}{\sqrt{b}}
\]
\[
\leq \frac{1 - \delta_{k+a} + \theta_{k+a,b}}{1 - \delta_{k+a} - \sqrt{\frac{k}{b}} \theta_{k+a,b}} \frac{2 \| \beta_{\text{max}}(k) \|_1}{\sqrt{b}}.
\]

If \( \beta \) is \( k \)-sparse, then \( \beta_{\text{max}}(k) = 0 \), which implies \( \beta = \hat{\beta} \).

The key argument used in the proof of Theorem 2 is the Shifting Inequality. This simple analysis requires a condition on the RIP that is weaker than other conditions on the RIP used in the literature.

In addition to Theorem 2, we also have the following result under a simpler condition.

**Theorem 3:** Let \( k \) be a positive integer. Suppose \( \beta \) is \( k \)-sparse and
\[
\delta_k + \sqrt{k} \theta_{k,1} < 1
\] (18)
then the \( \ell_1 \) minimization recovers \( \beta \) exactly.
Proof: The proof is similar to that of Theorem 2. For each $i \geq 1$, let $T_i = \{k + i\}$ and $h_{ki} = h|_{T_i}$. Note that $\Phi h = 0$. Hence,

$$0 = |\langle \Phi h, \Phi h_0 \rangle| = \left| \langle \Phi h_0, \Phi h_0 \rangle + \sum_{i \geq 1} \langle \Phi h_i, \Phi h_0 \rangle \right| \geq (1 - \delta_k) ||h_0||^2 - \sum_{i \geq 1} \theta_{ki,1} ||h_i||^2 = ||h_0||^2 \left( (1 - \delta_k) ||h_0||^2 - \theta_{ki,1} \sum_{i \geq 1} ||h_i||^2 \right) = ||h_0||^2 \left( (1 - \delta_k) ||h_0||^2 - \theta_{ki,1} \sum_{i \geq 1} ||h_i||^2 \right) \geq ||h_0||^2 \left( (1 - \delta_k) ||h_0||^2 - \theta_{ki,1} ||h_0||^2 \right) \geq (1 - \delta_k - \sqrt{k} \theta_{ki,1}) ||h||^2. \tag{20}$$

Since $1 - \delta_k - \sqrt{k} \theta_{ki,1} > 0$, this implies $h_0 = 0$ and, hence, $h = 0$. ■

Remark 2: Another commonly used condition in the sparse recovery literature is the mutual incoherence property (MIP), which can be seen from the facts that $h_0$ is the columns of $\Phi$, and the columns of $\Phi (\Phi^*)$’s are also assumed to be of length 1 in $\ell_2$-norm. A sharp MIP condition for the stable recovery of $k$-sparse signals in the noisy case is

$$(2k - 1) \mu < 1. \tag{21}$$

See [4]. It is easy to check that (18) is weaker than (20). This can be seen from the facts that $\delta_k \leq (k - 1) \mu$ and $\theta_{ki,1} \leq \sqrt{k} \mu$. (See, e.g., [5].)

For a general signal $\beta$, a slightly modified proof yields that the $\ell_1$ minimizer $\hat{\beta}$ satisfies

$$||\hat{\beta} - \beta||_2 \leq 2 \frac{(1 - \delta_k + \theta_{ki,1})}{1 - \delta_k - \sqrt{k} \theta_{ki,1}} \|\beta\|_{\ell_\infty} \tag{22}$$

under (18). Note that this bound is not as strong as the error bound given in (17) obtained under (16).

C. Recovery in the Presence of Errors

We now consider reconstruction of high dimensional sparse signals in the presence of bounded noise. Let $B \subset \mathbb{R}^n$ be a bounded set. Suppose we observe $(\Phi y)$ where $y = \Phi \beta + z$ with the error vector $z \in B$, and we wish to reconstruct $\beta$ by solving the $\ell_1$ minimization problem $(P_B)$. Specifically, we consider two types of bounded errors: $B_1(\eta) = \{z : \|\Phi^T z\|_\infty \leq \eta\}$ and $B_2(\eta) = \{z : \|z\|_2 \leq \eta\}$. We shall use $\beta^{DS}$ to denote the solution of the $\ell_1$ minimization problem $(P_B)$ with $B = B_1(\eta)$ and use $\beta^{\ell_2}$ to denote the solution of $(P_B)$ with $B = B_2(\eta)$.

The Shifting Inequality again plays a key role in our analysis in this case. In addition, the analysis of the Gaussian noise case follows easily from that of the bounded noise case.

Theorem 4: Suppose

$$\delta_{k+a} + \sqrt{\frac{k}{b}} \theta_{k+a,b} < 1 \tag{23}$$

holds for positive integers $k$, $a$, and $b$ where $a \leq b \leq 4a$. Then the minimizers $\beta^{DS}$ and $\beta^{\ell_2}$ satisfy

$$||\beta^{DS} - \beta||_2 \leq A \eta + B \|\beta\|_{\ell_\infty} \tag{24}$$

and

$$||\beta^{\ell_2} - \beta||_2 \leq C \eta + B \|\beta\|_{\ell_\infty} \tag{25}$$

where

$$A = 2 \sqrt{1 + \frac{k}{b} \frac{\sqrt{k + a}}{1 - \delta_{k+a} - \sqrt{k} \theta_{k+a,b}}} \tag{26}$$

$$B = \frac{2}{\sqrt{b}} \left( 1 - \frac{\theta_{k+a,b} \sqrt{1 + k/b}}{1 - \delta_{k+a} - \sqrt{k} \theta_{k+a,b}} \right) \tag{27}$$

$$C = 2 \sqrt{1 + \frac{k}{b} \frac{\sqrt{1 + \delta_{k+a}}}{1 - \delta_{k+a} - \sqrt{k} \theta_{k+a,b}}} \tag{28}$$

A proof of Theorem 4 based on the ideas of that for Theorem 2 is given in the Appendix.

Remark 3: As in the noiseless setting, an especially interesting case is $b = k$ and $a = [k/4]$. In this case, Theorem 4 yields that if $\beta$ is $k$-sparse and

$$\delta_{1.25k} + \theta_{k,1.25k} < 1 \tag{29}$$

holds, then the $\ell_1$ minimizers $\beta^{DS}$ and $\beta^{\ell_2}$ satisfy

$$||\beta^{DS} - \beta||_2 \leq \frac{\sqrt{10}}{1 - \delta_{1.25k} - \theta_{k,1.25k}} \cdot \sqrt{10} \tag{30}$$

and

$$||\beta^{\ell_2} - \beta||_2 \leq \frac{2 \sqrt{1 + \delta_{1.25k}}}{1 - \delta_{1.25k} - \theta_{k,1.25k}} \cdot \eta. \tag{31}$$

Again, (25) for stable recovery in the noisy case is weaker than the existing RIP conditions in the literature. See, for example, [9], [7], [61], and [5].

IV. GAUSSIAN NOISE

The Gaussian noise case is of particular interest in statistics and several methods have been developed. See, for example, [16], [12], and [9]. The results presented in Section III-C on the bounded noise case are directly applicable to the case where the noise is Gaussian. This is due to the fact that Gaussian noise is “essentially bounded.” Suppose we observe

$$y = \Phi \beta + z, \quad z \sim \mathcal{N}(0, \sigma^2 I_n) \tag{32}$$

where $\mathcal{N}(0, \sigma^2 I_n)$ is a standard multivariate Gaussian distribution with mean $0$ and covariance $\sigma^2 I_n$. The solution of the $\ell_1$ minimization problem (P) with $B = B_1(\eta)$ and $B_2(\eta)$ is denoted by $\beta^{DS}$ and $\beta^{\ell_2}$, respectively.

Theorem 5: Suppose

$$\delta_{k+a} + \frac{\sqrt{k}}{\sigma} \theta_{k+a,b} < 1 \tag{33}$$

holds for positive integers $k$, $a$, and $b$ where $a \leq b \leq 4a$. Then the minimizers $\beta^{DS}$ and $\beta^{\ell_2}$ satisfy

$$||\beta^{DS} - \beta||_2 \leq A \eta + B \|\beta\|_{\ell_\infty} \tag{34}$$

and

$$||\beta^{\ell_2} - \beta||_2 \leq C \eta + B \|\beta\|_{\ell_\infty} \tag{35}$$

where

$$A = 2 \sqrt{1 + \frac{k}{b} \frac{\sqrt{k + a}}{1 - \delta_{k+a} - \frac{\sqrt{k}}{\sigma} \theta_{k+a,b}}} \tag{36}$$

$$B = \frac{2}{\sqrt{b}} \left( 1 - \frac{\theta_{k+a,b} \sqrt{1 + k/b}}{1 - \delta_{k+a} - \frac{\sqrt{k}}{\sigma} \theta_{k+a,b}} \right) \tag{37}$$

$$C = 2 \sqrt{1 + \frac{k}{b} \frac{\sqrt{1 + \delta_{k+a}}}{1 - \delta_{k+a} - \frac{\sqrt{k}}{\sigma} \theta_{k+a,b}}} \tag{38}$$

A proof of Theorem 5 based on the ideas of that for Theorem 2 is given in the Appendix.

Remark 4: As in the noiseless setting, an especially interesting case is $b = k$ and $a = [k/4]$. In this case, Theorem 5 yields that if $\beta$ is $k$-sparse and

$$\delta_{1.25k} + \frac{\sqrt{k}}{\sigma} \theta_{k,1.25k} < 1 \tag{39}$$

holds, then the $\ell_1$ minimizers $\beta^{DS}$ and $\beta^{\ell_2}$ satisfy

$$||\beta^{DS} - \beta||_2 \leq \frac{\sqrt{10}}{1 - \delta_{1.25k} - \frac{\sqrt{k}}{\sigma} \theta_{k,1.25k}} \cdot \sqrt{10} \tag{40}$$

and

$$||\beta^{\ell_2} - \beta||_2 \leq \frac{2 \sqrt{1 + \delta_{1.25k}}}{1 - \delta_{1.25k} - \frac{\sqrt{k}}{\sigma} \theta_{k,1.25k}} \cdot \eta. \tag{41}$$

Again, (39) for stable recovery in the noisy case is weaker than the existing RIP conditions in the literature. See, for example, [9], [7], [61], and [5].
and wish to recover the signal $\beta$ based on $(\Phi, y)$. We assume that $\sigma$ is known and that the columns of $\Phi$ are normalized to have unit $\ell_2$ norm. Define two bounded sets

$$B_3 = \{z : ||\Phi^T z||_{\infty} \leq \sigma \sqrt{2 \log p}\} \quad \text{and}$$

$$B_4 = \{z : ||z||_{2} \leq \sigma \sqrt{n + 2 \sqrt{n \log n}}\}.$$

(29)

The following result, which follows from standard probability calculations, shows that Gaussian noise is essentially bounded. The readers are referred to [5] for a proof.

**Lemma 4:** The Gaussian error $z \sim N(0, \sigma^2 I_n)$ satisfies

$$P(z \in B_3) \geq 1 - \frac{1}{2 \sqrt{\pi \log p}} \quad \text{and} \quad P(z \in B_4) \geq 1 - \frac{1}{n}.$$  

(30)

Lemma 4 indicates that the Gaussian variable $z$ is in the bounded sets $B_3$ and $B_4$ with high probability. The results obtained in the previous sections for bounded errors can, thus, be applied directly to treat Gaussian noise. In this case, we shall consider two particular constrained $\ell_1$ minimization problems. Let $\hat{\beta}_{DS}^2$ be the minimizer of

$$\min_{\gamma \in \mathbb{R}^p} ||\gamma||_1 \quad \text{subject to} \quad y - \Phi \gamma \in B_3$$

(31)

and let $\hat{\beta}_{\ell^2}^2$ be the minimizer of

$$\min_{\gamma \in \mathbb{R}^p} ||\gamma||_1 \quad \text{subject to} \quad y - \Phi \gamma \in B_4.$$  

(32)

The following theorem is a direct consequence of Lemma 4 and Theorem 4.

**Theorem 5:** Suppose

$$\delta_{k+a} + \sqrt{k/b} \theta_{k+a,b} < 1$$

(33)

holds for some positive integers $k$, $a$, and $b$ with $a < b \leq 4a$. Then with probability at least $1 - (1/(2 \sqrt{\pi \log p}))$, the minimizer $\hat{\beta}_{DS}^2$ satisfies

$$||\hat{\beta}_{DS}^2 - \beta||_2 \leq A \sigma \sqrt{2 \log p} + B ||\beta_{\max}(k)||_1$$

and with probability at least $1 - (1/n)$, the minimizer $\hat{\beta}_{\ell^2}^2$ satisfies

$$||\hat{\beta}_{\ell^2}^2 - \beta||_2 \leq C \sigma \sqrt{n + 2 \sqrt{n \log n}} + B ||\beta_{\max}(k)||_1$$

where the constants $A$, $B$, and $C$ are given in as in Theorem 4.

**Remark:** Again, a special case is $b = k$ and $a = \lfloor k/4 \rfloor$. In this case, if $\beta$ is $k$-sparse and

$$\delta_{1.25k} + \theta_{k,1.25k} < 1$$

then, with high probability, the $\ell_1$ minimizers $\hat{\beta}_{DS}^2$ and $\hat{\beta}_{\ell^2}^2$ satisfy

$$||\hat{\beta}_{DS}^2 - \beta||_2 \leq \frac{\sqrt{10}}{1 - \delta_{1.25k} - \theta_{k,1.25k}} \sqrt{k \sigma \sqrt{2 \log p}}$$

(34)

$$||\hat{\beta}_{\ell^2}^2 - \beta||_2 \leq \frac{2 \sqrt{2(1 + \delta_{1.25k})}}{1 - \delta_{1.25k} - \theta_{k,1.25k}} \sigma \sqrt{n + 2 \sqrt{n \log n}}.$$  

(35)

The result given in (34) for $\hat{\beta}_{DS}^2$ improves [9, Th. 1.1] by weakening the condition from $\delta_{2k} + \theta_{k,2k} < 1$ to $\delta_{1.25k} + \theta_{k,1.25k} < 1$ and reducing the constant in the bound from $4/(1 - \delta_{2k} - \theta_{k,2k})$ to $\sqrt{10}/(1 - \delta_{1.25k} - \theta_{k,1.25k})$. The improvement on the error bound is minor. The improvement on the condition is more significant as it shows signals with larger support can be recovered accurately for fixed $\eta$ and $p$.

Candès and Tao [9] also derived an oracle inequality for $\hat{\beta}_{DS}^2$ in the Gaussian noise setting under the condition $\delta_{2k} + \theta_{k,2k} < 1$. Our method can also be used to improve [9, Ths. 1.2 and 1.3] by weakening the condition to $\delta_{1.25k} + \theta_{k,1.25k} < 1$.

V. DISCUSSIONS

The flexibility in the choice of $a$ and $b$ in the condition $\delta_{k+a} + \sqrt{k/b} \theta_{k+a,b} < 1$ used in Theorems 2, 4, and 5 enables us to deduce interesting conditions for compressed sensing matrices. We shall highlight several of them here and compare with the existing conditions used in the current literature. As mentioned in the introduction, it is sometimes more convenient to use conditions only involving the restricted isometry constant $\delta$ and for this reason we shall mainly focus on $\delta$. By choosing different values of $a$ and $b$ and using (10), it is easy to show that each of the following conditions is sufficient for the exact recovery of $k$-sparse signals in the noiseless case and stable recovery in the noisy case:

1) $\delta_{1.25k} + \theta_{k,1.25k} < 1$;
2) $\delta_{1.25k} + \sqrt{1.25\delta_{2k}} < 1$;
3) $\delta_{1.625k} < \sqrt{2} - 1 \approx 0.414$;
4) $\delta_{4k} < (1/(1 + \sqrt{1.25})) \approx 0.472$;
5) $\delta_{2k} < 2(2 - \sqrt{3}) \approx 0.535$;
6) $\delta_{k} < 2 - \sqrt{2} \approx 0.585$.

The derivation of Condition 1 has been discussed in the remarks of Theorem 2 and Theorem 4.

By Lemma 1 and (9), we have $\theta_{k,1.25k} \leq \sqrt{1.25k} \theta_{k,k} \leq \sqrt{1.25k}$. Therefore, Condition 1 implies Condition 2.

Condition 5 follows from Condition 1 and (10). In fact, if $\delta_{1.625k} < \sqrt{2} - 1$, then

$$\delta_{1.25k} + \theta_{k,1.25k} \leq \delta_{1.625k} + \sqrt{\delta_{k+0.025k}^2 + \delta_{k+0.025k}^2} < 1.$$  

Condition 4 is stronger than Condition 2. This is because if $\delta_{2k} < (1/(1 + \sqrt{1.25}))$, then

$$\delta_{1.25k} + \sqrt{1.25\delta_{2k}} < \delta_{2k} + \sqrt{1.25\delta_{2k}} < 1.$$  

To get Condition 5, let $a = 0.5k$, $b = 2k$. The condition $\delta_{k+a} + \sqrt{k/b} \theta_{k+a,b} < 1$ becomes $\delta_{1.5k} + \sqrt{2} \theta_{1.5k,2k} < 1$. This condition is met if $\delta_{3k} < 2(2 - \sqrt{3})$. In fact, by Lemma 1 and (9)

$$\delta_{1.5k} + \sqrt{1} \theta_{1.5k,2k} \leq \delta_{3k} + \sqrt{\frac{3}{2}} \theta_{k,2k} \leq \left(1 + \frac{\sqrt{3}}{2}\right) \delta_{3k} < 1.$$  

Condition 5 can be obtained in a similar manner. These conditions for stable recovery improve the conditions used in the literature, e.g., the conditions $\delta_{3k} + 3\delta_{4k} < 2$ in [7],

2As $\delta_{k+2k,2k}$ is for $\theta_{k+2k,2k}$, this condition should actually read $\delta_{1.25k} + \sqrt{1.25k}k\theta_{k+2k,2k} < 1$. There is a similar explanation for Condition 4.
\( \delta_{2k} + \theta_{2k} < 1 \) in [9], \( \delta_{1,2k} + \theta_{1,2k} < 1 \) in [5], \( \delta_{2k} < \sqrt{2} - 1 \) in [6], \( \delta_{1,72k} < \sqrt{2} - 1 \) in [5], and \( \delta_{2k} < 0.4531 \) in [13]. It is also interesting to note that Condition 5 allows \( \delta_{2k} \) to be larger than 0.5.

The flexibility in \( \delta_{k+a} + \sqrt{k/b} \theta_{k+a} < 1 \) also enables us to discuss the asymptotic properties of the RIP conditions. Letting \( a = tk; b = 4tk \) and using (7), (8), and (10), it is easy to infer that each of the following conditions is sufficient for stable recovery of \( k \)-sparse signals:

1) \( \delta_{(3t+1)k} < \sqrt{(2t(1+\sqrt{2t}))}, t \geq (1/3) \)
2) \( \delta_{((9t+1)/2)k} < \sqrt{(2t(1+\sqrt{2t}))}, t \in ((1/7), (1/3)) \)
3) \( \delta_{(1+t)k} < \sqrt{(2t(1+\sqrt{2t}))}, t \in (0, (1/7)] \)

These conditions reveal two asymptotic properties of the restricted isometry constant \( \delta \). The first is that \( \delta \) can be close to 1 (as \( t \) gets large), provided that checking the RIP for \( \Phi \) must be done for sets of columns whose cardinality is much bigger than \( k \), the sparsity for recovery. The second is that if \( \delta \) is allowed to be small, then checking the RIP for \( \Phi \) can be done for sets of columns whose cardinality is close to \( k \) (as \( t \) gets small). It is noted that checking the RIP remains quite impractical even in this case.

It is clear that with weaker RIP conditions, more matrices can be verified to be compressed sensing matrices. As aforementioned, for a given \( n \times p \) matrix, it is computationally difficult to check its restricted isometry property. However, it has been very successful in constructing random compressed sensing matrices which satisfy the RIP conditions with high probability, see [2], [3], [7], [8], [9], [15].

For example, Baronii et al. [3] showed that if \( \Phi \) is an \( n \times p \) matrix whose entries are drawn independently according to Gaussian distribution \( (\Phi_{i,j} \sim N(0, (1/n))) \) or Bernoulli distribution \( (\Phi_{i,j} \text{ is either } 1/\sqrt{n} \text{ or } -(1/\sqrt{n}), \text{ each with probability } 1/2) \), then \( \Phi \) fails to have RIP

\[ \delta_{p} < \alpha \]

with probability less than

\[ \tau_{\alpha} = 2e^{-\frac{1}{3}(\frac{2\alpha^{2}}{\sqrt{2}})}n \left( \frac{cP}{\alpha^{3}} \right)^{r} \]

where \( c \) is a constant.

It is not hard to see that the probability of failing drops at a considerable rate as the bound \( \alpha \) increases and/or the index \( r \) decreases. In fact, with a weaker condition \( \delta_{p} < \alpha + \epsilon < 1 \), this rate is

\[ \tau_{\alpha} = \frac{1}{\alpha^{r}} \left( 1 + e^{\frac{1}{\alpha^{r}}} \right) \]

This ratio is very large if \( n \) is large.

On the other hand, the improvement of RIP conditions can be interpreted as enlarging the sparsity of the signals to be recovered. For example, one of the previous results showed that the condition \( \delta_{2k} < \sqrt{2} - 1 \) ensures the recovery of a \( k \)-sparse signal. Replacing the condition by \( \delta_{1,252k} < \sqrt{2} - 1 \), we see that the sparsity of the signals to be recovered is relaxed \( 2(1/1.252) \approx 1.23 \) times.

**Appendix A**

**Proof of Lemma 1**

We just need to prove, for any vector \( c_1, c' \in \mathbb{R}^p \) with disjoint supports and sparsity \( k \) and \( ak' \), respectively, that

\[ |\langle \Phi c_1, \Phi c' \rangle| \leq \sqrt{a} \| \theta_{k,k'} \|_{2} \| c \|_{2} \| c' \|_{2}. \]

Without loss of generality, we assume that the support of \( c' \) is \( \{ 1, 2, \ldots, ak' \} \). For \( 1 \leq j \leq ak' \), when the \( ak' + j \)th coordinate of \( c' \) is mentioned, we actually mean the \( j \)th one.

For \( i = 1, 2, \ldots, ak' \), let \( c'_{i} \in \mathbb{R}^{p} \) be a vector such that \( c'_{i} \) keeps the \( i \)th, \((i+1)\)th, \ldots, \((a+1)\)th nonzero coordinates of \( c' \) and replaces other coordinates by zero.

\[ |\langle \Phi c_1, \Phi c' \rangle| = \left| \left\langle \Phi c_1, \Phi \left( \frac{1}{k} \sum_{i=1}^{ak'} c'_{i} \right) \right\rangle \right| \]

\[ \leq \frac{1}{k} \sum_{i=1}^{ak'} \left| \langle \Phi c_1, \Phi c'_{i} \rangle \right| \]

\[ \leq \frac{1}{k} \| \theta_{k,k'} \|_{2} \sum_{i=1}^{ak'} \| c'_{i} \|_{2} \]

\[ \leq \frac{1}{k} \| \theta_{k,k'} \|_{2} \sqrt{\frac{ak'}{\sum_{i=1}^{ak'} \| c'_{i} \|_{2}^{2}}} \]

\[ = \sqrt{a} \| \theta_{k,k'} \|_{2} \| c \|_{2} \| c' \|_{2}. \]

**Appendix B**

**Proof of the Shifting Inequality (Lemma 3)**

Let \( b_i = b_{i+1} + d_i \) for \( i = 1, 2, \ldots, q - 1 \). Then

\[ (q+r) \left( \sum_{i=1}^{q} b_{i}^{2} + r b_{q}^{2} \right) = (q+r)^2 \left( \sum_{i=1}^{q} d_{i}^{2} + r b_{q}^{2} \right) \]

\[ = (q+r)^2 b_{i}^{2} + 2(q+r)b_{q} \sum_{j=1}^{q-1} d_{j} + (q+r) \sum_{j=1}^{q-1} d_{j}^{2} \]

and

\[ r b_{i} + \sum_{i=1}^{q-1} d_{i}^{2} \]

\[ = (q+r)^2 b_{i}^{2} + 2(q+r)b_{q} \sum_{j=1}^{q-1} d_{j} + (q+r) \sum_{j=1}^{q-1} d_{j}^{2} \]

Since \( d_i \) is nonnegative for all \( i \), we know that

\[ 2(q+r) \sum_{i=1}^{q-1} (r+i)d_{i} \geq 2(q+r) \sum_{i=1}^{q-1} (r+i)d_{i}. \]

Also, it can be seen that for any \( 1 \leq i \leq j \leq q - 1 \), the coefficient of \( d_{i}d_{j} \) in \( \left( \sum_{i=1}^{q-1} (r+i)d_{i} \right)^{2} \) is \((1+I(i \neq j))(r+i)(r+i+1)\).
And the coefficient of $d_{l}d_{j}$ in $(q + r) \sum_{i=1}^{q-1} \sum_{i=j}^{r} d_{i}^{2}$ is $(1 + I(i \neq j))(q + r)i$. Since $q \leq 3r$, we know that

$$(r + i)(r + j) \geq (r + i)^{2} = (r - i)^{2} + 4ri \geq (q + r)i.$$ 

This means

$$\left( \sum_{i=1}^{q-1} (r + i)d_{i} \right)^{2} \geq (q + r) \left( \sum_{j=1}^{r} d_{j} \right)^{2}.$$ 

Hence, the inequality is proved.

**APPENDIX C**

**PROOF OF THEOREM 4**

Similar to the proof of theorem 2, we have

$$|\langle \Phi h, \Phi h_{0} + h_{a} \rangle| \geq \|h_{0} + h_{a}\|_{2} \left( 1 - \delta_{k+a} - \sqrt{\frac{k}{b}} \theta_{k+a,b} \right)$$

$$\times \|h_{0} + h_{a}\|_{2} - \|h_{0} + h_{a}\|_{2} \theta_{k+a,b} \frac{2 \|\beta_{\text{max}}(k)\|_{1}}{\sqrt{b}}.$$ 

**Case I**. $\mathcal{B} = \{ z : \| z \|_{2} \leq \eta \}$: It is easy to see that

$$|\langle \Phi h, \Phi (h_{0} + h_{a}) \rangle| \leq \|h_{0} + h_{a}\|_{2} \|\Phi (h_{0} + h_{a})\|_{2}$$

$$\leq 2\eta \sqrt{1 + \delta_{k+a}} \|h_{0} + h_{a}\|_{2}.$$ 

Therefore

$$\|h_{0} + h_{a}\|_{2} \leq \frac{2\eta \sqrt{1 + \delta_{k+a}} + \frac{2\theta_{k+a,b}}{\sqrt{b}} \|\beta_{\text{max}}(k)\|_{1}}{1 - \delta_{k+a} - \sqrt{\frac{k}{b}} \theta_{k+a,b}}.$$ 

Now

$$\|h\|_{2}^{2} = \|h_{0} + h_{a}\|_{2}^{2} + \sum_{i \geq 1} \|h_{i}\|_{2}^{2}$$

$$\leq \|h_{0} + h_{a}\|_{2}^{2} + \left( \sum_{i \geq 1} \|h_{i}\|_{2}^{2} \right)^{2}$$

$$\leq \|h_{0} + h_{a}\|_{2}^{2} + \left( \frac{k}{b} \|h_{0} + h_{a}\|_{2} + \frac{2 \|\beta_{\text{max}}(k)\|_{1}}{\sqrt{b}} \right)^{2}$$

$$= \left( \sqrt{1 + \frac{k}{b}} \|h_{0} + h_{a}\|_{2} + \frac{2 \|\beta_{\text{max}}(k)\|_{1}}{\sqrt{b}} \right)^{2}$$

$$\leq (C\eta + B \|\beta_{\text{max}}(k)\|_{1})^{2}.$$ 

**Case II**. $\mathcal{B} = \{ z : \| \Phi z \|_{\infty} \leq \eta \}$: By assumption, there is a $z \in \mathcal{B}$ such that $\Phi z = y - z$. So

$$|\langle \Phi h, \Phi (h_{0} + h_{a}) \rangle| = \|\Phi (\tilde{y} + z, \Phi_{T_{0},T_{n}}(h_{0} + h_{a}))\|_{2}$$

$$\leq \|\Phi_{T_{0},T_{n}}(\tilde{y} + z)\|_{2} \|h_{0} + h_{a}\|_{2}$$

$$\leq 2\sqrt{k + a\eta} \|h_{0} + h_{a}\|_{2}.$$ 

This implies

$$\|h_{0} + h_{a}\|_{2} \leq \frac{2\eta \sqrt{k + a + \frac{2\eta \theta_{k+a,b}}{\sqrt{b}} \|\beta_{\text{max}}(k)\|_{1}}{1 - \delta_{k+a} - \sqrt{\frac{k}{b}} \theta_{k+a,b}}.$$ 

Similar to Case I, we have

$$\|h\|_{2}^{2} \leq \left( \frac{1}{1 + \frac{k}{b}} \|h_{0} + h_{a}\|_{2} + \frac{2 \|\beta_{\text{max}}(k)\|_{1}}{\sqrt{b}} \right)^{2}.$$ 

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