DESIGN OF A NETWORKED CD-ROM FOR MULTIMEDIA APPLICATIONS

by

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Abstract

Allowing users of multimedia applications access to a large library of such items via the networking of PC CD-ROM drives (Net-CD) is investigated in this paper. A model for multimedia systems and an architecture for the implementation of this model are discussed. A simple Net-CD system is designed, constructed, and tested to determine the feasibility of networking CD-ROM drives on a PC.

Thesis Supervisor: Stephen Burns Title: Technical Director Of Biomedical Engineering Center

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Chapter 1

Introduction

Multimedia is the term for applications which appeal to more than one of the senses. These applications require a high-bandwidth information flow to the user. There is a wide variety of such applications available for PCs on CD-ROM. The amount of information available on one CD-ROM is hundreds of times greater than the capacity of standard 3-1/ 2 inch magnetic diskettes. Information in this large quantity lends itself to use in many types of applications such as video games, encyclopedias, video shopping, and the like.It would be an advantage to the user of multimedia CD-ROM software if his or her PC was networked to a large quantity of CD-ROM type applications. This paper discusses solutions to the networking of a PC CD-ROM and the implementation of one such system.

Chapter 2

Multimedia Systems

2.1 A Model for Multimedia Systems

Figure 2.1 diagrams a model of current multimedia systems. This model shows the general asymmetry of bandwidth between the user and the multimedia information source. Although one can envision a system in which the user input was of a high bandwidth, such as an application with video input, today's multimedia applications have low bandwidth user input such as a mouse, keyboard or remote control. The high bandwidth output of the system is required by applications which output video or audio. This bandwidth asymmetry plays a key role in the system architecture that will be discussed.

The system model in Figure 2.1 is separated into the user interface and the multimedia information source. One example of a user interface is a personal computer, capable of playing audio and video, through which multimedia applications a used. Another is the combination of a television (for audio and video output) and a set-top (for user interface) which would be used for interactive multimedia applications, including video on demand.

By separating the model in this manner the user interface can be made almost independent of the content of the multimedia experience. The exception to this is the necessity of a standard in terms of the data communication and the types of input and output media. In this way, the user can have access to a large multimedia library through a constant interface.

The multimedia information source can be represented in many ways. The source can be a video library, a source of processing power for applications such as central computing and video games, a CD-ROM, or any source of multimedia applications. The information source can also be implemented as a network of many different information providers (IP)

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which will provide the user with a variety of multimedia and will also allow for competition between IP's.



Figure 2.1: Multimedia System Model

2.2 Video Dial Tone Architecture

The Video Dial Tone (VDT) multimedia system architecture takes advantage of, (1) a network already in existence and (2) the asymmetric nature of contemporary multimedia.

The telephone network is used as the channel between the user interface and information providers. By the use of this network, ordinary telephone subscribers will already have access to a data communications channel. The disadvantage of this system is that the user is connected to the network via twisted pair copper wire. The twisted pair, in its current configuration, has quite a low bandwidth relative to the requirements of multimedia output.

The solution to this problem comes in the form of a technology known as the Asymmetric Digital Subscriber Line (ADSL). ADSL allows the normally low bandwidth twisted pair to act as a data communications channel with low bandwidth in one direction (9600 bits/sec) and a high bandwidth in the opposite direction (1.522 Megabits/sec). Because of the inherent asymmetry of today's multimedia applications, ADSL allows the telephone network to be used as a multimedia information source with a small modification to the twisted pair terminations at the subscriber and at the central office (CO).

A diagram of the VDT system architecture is shown in Figure 2.2. The ADSL twisted pair line links the user interface hardware, or Customer Premises Equipment (CPE), with a dedicated VDT line card in the central office. Through the CO, the line card is connected to other CO's and the information providers through a Synchronous Optical NETwork (SONET) using the Asynchronous Transfer Mode (ATM) protocol. SONET is a high bandwidth fiber-optic network which links distant portions of the telephone network.

The VDT line card can be considered the dividing line between the user interface and the multimedia information source in the model shown in Figure 2.1. This line card contains a processor capable of running scripts downloaded into it, as well as a buffer used as a temporary storage area for multimedia information. The line card is capable of executing protocols necessary for communication with both the user interface and the IP. Because of the much higher bandwidth of SONET relative to the ADSL, multimedia information can

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be sent from the IP in large bursts, buffered in the line card and then sent over the twisted pair at the proper speed. This allows SONET to be time-multiplexed on a per-user basis.

The VDT architecture allows a currently existing network to implement a multimedia information source with little modification to the existing line. This will give users an easy connection to a large library of multimedia titles.



Information Provider



Chapter 3

Design of Networked CD-ROM

3.1 Networked CD-ROM on VDT

One application that would be well suited for VDT would be that of supplying the user with a large library of PC CD-ROM titles. A diagram for the implementation of such a system is shown in Figure 3.1. The CPE for the networked CD-ROM (Net-CD) would consist of the customer's personal computer and an interface peripheral. The interface would create a data link between the PC and line card, through the SONET and eventually to the IP.

There are many ways to implement the CD-ROM titles in the IP. The IP could be simply another PC with a CD-ROM drive and an interface to the network. A more practical solution might consist of a server which has fast access to many CD-ROM titles. The line card might buffer and cache data in order to improve performance.

The user load on the IP is limited by the bandwidth of its connection to the network. One 100 Megabit/sec line, for instance, could service up to 65 users. In order to increase its capacity, the IP could purchase more lines to increase its bandwidth. This would, however, also require equipment which could handle the added load. The load on the IP is limited by its own resources and by the market.

The architecture of the VDT network has already been developed by engineers at Bell Communications Research (Bellcore). In order to implement Net-CD, software drivers and VDT interface hardware must designed for the CPE. The following section describes the design of the required software for the customer's PC.

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3.2 Net-CD Software Requirements

The design of the necessary software to implement Net-CD required research into MS-DOS's implementation of CD-ROM drives in general. Unfortunately, DOS cannot handle CD-ROM drives in the same manner as it does for regular disk drives, due to the large size of data files that is allowed on CD-ROM drives.

DOS requires CD-ROM drives to be supplied with a software driver which will allow communication between the drive and Microsoft's CD-ROM Extensions (MSCDEX). MSCDEX masks the workings of the CD-ROM drive from DOS in a manner that makes the CD appear as if it were a network drive. This is done because DOS allows an exception for network drives in terms of maximum file size.

Figure 3.2 shows the links between DOS, MSCDEX, DRIVER.SYS (the CD-ROM drive specific device driver), and the CD-ROM drive. Both of the first two sections, DOS and MSCDEX, are independent of the type of CD-ROM drive in the system. The final two sections are device dependent and therefore cannot be implemented in the PC if the CPE is to be made as general as possible. This division into device independent and dependant sections is illustrated by a vertical line in Figure 3.2.

The feasibility of implementing the Net-CD on a PC can be shown by writing and testing software which allows us to replace the vertical line in Figure 3.2 with a division in hardware, such as the ADSL line in the VDT architecture. This is shown in Figure 3.3.

In this configuration, MSCDEX communicates with LOC-CD.SYS, which, to MSC-DEX, looks like a CD-ROM device driver. LOC-CD.SYS is linked to REM-CD.EXE in the remote system via a communications channel. REM-CD.EXE, on the other hand, looks like MSCDEX to the device dependent DRIVER.SYS.

The communications channel would be replaced by the telephone network between the CPE and the server in the actual implementation of Net-CD using VDT. However, for this experiment, it is a 115.2kbit/sec serial channel. Although this is roughly 12% of the speed of system limiting ADSL speed of 1.522 Megabits/sec, it will not effect the feasibility of a PC Net-CD, as long as a robust CPE interface to the ADSL line can be created.



Figure 3.1: VDT Implementation of a Networked PC CD-ROM

The LOC-CD.SYS software driver will receive commands from MSCDEX, and group them into one of two categories. The first consists of all commands which do not require any interaction with the CD-ROM device. All requests of this type will be answered locally. The remaining commands fall into the second category. These requests are placed into a packet for shipment to the remote PC's REM-CD.EXE.



Figure 3.2: PC CD-ROM Software Interaction



Figure 3.3: Net-CD Software Interaction

The REM-CD.EXE receives commands from the LOC-CD.SYS and passes them to DRIVER.SYS as if they were coming directly from MSCDEX. When the driver returns the request, REM-CD.EXE sends the result of the request, which may contain large amounts of data, back to LOC-CD.SYS. At this point, LOC-CD.SYS returns the result to MSCDEX as if the data had come from a device dependant driver. In this way, the DOS in the local PC will accept LOC-CD.SYS as a valid device driver for an internal CD-ROM drive.

Chapter 4

Design of Net-CD Software

4.1 Operation of MSCDEX

Microsoft's CD Extensions (MSCDEX) require the presence of a CD-ROM device driver. This driver must be able to properly handle CD-ROM requests, while still meeting the requirements of a DOS device driver. Except for the initialization stage, all requests of the device driver are made through MSCDEX. These requests are in the DOS standard form of the *request header*. This request header varies in length dependant on the type of request made. The type of request is specified in the *command* field of the request header. The commands that are supported by LOC-CD.SYS are listed in Table 4.1. All commands with an index of 128 or greater are CD-ROM device driver specific commands.

 Table 4.1: Driver Commands Supported by Net-CD

Index	Command	
0	Init	
3	IO Control Input	
7	IO Control Flush	
12	IO Control Output	
13	Device Open	
14	Device Close	
128	Read Long	
130	Read Long Prefetch	
131	Seek	

Normal operation of a CD-ROM drive consists of DOS making calls to MSCDEX which, in turn, makes requests of the CD-ROM device driver. However, DOS allows software to make direct requests of the device driver via its *multiplex interrupt* (INT 0x2F). These requests have a form and function identical to the request header. The multiplex interrupt will be shown to play a key role in the operation of REM-CD.EXE.

4.2 Design of LOC-CD.SYS

In order to implement LOC-CD.SYS, the source code for this driver must conform to DOS device driver requirements. Every DOS device driver must begin with the device header. This 13 byte header is used to inform DOS of the characteristics of the device, as well as pointers to two required procedures: *strategy* and *interrupt*. When DOS makes requests of the device, it first calls the strategy procedure, then follows with a call to the interrupt procedure.

Beyond the DOS device driver requirements, LOC-CD.SYS must be able to handle the extended CD-ROM device driver commands of MSCDEX. These commands, along with all supported regular commands except for initialization, must be sent to the remote CD-ROM site via the serial communications channel. This requires serial communication routines and a protocol for executing the command remotely. The baud rate for the serial line for this application is set to 115.2 kilobits per second, which is the maximum rate that DOS allows.

4.3 Communication Protocol

Table 4.2 depicts the protocol used for the Net-CD's data transmission. The local system begins the communication with a wakeup signal. After this signal is received and acknowledged by REM-CD.EXE, the system then begins a loop which begins when the local system's software receives a CD-ROM request from MSCDEX. The request is then sent to the remote system. The remote system then performs the MSCDEX request. The results of the request are then transmitted back to LOC-CD.SYS which, in turn, gives the results to MSCDEX. When an error occurs, the process detecting the error will transmit the error byte until the offending system returns the error acknowledge signal. At this point, the protocol returns to the original transmission of the request header length (denoted by a "*" in Table 4.2).

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Local System	Remote System
Send Wakeup	
	Receive Wakeup
	Send Acknowledge(ACK)
Receive ACK	
Get Request from MSCDEX	
Send Request Header (RH) Length	
	Receive RH Length
	Send ACK
Receive ACK	
Send RH Data Packet (RHP)	
	Receive RHP
	Send ACK
Receive ACK	
	Perform Request
	Send RH Result Length
Receive RH Result Length	
Send ACK	
	Receive ACK
	Send RH Result Data Packet
Receive RH Result Data Packet	
Send ACK	
	Receive ACK
	Send CD Data Length
Receive CD Data Length	
Send ACK	
	Receive ACK
	Send Data Packet
Receive Data Packet	
Send ACK	
	Receive ACK
Process Data	
Return to *	Return to *

Table 4.2: Protocol for Net-CD Data Transmission

* **

4.4 Design of REM-CD.EXE

REM-CD.EXE must be able to receive data via the serial line and execute commands sent from LOC-CD.SYS via the given protocol. Since the requests of MSCDEX to LOC-CD.SYS that are received by REM-CD.EXE are in a form standard to CD-ROM device driver requests, REM-CD.EXE can take advantage of the multiplex interrupt mentioned earlier to call the actual CD-ROM device driver, DRIVER.SYS, with ease.

Once DRIVER.SYS has processed the request, REM-CD.EXE can then send the results back to LOC-CD.SYS via the serial line. REM-CD.EXE is simply an interface between LOC-CD.SYS and the DRIVER.SYS. All data associated with each request returned according to the protocol, and is unchanged by REM-CD.EXE.

Chapter 5

Results

The implementation of Net-CD described in the previous chapters is completed and functional. The project was a success. Except for the slow CD-ROM data transfer delay, the CPE was indistinguishable from a PC containing a CD-ROM drive. Several CD-ROM applications were tested, both from DOS and from the Microsoft Windows environment. Files and directories were tested for integrity via a comparison between the information obtained by accessing the CD-ROM drive both locally and remotely.

The Net-CD data transfer rate on the serial line is 13.2 times slower than would be possible by using an ADSL line instead of a serial one. Assuming that a normal CD-ROM drive has a data transfer bandwidth of 300kbits/sec, the transfer of one data sector on a CD-ROM (2048 bytes) will take 54.6 milliseconds on a PC with a local CD-ROM drive. One sector of data will take 10.8 milliseconds to travel on an ADSL line, yielding a 19.8% increase in total data access time for one sector on the Net-CD system. The system implemented in this paper, however, has a data sector transfer time of 142 ms -- a 260% increase in data access time.

The 19.8% increase in data access time is not insignificant. This can be reduced, however, by copying the information on a CD-ROM to a large high-speed magnetic disk which could actually make the Net-CD faster to use than a local CD-ROM drive, and would allow multiple users to simultaneously access a single copy of the application.

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Chapter 6

Conclusion

The networking of PC CD-ROM drives is one of many possible types of multimedia systems. The success of the Net-CD system described in this paper shows that PC CD-ROM networking is feasible. The implementation of a Net-CD, perhaps over a VDT type system, will give multimedia users access to a large library of multimedia applications and will help open the market for Video Dial Tone and other interactive multimedia systems

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Appendix A

LOC-CD.ASM

.286 COM1 = 03F8HCOM2 = 02F8HCOM3 = 03E8HCOM4 = 02E8HCOM = COM1; RETURNS ZERO FLAG TRUE (JZ WILL BE TAKEN) IF OK TO XMIT X_OK MACRO MOV DX, COM+5 IN AL, DX AND AL,020H CMP AL,020H ENDM ; RETURNS ZERO CONDITION TRUE (JZ BRANCHES) IF OK TO READ REC. REGISTER R_OK MACRO MOV DX, COM+5 IN AL, DX AND AL,01H CMP AL,01H ENDM X_RDY MACRO ; MOV AL, '*' ; CALL P_AL MOV DX, COM MOV AL, RCV_READY OUT DX,AL ENDM

R_RDY MACRO MOV DX,COM IN AL,DX CALL P_BYTE CMP AL,RCV_READY ENDM

RECV_ALM MACRO MOV DX,COM IN AL,DX ENDM

R_WAITM MACRO @@: R_OK JNE @B ENDM

_TEXT SEGMENT BYTE PUBLIC `CODE' ASSUME CS:_TEXT

; ; DEVICE DRIVERS ORIGINATE AT 0 (NOT 100H) ;

ORG 0

VERY_TOP LABEL BYTE
;
; ;; FIRST CHARACTER DEVICE HEADER (MUST BE AT OFFSET 0)

DEVICEHEADER:

;

DW -1, -1 ; POINTER TO NEXT DEVICE DW 0C800H ; (CHARACTER, IOCTL, OPEN/CLOSE/RM) DW STRATEGY ;STRATEGY ROUTINE OFFSET DW INTERRUPT ;INTERRUPT ROUTINE OFFSET DB `MSCD000 ` ;DEVICE NAME (MUST BE 8 VALID CHARS) DW 0 ;RESERVED DB 0 ;DRIVE LETTER UNITS DB 1 ;NUMBER OF UNITS

REQUESTHEADER:

DD 0 ; STRATEGY KEEPS REQUEST HEADER PTR HERE

; REQUEST HEADER LENGTH TABLE RHLT: DB 23 ; ; 0-INIT DB 13 ; ; 1-ERROR DB 13 ; ; 2-ERROR DB 26 ; ; 3-IOCTL INPUT DB 13 ; ; 4-ERROR DB 13 ; ; 5-ERROR DB 13 ; ; 6-ERROR DB 13 ; ; 7-IOCTL FLUSH DB 13 ; ; 8-ERROR DB 13 ; ; 9-ERROR DB 13 ; ;10-ERROR DB 13 ; ;11-ERROR DB 26 ; ;12-IOCTL OUTPUT DB 13 ; ;13-DEVICE OPEN DB 13 ; ;14-DEVICE CLOSE DB 27 ; ;128-READ LONG DB 13 ; ;129-ERROR DB 27 ; ;130-READ LONG PREFETCH DB 24 ; ;131 SEEK DB 22 ; ;132 PLAY AUDIO DB 13 ; ;133-STOP AUDIO DB 13 ; ;134-ERROR DB 13 ; ;135-ERROR DB 13 ; ;136-RESUME AUDIO

PROCESSING_MSG DB `PROCESSING MSCDEX REQUEST.',0DH,0AH,'\$'
SENDING_MSG DB `SENDING PACKET...','\$'
SENDING_MSG2 DB `SENDING RCV_READY','\$'
SENT_MSG DB `PACKET SENT',0DH,0AH,'\$'
RECEIVED_MSG DB `RCVED.',0DH,0AH,'\$'
RCVING_MSG DB `RCVING.','\$'

```
SENT_BYTE DB 'BYTE SENT...','$'
RECEIVED_BYTE DB 'BYTE RECEIVED...','$'
OUT_DATA_MSG DB `(OUTDATA)','$'
DONE_MSG DB 'DONE WITH REQUEST', ODH, OAH, '$'
IOCTL_MSG DB 'IOCTL INPUT REQUEST #','$'
GETHDR DB 'GET HEADER ADDRESS', 0DH, 0AH, '$'
OTHERMSG DB 'OTHER OPERATION', 0DH, 0AH, '$'
EXITMSG DB 'EXITING OPERATION', ODH, OAH, '$'
HEXES DB '0123456789ABCDEF'
INTMSG DB 'INTERRUPT', ODH, OAH, '$'
STRMSG DB `STRATEGY', ODH, OAH, '$'
DBG_MSG DB `ERROR! OFFSET!=0','$'
OK_MSG DB "LOC-CD READY.", 0AH, 0DH, '$'
COMMAND DB 8
DATA LEN DW ?
DTA_OFS DW ?
DTA SEG DW ?
DOS_DTA_OFS DW ?
DOS DTA SEG DW ?
TMP COUNT DW 3
=================
.
; STRATEGY
;
; DEVICE STRATEGY ROUTINE. SAVES THE POINTER TO THE CALLER'S
   REOUEST
; HEADER STRUCTURE FOR THE SUBSEQUENT CALL INTO INTERRUPT
   BELOW.
;
; ENTRY:
; ES:BX - FAR PTR TO THE CALLER'S REQUEST HEADER.
;
; EXIT:
; REQUESTHEADER - FAR PTR TO THE CALLER'S REQUEST HEADER.
```

```
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```

```
_____
PR_RH PROC NEAR
PUSH AX
PUSH BX
PRH_LOOP:
MOV AX, BX
CALL P_BYTE
MOV AL, ':'
CALL PRNT_AL
MOV AL, BYTE PTR DS: [BX]
CALL P_BYTE
INC BX
MOV AL, ' '
CALL PRNT_AL
CALL PRNT_AL
LOOP PRH_LOOP
POP BX
POP AX
RET
PR_RH ENDP
SAVE_XFER PROC NEAR
```

PUSH AX MOV AX,WORD PTR DS:[BX+14] MOV CS:DOS_DTA_OFS,AX CMP AX,0 JE @F MOV DI,OFFSET CS:DBG_MSG CALL PR_MSG @@: MOV AX,WORD PTR DS:[BX+16] MOV CS:DOS_DTA_SEG,AX POP AX RET SAVE_XFER ENDP RESTORE_XFER PROC NEAR PUSH AX MOV AX, WORD PTR CS:DOS_DTA_OFS MOV CS:DTA_OFS,AX MOV WORD PTR DS: [BX+14], AX MOV AX, CS: DOS_DTA_SEG MOV CS:DTA_SEG,AX MOV WORD PTR DS:[BX+16],AX POP AX RET RESTORE_XFER ENDP PRINT_DATA PROC NEAR PUSH CX PUSH BX PUSH AX PUSH ES MOV AX, CS: DOS_DTA_SEG MOV ES, AX MOV BX, CS: DOS_DTA_OFS MOV CX,256 PDATA_LOOP: MOV AL, BYTE PTR ES: [BX] CALL P_BYTE MOV AL, ' ` CALL PRNT_AL CALL PRNT_AL INC BX LOOP PDATA_LOOP POP ES POP AX POP BX POP CX RET PRINT_DATA ENDP OUT_DATA PROC NEAR PUSH BX PUSH ES

PUSH CX PUSH AX ; PUSH DI ; MOV DI, OFFSET CS: OUT_DATA_MSG ; CALL PR_MSG MOV CX, 128 MOV ES, CS: DOS_DTA_SEG MOV BX, CS: DOS_DTA_OFS MOV AH, 0 **@@**: MOV AL, BYTE PTR ES: [BX] CALL P_BYTE INC BX LOOP @B POP DI POP AX POP CX POP ES POP BX RET OUT_DATA ENDP ; ; ON ENTRY: ; DS:BX POINTS TO REQUEST HEADER SEND_REQUEST_HEADER PROC NEAR MOV CS:COMMAND, RHP PUSH DI PUSH AX MOV AL, DS:RQH_CMD[BX] CMP AL,14

CMP AL,14 JBE @F SUB AL,113 @@: MOV AH,0

MOV DI, AX MOV AL, BYTE PTR CS:RHLT[DI] MOV AL, BYTE PTR DS: [BX] MOV CS:DATA_LEN,AX MOV CS:DTA_OFS, BX MOV CS:DTA_SEG,DS CALL SND_PKT CALL SAVE XFER MOV AL, DS: [BX].RQH_CMD ; MOV CX, CS: DATA_LEN ; CALL PR_RH CMP AL,3 JNE @F ; IOCTL_I LES DI, IOCTL_XFER[BX] ; APPLICATION'S TRANSFER ADDRESS MOV CS:DTA_OFS, DI MOV CS:DTA_SEG,ES MOV CS:DATA_LEN,1 CALL SND_PKT @@: CMP AL,12 JNE @F LES DI, IOCTL_XFER[BX] ; APPLICATION'S TRANSFER ADDRESS MOV CS:DTA_OFS, DI ; IOCTL_O MOV CS:DTA_SEG,ES MOV CS:DATA_LEN,1 CALL SND_PKT **@@**: POP AX POP DI RET SEND_REQUEST_HEADER ENDP

RECEIVE_REQUEST_HEADER PROC NEAR PUSH AX PUSH DI MOV CS:DTA_OFS,BX MOV CS:DTA_SEG,DS

CALL RCV_PKT CALL RESTORE_XFER CALL RCV_PKT MOV AX, CS: DATA_LEN CMP AX,256 JB @F CALL PRINT_DATA 66: POP DI POP AX RET RECEIVE_REQUEST_HEADER ENDP STRATEGY PROC FAR MOV WORD PTR CS:REQUESTHEADER.LO, BX MOV WORD PTR CS:REQUESTHEADER.HI, ES MOV DI, OFFSET CS: STRMSG CALL PR_MSG RET STRATEGY ENDP

;=========; ; ; INTERRUPT ; ; MAIN ENTRY POINT TO THE DEVICE INTERRUPT HANDLER. ; ; ENTRY: ; REQUESTHEADER - FAR PTR TO THE CALLER'S REQUEST HEADER. ; (SEE MSCDEX.INC) FOR FORMAT OF THE REQUEST HEADER) ; ;============

INTERRUPT PROC FAR

```
PUSH AX
    PUSH BX
    PUSH CX
    PUSH DX
    PUSH SI
    PUSH DI
    PUSH BP
    PUSH DS
    PUSH ES
MOV DI, OFFSET CS: INTMSG
 CALL PR_MSG
CALL SET_BD
;
; POINT DS:BX AT THE REQUEST HEADER
;
    LDS BX, DWORD PTR CS:REQUESTHEADER
    MOV AL, DS: [BX].RQH_CMD ;GET COMMAND
    CMP AL, DVRQ_NCMD_MAX ; EXTENDED COMMAND?
    JBE OK_CMD ; NOPE
;
; MUST BE EXTENDED CDROM DEVICE DRIVER COMMAND, CONVERT INTO
    INDEX AT
; END THE NORMAL CDROM DEVICE DRIVER FUNCTION TABLE
;
    CMP AL, DVRQ_ECMD_MIN
    JB @F
    SUB AL, (DVRQ_ECMD_MIN - DVRQ_NCMD_MAX - 1)
    CMP AL, (DVRQ_NCMD_MAX + DVRQ_ECMD_MAX - DVRQ_ECMD_MIN +
    1)
    JBE OK_CMD ; YES
@@:
    JMP ERROR ; NOPE, COMMAND OUT OF RANGE
OK CMD:
    CBW
    OR AX, AX ; INIT TIME?
    JZ INIT ; YES
    ; NO, DISPATCH TO COMMAND
    CMP AX,3
```

```
JE IOCTL_I
```

SEND:

MOV DI, OFFSET CS: PROCESSING_MSG CALL PR_MSG CALL SEND_REQUEST_HEADER CALL RECEIVE_REQUEST_HEADER

JMP EXIT_INIT

```
===========
;
; INIT
;
; INIT, CD-ROM DEVICE DRIVER ROUTINE TO INITIALIZE THE DRIVER.
;
; THIS IS THE ONLY DEVICE DRIVER CALL COMING DIRECTLY FROM
   DOS, AND IS
; ONLY MADE ONCE. INIT SHOULD INITIALIZE CERTAIN FIELDS IN THE
   REQUEST
; HEADER:
; - SET INIT_UNITS AND INIT_DEVNO TO 0 SINCE DOS VIEWS THIS
   AS A CHARACTER
; DEVICE (MSCDEX MAKES ITS OWN DETERMINATION OF THE NUMBER OF
   UNITS THROUGH
; THE DEVICE HEADER).
; - RETURN THE ADDRESS OF THE END OF THE RESIDENT CODE/DATA
   SECTION IN
; INIT_ENDADDR. CODE/DATA AFTER THIS POINTER IS DISCARDED.
; - PARSE THE CONFIG.SYS LINE AFTER THE `=' CHARACTER (POINTED
   TO BY
; INIT_BPBARR) FOR THE DEVICE NAME AND FILL IN THE DEV_NAME
   FIELD IN
; THE DEVICE HEADER MAKING SURE IT IS A LEGAL 8-CHARACTER
   FILENAME
; (PADDED OUT TO 8 CHARACTERS WITH SPACES IF NECESSARY).
ï
```

INIT:

MOV DS:[BX].INIT_UNITS,0 MOV DS:[BX].INIT_DEVNO,0 MOV WORD PTR DS:[BX+2].INIT_ENDADDR,CS MOV WORD PTR DS:[BX].INIT_ENDADDR,OFFSET VERY_END JMP EXIT_INIT

;_IOCTLHEADER STRUC ;IOCTL_RQH DB SIZE _REQUESTHEADER DUP (?) ;IOCTL_MEDIA DB ? ;IOCTL_XFER DD ? ;IOCTL_NBYTES DW ? ;IOCTL_SECTOR DW ? ;IOCTL_VALID DD ? ;_IOCTLHEADER ENDS

IOCTL_I:

LES DI, IOCTL_XFER[BX] ;APPLICATION'S TRANSFER ADDRESS MOV AL, BYTE PTR ES:[DI]

MOV AX,0

CMP AX,0 JNE SEND INC DI

```
MOV WORD PTR ES: [DI], OFFSET CS: DEVICEHEADER
MOV WORD PTR ES: [DI+2], CS
JMP EXIT_NOT_BUSY
```

```
=
;
; IOCTL INPUT COMMANDS...
;
=
;
; RETURN_ADDRESS
;
; IOCTL INPUT SUB-FUNCTION #0.
;
; RETURN THE ADDRESS OF THE DEVICE HEADER.
;
_____
```

RETURN_ADDRESS:

MOV WORD PTR ES: [DI].IO_DEVADDR,OFFSET DEVICEHEADER MOV WORD PTR ES: [DI+2].IO_DEVADDR,CS ; JMP SHORT IOCTL_IO_EXIT

```
; =========;
;
; DEVICE_OPEN
;
; DEVICE OPEN, CD-ROM DEVICE DRIVER ROUTINE INDICATING TO THE
DEVICE
```

```
; DRIVER THAT AN APPLICATION IS BEGINNING TO USE IT.
;
_____
_______
ï
; EXIT POINTS.
;
ERROR:
;
; EXIT INDICATING UNKNOWN COMMAND ERROR
;
  MOV AX, (ERRBIT + DONEBIT + DRVERR_UNKNOWN_COMMAND)
  JMP SHORT EXIT
EXIT_BUSY:
;
; EXIT INDICATING NO ERRORS AND AUDIO IS PLAYING
;
  MOV AX, (DONEBIT + BUSYBIT)
  JMP SHORT EXIT
EXIT_NOT_BUSY:
;
; EXIT INDICATING NO ERRORS AND AUDIO IS NOT PLAYING
;
  MOV AX, DONEBIT
  LDS BX, DWORD PTR CS: REQUESTHEADER
  MOV [BX].RQH_STATUS,AX
```

EXIT:

;

; EXIT AND POKE THE STATUS WORD INTO THE REQUEST HEADER. ; THIS IS ALWAYS THE FINAL EXIT POINT. ; ; LDS BX,DWORD PTR CS:REQUESTHEADER ; MOV [BX].RQH_STATUS,AX

EXIT_INIT: ; ; EXIT POINT FROM INIT ; MOV DI, OFFSET CS: DONE_MSG CALL PR_MSG POP ES POP DS POP BP POP DI POP SI POP DX POP CX POP BX POP AX RET INTERRUPT ENDP DELAY PROC NEAR PUSH CX MOV CX, 0FFFFH 66: JMP SHORT \$+2 LOOP @B POP CX

RET

DELAY ENDP

; FUNCTIONS IN SERIAL.ASM

CONSOLE

SET_BD PROC NEAR

MOV DX, COM+3 OUT DX,AL

JMP SHORT \$+2

JMP SHORT \$+2

MOV DX, COM+1 OUT DX,AL

JMP SHORT \$+2

MOV DX, COM+3 OUT DX,AL

JMP SHORT \$+2

RET

MOV DX, COM OUT DX,AL

```
; MACRO X_OK, R_OK => XMIT OR RCV OK TEST. JZ WILL BRANCH IF OK
; PROC SET_BD => SET BAUD RATE TO 1200 BAUD, 8N1
; PROC XMIT_AL => PLACE [AL] IN XMIT REGISTER
```

; PROC PRNT_AL => PRINTS THE ASCII REPRESENTATION OF AL TO

; PROC R_WAIT => LOOP UNTIL OK TO READ RCV REGISTER

MOV AL,080H ; SET DLAB IN REGISTER 3 OF COM1

MOV AL,01H ; SET BAUD LSB TO 1 (REG. 0)

MOV AL, 0 ; SET BAUD MSB TO 0 (REG. 1)

MOV AL,03H ; SET DLAB=0, REG. 3 = 8-N-1

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; PROC X_WAIT => LOOP UNTIL OK TO XMIT

; PROC RECV_AL => READ RCV REGISTER TO AL

; PROC R_FLUSH => FLUSHES RCV BUFFERS

SET_BD ENDP

XMIT_AL PROC NEAR

MOV DX,COM OUT DX,AL ; CALL P_BYTE RET

XMIT_AL ENDP

X_WAIT PROC NEAR

PUSH AX

TBE:

X_OK JNE TBE POP AX RET X_WAIT ENDP

RECV_AL PROC NEAR

MOV DX,COM IN AL,DX RET

RECV_AL ENDP P_AL PROC NEAR PUSH AX PUSH BX MOV AH, 0EH

MOV BH,0 ;INT 10H POP BX POP AX RET P_AL ENDP PRNT_AL PROC NEAR ; PUSH AX ; PUSH BX ; MOV AH, OEH ; MOV BH,0 ; INT 10H ; POP BX ; POP AX RET PRNT_AL ENDP ; PRINT MESSAGE PROCEDURE ; ON ENTRY: ; DI POINTS TO MESSAGE IN THE DATA ; SEGMENT WHICH IS TERMINATED BY A \$ PR_MSG PROC NEAR PUSH AX PUSH BX MSGLOOP: MOV AL, CS: [DI] CMP AL,'\$' JE MSGDONE PUSH DI CALL P_AL POP DI INC DI JMP MSGLOOP MSGDONE: POP BX POP AX RET PR_MSG ENDP

; P_WORD PRINTS THE VALUE IN AX AS A HEXADECIMAL P_WORD PROC NEAR PUSH AX CALL P_BYTE SHR AX,8 CALL P_BYTE POP AX RET P_WORD ENDP ; P_BYTE PRINTS THE VALUE IN AL AS A HEXADECIMAL P_BYTE PROC NEAR PUSH BX PUSH AX MOV BX, AX SHR BX,4 AND BX,000FH MOV AL, CS: HEXES [BX] CALL P_AL POP BX PUSH BX AND BX,000FH MOV AL, CS: HEXES [BX] CALL P_AL POP AX POP BX RET P_BYTE ENDP

```
R_WAIT PROC NEAR
; MOV AL, '+'
; CALL P_AL
RXRDY:
   R_OK
    JNE RXRDY
    RET
R_WAIT ENDP
R_FLUSH PROC NEAR
FLUSHER:
   R_OK
    JNE FL_DONE
    CALL RECV_AL
 JMP FLUSHER
FL_DONE:
    RET
R_FLUSH ENDP
ONE_FLUSH PROC NEAR
PUSH CX
MOV CX,10
@@:
R_OK
 JNE @F
CALL RECV_AL
LOOP @B
@@:
 POP CX
    \operatorname{RET}
ONE_FLUSH ENDP
RCV_PKT PROC NEAR
```

PUSH ES PUSH DI PUSH AX PUSH BX PUSH CX CALL R_FLUSH MOV DI, OFFSET CS: SENDING_MSG2 CALL PR_MSG MOV AX, CS: DTA_OFS SHR AX,4 ADD AX, CS: DTA_SEG MOV CS:DTA_SEG,AX MOV AX, CS: DTA_OFS AND AX,000FH MOV CS:DTA_OFS,AX MOV AX, CS: DTA_SEG MOV ES, AX MOV AX, CS: DTA_OFS MOV DI,AX CALL R_FLUSH RCV_GO: MOV AL, '-'CALL P_AL SND_CHK: CALL X_WAIT X RDY MOV CX,65535 @@: R_OK JE RCV LOOP @B R_OK JNE SND_CHK RCV:

RECV_ALM

CMP AL,CS:COMMAND JNE RCV_GO CLI R_WAITM RECV_ALM MOV BYTE PTR CS:DATA_LEN[0],AL R_WAITM RECV_ALM MOV BYTE PTR CS:DATA_LEN[1],AL

MOV CX,CS:DATA_LEN CMP CX,0 JE RCV_DONE

RCV_LOOP:

R_WAITM RECV_ALM MOV BYTE PTR ES:[DI],AL INC DI LOOP RCV_LOOP

RCV_DONE: STIMOV DI, OFFSET CS:RECEIVED_MSG CALL PR_MSG POP CX POP BX POP AX POP DI POP ES RET RCV_PKT ENDP SND_PKT PROC NEAR PUSH ES PUSH DI PUSH AX PUSH BX PUSH CX

CALL DELAY MOV AX, CS: DTA_OFS SHR AX,4 ADD AX, CS: DTA_SEG MOV CS:DTA_SEG,AX MOV AX, CS: DTA_OFS AND AX,000FH MOV CS:DTA_OFS,AX MOV DI, OFFSET CS: SENDING_MSG CALL PR_MSG @@: CALL R_WAIT R RDY JNE @B MOV DI, OFFSET CS: OK_MSG CALL PR_MSG MOV AL, CS: COMMAND CALL P_BYTE CALL DELAY CALL X_WAIT CALL XMIT_AL MOV AL, BYTE PTR CS:DATA_LEN[0] CALL DELAY CALL X_WAIT CALL XMIT_AL MOV AL, BYTE PTR CS: DATA_LEN[1] CALL DELAY CALL X_WAIT CALL XMIT_AL MOV AX, CS: DATA_LEN CALL P_WORD CALL ONE_FLUSH MOV CX, CS: DATA_LEN CMP CX,0 JE SND DONE MOV ES, CS: DTA_SEG MOV DI, CS: DTA_OFS SND LOOP: CALL DELAY CALL X WAIT MOV AL, BYTE PTR ES: [DI]

CALL XMIT_AL INC DI LOOP SND_LOOP

SND_DONE:

MOV DI, OFFSET CS:SENT_MSG

CALL PR_MSG

 POP
 CX

 POP
 BX

 POP
 AX

POP DI

POP ES

RET

SND_PKT ENDP

VERY_END:

_TEXT ENDS

END

Appendix B

REM-CD.ASM

```
; Rem-CD.asm
  ; .
  ; Remote executable to send data
  ; and interpret requests for the
  ; Net-CD system.
  ĩ
  ; Steve Levis
  ; 12/19/94
  .286
  .MODEL LARGE
  COM1=03F8h
  COM2 = 02F8h
  COM3 = 03E8h
  COM4=02E8h
  COM=COM4
  INCLUDE CMDS.ASM
  ; Returns Zero Flag True (JZ will be taken) if OK to xmit
  X_OKMACRO
    MOVDX, COM+5
     INAL, DX
     ANDAL,020h
     CMPAL,020h
  ENDM
  ; Returns Zero Condition True (JZ Branches) if OK to read
rec. register
  R_OKMACRO
     MOVDX, COM+5
```

```
INAL, DX
  ANDAL,01h
  CMPAL,01h
  ENDM
X_RDY MACRO
; MOV AL, '*'
; CALL P_AL
MOV DX,COM
 MOV AL, Rcv_Ready
 OUT DX,AL
; CALL P_BYTE
 ENDM
R_RDY MACRO
 MOV DX, COM
 IN AL, DX
 CMP AL, Rcv_Ready
 ENDM
RECV_ALM MACRO
 MOV DX, COM
  IN AL, DX
  ENDM
 R_WAITM MACRO
 ; PUSH BX
 ; MOV CX, 0FFFFh
 ; MOV BX,0
 @@: R_OK
 JNE @B
 ; JE @F
 ; LOOP @B
 ; MOV BX,1
 ;@@:
 ;
 ; CMP BX,0
 ; POP BX
  ENDM
```

```
; Functions in SERIAL.ASM
  ; MACRO X_OK,R_OK => Xmit or Rcv OK test.JZ will branch
if ok
  ;
  ; PROCSET BD => Set Baud Rate to 1200 Baud, 8N1
  ; PROCXMIT_AL => Place [AL] in Xmit Register
  ; PROCX_WAIT => Loop until OK to Xmit
  ; PROCRECV_AL => Read Rcv Register to AL
  ; PROCR_WAIT => Loop until OK to read Rcv Register
  ; PROCPRNT_AL => Prints the ASCII representation of AL to
console
  ; PROCR FLUSH => Flushes Rcv Buffers
  SSEGSEGMENT STACK
     DB32 DUP("STACK---")
  SSEGENDS
  DSEGSEGMENT
  OK_MSGDB"Ok.",'$'
  command db 8
  data len dw 31
  DTA_ofs dw ?
  DTA_seg dw ?
  dbg_msg db "Non-Cooked sector requested!",'$'
  RH_MSG db "Request Header Command:",'$'
  Special_MSG db "No special packet required", 0ah, 0dh, '$'
  sent_msg db "Packet sent.",0ah,0dh,'$'
  sending_msg db "Sending packet...",'$'
  received_msg db "Packet received.", 0ah, 0dh, '$'
  rcving_msg db "Receiving packet...",'$'
  waiting_msg db "Waiting for packet...",'$'
  received_byte db "Byte received...",'$'
  sent_byte db "Byte sent...",'$'
  cmd_msg db "Cmd Rec.",'$'
  hexes db '0123456789ABCDEF'
  unx_msg db 'Unexpected command:','$'
  s_req_msg db `Sending Request', 0ah, 0dh, '$'
  s_rh_msg db 'Sending Request Header', 0ah, 0dh, '$'
  delay_msg db` $'
```

```
s_data_msg db `Sending Data',0ah,0dh,'$'
inting_msg db `Calling Interrupt 2F','$'
newline db 0ah,0dh,'$'
ReqHdr db 2 Dup("Request Header__")
RH_Len dw ?
DT_Len dw ?
DSEGENDS
```

```
BSEG1 SEGMENT
ORG 0
Buffer1 db 8192 DUP("BUFFER___")
BSEG1 ENDS
```

CSEGSEGMENT `CODE' ASSUMECS:CSEG,DS:DSEG,SS:SSEG

MAINPROCFAR

PUSHDS PUSH0 MOVAX,DSEG MOVDS,AX

CALLSET_BD CALLR_FLUSH MOV DI,OFFSET OK_MSG CALL PR_MSG

PROC_CMD: MOV COMMAND, RHP CALL RCV_RH CMP DT_LEN,63488 JA ENDMAIN MOV DI,OFFSET DS:S_REQ_MSG CALL PR_MSG CALL SND_REQ MOV DI,OFFSET DS:S_RH_MSG

```
CALL PR_MSG
CALL SND_RH
MOV DI, OFFSET DS:S_DATA_MSG
CALL PR_MSG
CALL SND_DATA
JMP PROC_CMD
ENDMAIN:
MOV AX, DT_LEN
CALL P_WORD
RET
MAINENDP
delay proc near
push cx
push di
mov di, offset ds:delay_msg
call pr_msg
mov cx, 0FFFFh
@@:
 jmp short $+2
loop @b
pop di
pop cx
 ret
delay endp
RCV_RH PROC NEAR
 MOV DTA_ofs, OFFSET ReqHdr
MOV DTA_seg,DS
 CALL RCV_PKT
 MOV AX, DS: Data_Len
 MOV DS:RH_Len,AX
 MOV AL, DS: ReqHdr[2] ; Command Code Field
 CMP AL,3
 JNE @f
 ; IOCTL_I
```

```
MOV WORD PTR DS:ReqHdr[14], OFFSET BSEG1:Buffer1
MOV WORD PTR DS:ReqHdr[16], BSEG1
MOV AX, WORD PTR DS:RegHdr[18]
MOV DT_LEN, AX
MOV DTA_ofs, OFFSET BSEG1:Buffer1
MOV DTA_seg, BSEG1
 CALL RCV_PKT
 JMP RH DONE
@@:
CMP AL,12
 JNE @f
 ; IOCTL_O
MOV WORD PTR DS:ReqHdr[14], OFFSET BSEG1:Buffer1
MOV WORD PTR DS:ReqHdr[16], BSEG1
MOV AX, WORD PTR DS:ReqHdr[18]
MOV DT_LEN, AX
MOV DTA_ofs, OFFSET BSEG1:Buffer1
MOV DTA_seg, BSEG1
CALL RCV_PKT
JMP RH DONE
66:
```

CMP AL,128 JNE @f

```
READ_LONG: ; READ LONG
MOV WORD PTR DS:ReqHdr[14],OFFSET BSEG1:Buffer1
MOV WORD PTR DS:ReqHdr[16],BSEG1
MOV AX,WORD PTR DS:ReqHdr[18]
SHL AX,11 ; AX = AX * 2048
MOV DT_LEN,AX
MOV AL,BYTE PTR DS:ReqHdr[24]
CMP AL,0
JE RLDONE
MOV DI,OFFSET DS:DBG_MSG
CALL PR_MSG
STAY: JMP STAY
```

RLDONE: MOV DTA_ofs,OFFSET BSEG1:Buffer1

```
MOV DTA_seg, BSEG1
 JMP RH_DONE
00:
MOV DS:DT_LEN,0
RH_DONE:
MOV AL, DS:ReqHdr[2]
CALL P_BYTE
MOV AL, ' '
CALL P_AL
RET
RCV_RH ENDP
SND_REQ PROC NEAR
MOV CX,50
@@: PUSHCX
MOV AL, '#'
CALL P_AL
MOV AX, 1510h
MOV CX,4 ; Drive E:
MOV BX, DS
MOV ES, BX
MOV bx, OFFSET BSEG1:Buffer1
MOV word ptr DS:ReqHdr[14], bx
MOV bx, BSEG1
MOV word ptr DS:ReqHdr[16], Bx
MOV BX, OFFSET DS:ReqHdr
INT 02Fh
MOV CX, word ptr DS:ReqHdr[3]
AND CX,8300h
CMP CX,0100h
JE @F
POP CX
LOOP @B
RET
@@:
POP CX
RET
```

SND_REQ ENDP

SND_RH PROC NEAR MOV DTA_ofs,OFFSET DS:ReqHdr MOV DTA_seg,DS MOV AX,DS:RH_Len MOV Data_Len,AX CALL SND_PKT RET SND_RH ENDP

SND_DATA PROC NEAR MOV AX,DT_Len CALL P_WORD MOV Data_Len,AX MOV DTA_seg,BSEG1 MOV DTA_ofs,OFFSET BSEG1:Buffer1 CALL SND_PKT RET SND_DATA ENDP

SET_BDPROCNEAR

MOV AL,080h ; Set DLAB in Register 3 of COM1 MOVDX,COM+3 OUT DX,AL JMP SHORT \$+2 MOV AL,01h ; Set baud LSB to 1 (Reg. 0)

MOVDX,COM OUT DX,AL JMP SHORT \$+2

MOV AL,0 ; Set baud MSB to 0 (Reg. 1) MOVDX,COM+1 OUT DX,AL JMP SHORT \$+2

MOV AL,03h ; Set DLAB=0, Reg. 3 = 8-N-1

```
MOVDX, COM+3
  OUT DX,AL
  JMP SHORT $+2
  RET
SET_BD ENDP
XMIT_AL PROC NEAR
; PUSH DI
; MOV DI, offset sent_byte
; CALL PR_MSG
; POP DI
  MOVDX, COM
  OUT DX, AL
  RET
XMIT_AL ENDP
X_WAIT PROC NEAR
PUSH AX
TBE:
; MOV AL, '+'
; CALL P_AL
  X_OK
  JNETBE
 POP AX
  RET
X_WAITENDP
```

RECV_AL PROC NEAR

MOVDX, COM

IN AL, DX RET RECV_AL ENDP P_AL PROC NEAR PUSH AX PUSH BX MOV AH, OEH MOV BH,0 INT 10h POP BX POP AX RET P_AL ENDP PRNT_AL PROC NEAR ; PUSH AX ; PUSH BX ; MOV AH, OEH ; MOV BH,0 ; INT 10h ; POP BX ; POP AX RET PRNT_AL ENDP ; Print Message Procedure ; On Entry: ; DI Points to Message in the data ; segment which is terminated by a \$ PR_MSG PROC NEAR PUSH AX PUSH BX MSGLOOP: MOV AL, DS: [DI] CMP AL,'\$' JE MSGDONE CALL P_AL

INC DI JMP MSGLOOP MSGDONE: POP BX POP AX RET PR_MSG ENDP ; p_word prints the value in ax as a hexadecimal p_word proc near push ax call p_byte shr ax,8 call p_byte pop ax ret p_word endp ; p_byte prints the value in al as a hexadecimal p_byte proc near push bx push ax mov bx,ax shr bx,4 and bx,000Fh mov al,ds:hexes[bx] call p_al pop bx push bx and bx,000Fh mov al,ds:hexes[bx] call p_al pop ax pop bx ret p_byte endp

R_WAIT PROC NEAR PUSH AX RxRDY: R_OK JNERxRDY POP AX RET R_WAIT ENDP R_FLUSH PROC NEAR FLUSHER: R_OK JNEFL_DONE CALL RECV_AL JMP FLUSHER FL_DONE: RET R_FLUSH ENDP ONE_FLUSH PROC NEAR R_OK JNE @F CALL RECV_AL 00: RET ONE_FLUSH ENDP rcv_pkt proc near push es push di push ax push bx push cx

MOV DI, OFFSET DS:waiting_msg

```
CALL PR_MSG
mov ax,CS:DTA_ofs
shr ax,4
add ax,CS:DTA_seg
mov CS:DTA_seg,ax
mov ax,CS:DTA_ofs
and ax,000Fh
mov CS:DTA_ofs,ax
mov es,ds:DTA_seg
mov di,ds:DTA_ofs
call r flush
mov al, '-'
call p_al
snd_chk: ;(or money order)
 call x_wait
x_rdy
mov cx,65535
@@:
r_ok
 je rcv
 loop @b
 r_ok
 jne snd_chk
rcv:
 recv_alm
 cmp al,ds:command
 jne snd_chk
xfer_com:
 cli
 r_waitm
 recv_alm
 mov byte ptr ds:data_len[0],al
 r_waitm
 recv_alm
 mov byte ptr ds:data_len[1],al
 mov cx,ds:data_len
 cmp cx,0
```

```
je rcv_done
rcv_loop:
 sti
 r_waitm
 cli
 recv_alm
mov byte ptr es:[di],al
 inc di
 loop rcv_loop
rcv_done:
 sti
mov di, offset received_msg
call pr_msg
pop cx
pop bx
pop ax
pop di
pop es
ret
rcv_pkt endp
snd_pkt proc near
push es
push di
push ax
push bx
push cx
mov di, offset waiting_msg
call pr_msg
mov ax,CS:DTA_ofs
 shr ax,4
add ax,CS:DTA_seg
mov CS:DTA_seg,ax
mov ax,CS:DTA_ofs
 and ax,000Fh
mov CS:DTA_ofs,ax
```

```
@@: call r_wait
r_rdy
jne @B
call delay
mov di, offset ok_msg
call pr_msg
mov al,ds:command
call delay
call x_wait
call xmit_al
mov al,byte ptr ds:data_len[0]
call delay
call x_wait
call xmit_al
mov al,byte ptr ds:data_len[1]
call delay
call x_wait
call xmit_al
call one_flush
mov cx,ds:data_len
cmp cx, 0
 je snd_done
mov es,ds:DTA_seg
mov di,ds:DTA_ofs
snd_loop:
call x_wait
mov al, byte ptr es: [di]
call xmit_al
 inc di
 loop snd_loop
snd_done:
mov di, offset sent_msg
call pr_msg
pop cx
pop bx
pop ax
pop di
pop es
 ret
```

snd_pkt endp

CSEGENDS

ENDMAIN

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