Graduate School Introductory Computational Simulation Course Pedagogy

by

Laura L. Proctor

Submitted to the Department of Computation for Design and Optimization in partial fulfillment of the requirements for the degree of

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Abstract

Numerical methods and algorithms have developed and matured vastly over the past three decades now that computational analysis can be performed on almost any personal computer. There is a need to be able to teach and present this material in a manner that is easy for the reader to understand and be able to go forward and use. Three popular course at MIT were without lecture notes; in this thesis the lecture notes are presented. The first chapter covers material taught in Numerical Methods for Partial Differential Equations (2.097/6.339/16.920) specifically the Integral Equation Methods section of this course, chapter two shows the notes for the course Introduction to Numerical Simulation (2.096/6.336/16.910), and chapter three contains the notes for the class Foundations of Algorithms and Computational Techniques in Systems Biology (6.581/20.482). These course notes give a broad overview of many algorithms and numerical methods that one can use to solve many problems that span many fields - from biology to aerospace to electronics to mechanics.

Thesis Supervisor: Jacob K. White Title: Cecil H. Green Professor

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Introduction

Numerical methods and algorithms have developed and matured vastly over the past three decades now that computational analysis can be performed on almost any personal computer. There is a need to be able to teach and present this material in a manner that is easy for the reader to understand and be able to go forward and use. Three popular course at MIT were without lecture notes; in this thesis the lecture notes are presented. The first chapter covers material taught in Numerical Methods for Partial Differential Equations (2.097/6.339/16.920) specifically the Integral Equation Methods section of this course, chapter two shows the notes for the course Introduction to Numerical Simulation (2.096/6.336/16.910), and chapter three contains the notes for the class Foundations of Algorithms and Computational Techniques in Systems Biology (6.581/20.482). These course notes give a broad overview of many algorithms and numerical methods that one can use to solve many problems that span many fields - from biology to aerospace to electronics to mechanics.

Chapter 1

Integral Equation Methods

Numerical Methods for Partial Differential Equation is a course that covers several techniques for discretizing partial differential equations in order to solve them. Very often partial differential equations do not have analytic solutions and one needs to apply an appropriate method to solve them. In this thesis, the Integral Equation Methods used for solving Partial Differential Equations is covered.

1.1 Discretization of Boundary Integral Equations

Numerical Methods for PDEs

Boundary Element Methods, Lecture 1 Introduction to Discretization of Boundary Integral Equations

L. Proctor, S. De, C. Coelho, D. Willis, X. Wang, & J. White

November 23, 2008

1 Module Outline - 6 Lectures

2 Outline for Today

Background

Exterior versus interior problems Point source approach

Test Function Selection

Collocation Method

Galerkin Method

Some issues in 3D

Singular integrals

3 Background

3.1 Interior vs Exterior Problems

Interior

Exterior



Temperature in a tank Ice cube in a bath

 $\frac{\text{What is the heat distribution?}}{\text{Heat flow} = \text{Thermal conductivity } \int_{surface} \frac{\partial T}{\partial n}$

3.1.1 The Interior Problem

Example: Heat Distribution in a Tank



How does one determine the heat distribution using the finite difference method?

Use the above example of heat distribution in a tank, to see what is meant by an interior problem. This is a Dirichlet problem because $T(\vec{x})$, the temperature, is defined on the surface $\vec{x} \in \Gamma$. The steady state two dimensional heat flow equation is defined using the Laplace Equation, $\nabla^2 T(\vec{x}) = 0$. The domain of this problem is clearly the <u>interior</u> of the tank, $\vec{x} \in \Omega$.

3.1.2 The Exterior Problem

Up until this time, all problems that have been studied using finite-element and finite-difference methods have been *interior* problems, but now, we begin to wonder how to form an <u>exterior</u> problem. This poses some problems for these other methods such as generating the grid as well as mesh truncation.



Example: Ice Cube in a Bath

Above is a problem showing an *ice cube in a bath* which is used to indicate an exterior problem. Again, this is a Dirichlet problem because $T(\vec{x})$ is defined on the surface $\vec{x} \in \Gamma$ and described by the steady state two dimensional heat flow equation defined using Laplace's Equation, $\nabla^2 T(\vec{x}) = 0$. The problem domain is the infinitely extending region exterior to the ice cube. A point that we will expand in a later lecture is that with exterior problems, an additional boundary condition is needed to specify what happens at a large distance away from our point source. Assuming there are no heat sources exterior to the cube will impose the following *radiation* boundary condition

$$\lim_{\|\vec{x}\| \to \infty} T(\vec{x}) \to 0.$$

Suppose that for this specific problem, we are only interested in what occurs at the surface of the cube. It seems inefficient to use the finite-difference or finite-element methods where one needs to compute the temperature everywhere in Ω .



3.2 Examples

3.2.1 Computation of Capacitance

What is the capacitance?





In the example in the slide, the yellow plates form a parallel-plate capacitor with an applied voltage V. In this 3-D electrostatics problem, the electrostatic potential Ψ satisfies Laplace's equation $\nabla^2 \Psi(x) = 0$ in the region exterior to the plates, and the potential is known on the surface of the plates (Dirichlet boundary condition). Furthermore, far from the plates,

$$\lim_{\|\vec{x}\| \to \infty} \Psi(\vec{x}) \to 0.$$

(Exterior Radiation Boundary Condition to be studied further in a future lecture). The value of interest is the capacitance, C, which satisfies

q = CV

where q, the net charge on one of the plates, is given by the surface normal of the potential integrated over one plate and scaled by a dielectric permittivity.

Note 1

This is a typical application example, determining the charge density on the surface of conducting plates given an applied voltage. In this particular example, the top plate potential is $\Psi = 0.5V$ and the bottom plate potential is $\Psi = -0.5V$, where V is the voltage noted in the figure.

For this exterior Dirichlet problem, one can write an integral equation that relates the surface charge density on the plates σ to the potential on the plates. This integral equation, $\Psi(\vec{x}) = \int_{\Gamma} \frac{1}{\|\vec{x}-\vec{x}'\|} \sigma(\vec{x}') dS'$, is often referred to by physicists as the superposition integral. In the integral equation, x is any point on the plate surfaces and the surface being integrated over is the union of the top and bottom plate surfaces. Note that the integration surface is not a connected domain, but this presents no difficulties.

3.2.2 Drag Force in a Microresonator



Note 2

Example 2: Drag force in a MEMS device

The example in the slide is a microresonator, it is a structure that can be made to vibrate using electrostatic forces. The changing character of those vibrations can be used to sense rotation. The particulars of how the microresonator operates is not directly relevant to our discussion of integral equations, except for one point. In order to determine how much energy is needed to keep the microresonator vibrating, it is necessary to determine the fluid drag force on comb structures shown in the bottom part of the slide. The fluid is the air surrounding the structure, and at the micron-scale of these devices, air satisfies the incompressible Stokes equation,

$$\nabla^2 u(x) = \nabla p(x) \tag{1}$$
$$\nabla \cdot u(x) = 0$$

where u is the fluid velocity and p is the pressure. By specifying the comb velocity, and then computing the surface pressure and the normal derivative of velocities tangent to the surface, one can determine the net drag force on the comb. Once again, this is a problem in which the known quantities (the comb velocity) and the quantities of interest (the derivative of the tangential components of fluid velocity) are on the surface.

3.2.3 Aircraft Drag



Discretization for F-18 pressure simulation (no lift) Inviscid, Irrotational, Steady Flow Potential flow: $\nabla^2 u(x) = 0$ $\nabla u =$ velocity

Note 3

Example 3: Aircraft Drag

The potential flow model for aircraft drag computation will be discussed in more detail in subsequent lectures, so we only give a brief description here. In order to compute the drag on the wing of an aircraft, one must determine the difference between the wing velocity and the velocity of the air very close to the wing. If the air can be assumed inviscid, irrotational, and incompressible, the velocity is given by the gradient of a scalar potential which satisfies Laplace's equation. The boundary conditions for the Laplace's equation are given as a velocity boundary condition on the aircraft surface, equivalently a Neumann condition on the potential, and it is usually assumed that the potential approaches zero at infinity. The boundary condition at infinity is more subtle than it may seem, as we shall see in later lectures. Finally, it is common to introduce an artificial boundary in the domain, and specify a condition on that boundary to introduce rotational effects.

3.2.4 Capacitance of Microprocessor Signal Lines



Note 4 Example 4: Capacitance of microprocessor signal lines

This last example in the above slide is a picture of the wiring on a microprocessor integrated circuit. A typical microprocessor has millions of wires, so we are only looking at a small piece of a processor. The critical problem in this example is determining how long signals take to get from the output of a logical gate to the input of the next gate. To compute that delay, one must determine the capacitance on each of the wires given in the slide picture. To do so requires computing charges given electrostatic potentials as noted above.

3.3 Advantages of Integral Equation Method

3.3.1 What is common about these examples?

Exterior Problems

MEMS device - fluid (air) creates drag
Aircraft Design - exterior air flow
Signal Line - Exterior fields.
Quantities of interest are on surface
MEMS device - Just want surface traction force
Aircraft Design - Just want surface tangent velocities
Signal Line - Just want surface charge.
Exterior problem is linear and space-invariant
MEMS device - Exterior Stoke's flow equation (linear)
Aircraft Design - Laplace's equation, plus wakes.
Signal line - Laplace's equation in free spee (linear)
But problems are geometrically very complex

3.3.2 Why not use FDM / FEM?



Only need $\frac{\partial T}{\partial n}$ on the surface, but T is computed everywhere. Must truncate the mesh, $\Rightarrow T(\infty) = 0$ becomes T(R) = 0.

Consider the two dimensional exterior heat conduction problem in the above figure in which the temperature is known on the surface of the square. Suppose the quantity of interest is the total heat flow out of the square. The temperature T satisfies

$$\nabla^2 T(x) = 0 \qquad x \in \Omega$$

$$T(x) \text{ given } x \in \Gamma$$

$$\lim_{\|x\| \to \infty} T(x) = 0$$
(2)

where Ω is the infinite domain outside the square and Γ is the region formed by the edges of the square.

Using finite-element or finite-difference methods to solve this problem requires introducing an additional approximation beyond discretization error. It is not possible to discretize all of Ω , as it is infinite, and therefore the domain must be truncated with an artificial finite boundary. In the slide, the artificial boundary is a large ellipse on which we assume the temperature is zero. Clearly, as the radius of the ellipse increases, the truncated problem more accurately represents the domain problem, but the number of unknowns in the discretization increases.

3.4 Point Source Approach

3.4.1 Green's Function

Heat Distribution in 2-D



Green's Function:
$$T = log\left(\sqrt{(x-x_0)^2 + (y-y_0)^2}\right)$$

From basic electrostatics, one knows that in 3-D, the potential field produced by a point charge decays inversely with the distance to the point charge. Since, roughly, one can represent any charge distribution using a sum of point charges, one can express the potential due to a charge density as a sum of point charge potentials. Therefore, point charge potentials play a special role, and are often referred to as Greens' functions for the problem.

In 2-D The potential due to a point charge is:

$$u = \log\left(\sqrt{(x - x_0)^2 + (y - y_0)^2}\right)$$
(3)
$$\forall (x, y) \neq (x_0, y_0)$$

<u>In 3D</u> The potential due to a point charge is:

$$u = \frac{1}{\sqrt{(x - x_0)^2 + (y - y_0)^2 + (z - z_0)^2}}$$

$$\forall (x, y, z) \neq (x_0, y_0, z_0)$$
(4)

$$\underbrace{\operatorname{In} \ 2D}_{\text{then}} \quad \operatorname{If} \ u = \log\left(\sqrt{(x-x_0)^2 + (y-y_0)^2}\right) \\
\operatorname{then} \ \frac{\partial^2 u}{\partial x^2} + \frac{\partial^2 u}{\partial y^2} = 0 \quad \forall \quad (x,y) \neq (x_0,y_0)$$

$$\underbrace{\operatorname{In} \ 3D}_{\frac{\partial^2 u}{\partial x^2} + \frac{\partial^2 u}{\partial y^2} + \frac{\partial^2 u}{\partial z^2} = 0 \quad \forall \quad (x,y,z) \neq (x_0,y_0,z_0)$$

Proof: Just differentiate and see!

In the next few slides, we will use an informal semi-numerical approach to derive the integral form of Laplace's equation. We do this in part because such a derivation lends insight to the subsequent numerical procedures.

To start, recall from basic physics that the potential due to a point charge is related only to the distance between the point charge and the evaluation point. In 2-D the potential is given by the log of the distance, equation (3), and in 3-D the potential is inversely proportional to the distance, equation (4). These functions are sometimes referred to as Green's functions for Laplace's equation, but have the physical interpretation as the potential due to a point charge. We will be studying Green's functions in more depth later on.

▷ **Exercise 1** Show by direct differentiation that the functions in equations (3) and (4) satisfy $\nabla^2 u = 0$, in the appropriate dimension almost everywhere.

3.4.2 Scaling Green's Function

æ

u is given on surface
Surface
$$\frac{\partial^2 u}{\partial x^2} + \frac{\partial^2 u}{\partial y^2} = 0$$
 outside
Let $u = \log\left(\sqrt{(x - x_0)^2 + (y - y_0)^2}\right)$
 $\frac{\partial^2 u}{\partial x^2} + \frac{\partial^2 u}{\partial y^2} = 0$ outside Problem Solved

Boundary conditions are not satisfied!

A simple idea for computing the solution of Laplace's equation outside the square is to let

$$u(x,y) = \alpha \log \sqrt{(x-x_0)^2 + (y-y_0)^2}$$

where (x_0, y_0) is a point inside the square. Clearly u will always satisfy $\nabla^2 u = 0$ outside the square, but u may not match the boundary conditions. By adjusting

 $\alpha,$ it is possible to make sure to match the boundary conditions at at least one point.

This concept is applied to a circle as a simple example of how to match the boundary conditions.



 \triangleright **Exercise 2** Suppose the potential on the surface of the square is a constant. Can you match that constant potential everywhere on the perimeter of the square by judiciously selecting α ?

$$u \text{ is given on surface} \qquad u = \sum_{i=1}^{n} \alpha_i \log \left(\sqrt{(x-x_i)^2 + (y-y_i)^2} \right) = \sum_{i=1}^{n} \alpha_i G(x-x_i, y-y_i)$$
Pick the α_i 's to match the boundary conditions!

To construct a potential that satisfies Laplace's equation and matches the boundary conditions at more points, let u be represented by the potential due to a sum of n weighted point charges in the square's interior. As shown in the slide, we can think of the potential due to a sum of charges as a sum of Green's functions. Of course, we have to determine the weights on the n point charges, and the weight on the i^{th} charge is denoted hereby α_i .

$$\begin{bmatrix} G(x_{t_1} - x_1, y_{t_1} - y_1) & \cdots & G(x_{t_1} - x_n, y_{t_1} - y_n) \\ \vdots & \ddots & \vdots \\ G(x_{t_n} - x_1, y_{t_n} - y_1) & \cdots & G(x_{t_n} - x_n, y_{t_n} - y_n) \end{bmatrix} \begin{bmatrix} \alpha_1 \\ \vdots \\ \alpha_n \end{bmatrix} = \begin{bmatrix} \Psi(x_{t_1}, y_{t_1}) \\ \vdots \\ \Psi(x_{t_n}, y_{t_n}) \end{bmatrix}$$

To determine a system of n equations for the $n \alpha_i$'s, consider selecting a set of n test points, as shown in the slide above. Then, by superposition, for each test point (x_{t_i}, y_{t_i}) ,

$$u(x_{t_i}, y_{t_i}) = \sum_{i=1}^n \alpha_i \log \sqrt{(x_{t_i} - x_0)^2 + (y_{t_i} - y_0)^2} = \sum_{i=1}^n \alpha_i G(x_{t_i} - x_0, y_{t_i} - y_0).$$
(5)

Writing an equation like (5) for each test point yields the matrix equation

$$\begin{bmatrix} G(x_{t_1}-x_1,y_{t_1}-y_1) & \cdots & G(x_{t_1}-x_n,y_{t_1}-y_n) \\ \vdots & \ddots & \vdots \\ G(x_{t_n}-x_1,y_{t_n}-y_1) & \cdots & G(x_{t_n}-x_n,y_{t_n}-y_n) \end{bmatrix} \begin{bmatrix} \alpha_1 \\ \vdots \\ \alpha_n \end{bmatrix} = \begin{bmatrix} \Psi(x_{t_1},y_{t_1}) \\ \vdots \\ \Psi(x_{t_n},y_{t_n}) \end{bmatrix}$$
(6)

The matrix A in equation (6) has some properties worth noting:

- A is dense, that is $A_{i,j}$ never equals zero. This is because every charge contributes to every potential.
- If the test points and the charge points are ordered so that the i^{th} test point is nearest the i^{th} charge, then $A_{i,i}$ will be larger than $A_{i,j}$ for all j.

The 2^{nd} item above seems to suggest that A is diagonally dominant, but this is *not* the case. Diagonal dominance requires that the absolute *sum* of the offdiagonal entries is smaller than the magnitude of the diagonal. The matrix above easily violates that condition.

 \triangleright **Exercise 3** Determine a set of test points and charge locations for the 2-D square problem that generates an A matrix where the magnitude of the diagonals are bigger than the absolute value of the off-diagonals, but the magnitude of the diagonal is smaller than the absolute sum of the off-diagonals.

3.4.3 Source Point Locations

Where should the sources be located?





Close to the boundary



Problems with these placements:



Close to the boundary

Clustered in the center

3.4.4 Computational Results



It is possible to construct a numerical scheme for solving exterior Laplace problems by adding progressively more point charges so as to match more boundary conditions. In the above graph, we show an example of using such a method to compute the potential exterior to a circle of radius 10, where the potential on the circle is given to be unity. In the example, charges are placed uniformly on a circle of radius 9.5, and test points are placed uniformly on the radius 10 circle. If 20 point charges are placed in a circle of radius 9.5, then the potential produced will be exactly one only at the 20 test points on the radius 10 circle. The potential produced by the twenty point charges on the radius 10 circle is plotted in the lower left corner of the slide above. As might be expected, the potential produced on the radius 10 circle is exactly one at the 20 test points, but then oscillates between 1 and 1.2 on the radius 10 circle. If 40 charges and test points are used, the situation improves. The potential on the circle still oscillates, as shown in the lower right hand corner, but now the amplitude is only between 1 and 1.004.

3.5 Charge Density

Want to smear point charges to the surface



Results in an integral equation

$$\Psi(\vec{x}) = \int_{\Gamma} G(\vec{x}, \vec{x}') \sigma(\vec{x}') dS'$$
(7)

How do we solve the integral equation?

In equation (7) for which variable are we trying to solve?



$$\Psi(x,y) = \int_{\Gamma} \sigma(x',y') \log \sqrt{(x-x')^2 + (y-y')^2} dx' dy'$$

Single Layer Potential

The oscillating potential produced by the point charge method is due to the rapid change in potential as the separation between evaluation point and point charge shrinks. If the point charges could be smeared out, so that the produced potential did not rise to infinity with decreasing separation, then the resulting computed potential would not have the oscillation noted on the previous slide. In addition, it makes the most sense to smear the point charges onto the surface, as then the charge density and the known potential have the same associated geometry. The result is the integral equation (7), where now the unknown is a charge density on the surface and the potential due to that charge density is given by the well-known superposition integral. In the case of two or three dimensional Laplace problems, $G(\vec{x}, \vec{x}')$ can be written as $\hat{G}(\vec{x} - \vec{x}')$, as the potential is only a function of distance to the charge density and not a function of absolute position. For such a Green's function, this equation is,

$$\Psi(\vec{x}) = \int_{\Gamma} \hat{G}(\vec{x} - \vec{x}')\sigma(\vec{x}')dS', \qquad (8)$$

which one may recognize from system theory as a convolution integral. This connection is quite precise. A space-invariant system has an impulse response, which is usually referred to as a Green's function. The output, in this case, the potential, is a convolution of the impulse response with the input, in this case, the charge density. Such an integral form of the potential is referred to as a single layer potential.

The single layer potential is an example of a class of integral equations known as "Fredholm integral equation of the First Kind". A Fredholm integral equation of the Second Kind results when the unknown charge density exists not only under the integral sign but also outside it. An example of such an equation is

$$\Psi(\vec{x}) = \sigma(\vec{x}) + \int_{\Gamma} K(\vec{x} - \vec{x}')\sigma(\vec{x}')dS'.$$
(9)

Fredholm integral equations, in which the domain of integration is fixed, usually arise out of boundary value problems. Initial value problems typically give rise to the so-called Volterra integral equations, where the domain of integration depends on the output of interest. For example, consider the initial value problem

$$\frac{dx(t)}{dt} = tx(t); \quad t \in [0, T], \ T > 0.$$
$$x(t = 0) = x_0$$

The "solution" of this equation is the following Volterra integral equation:

$$x(t) = x_0 + \int_0^t \xi x(\xi) d\xi.$$

4 Basis Functions

4.1 Basic Idea



Basis Functions can be used to approximate the surface charge density in a similar way in which they approximate geometry for finite elements.

Numerical solution of the single layer potential

As we have studied extensively in the finite-element section of the course, one approach to numerically computing solutions to partial differential equations is to represent the solution approximately as a weighted sum of basis functions. Then, the original problem is replaced with the problem of determining the basis function weights. In finite-element methods, the basis functions exist in a volume, for integral equations they typically exist on a surface. For 2-D problems that means the basis functions are restricted to curves and in 3-D the basic functions are on physical surfaces.

As an example, consider the circle in the above figure. One could try to represent the charge density on the circle by breaking the circle into n sub-arcs, and then assume the charge density is a constant on each sub-arc. Such an approach is not commonly used. Instead the geometry is usually approximated along with the charge density. In this example case, shown in the center right of the slide, the sub-arcs of the circle are replaced with straight sections, thus forming a polygon. The charge density is assumed constant on each edge of the polygon. The result is a piecewise constant representation of the charge density on a polygon.

4.2 Geometric Approximation



The idea that both the geometry and the unknown charge density has been approximated is not actually a new issue. As shown in the figure in the above slide, if FEM methods are used to solve an interior problem, and triangular elements are used, then the circle is approximated to exactly the same degree as when straight sections replace the sub-arcs for the surface integral equation. As shown at the bottom of the above slide, we can substitute the basis function representation into the integral equation, but then we should also note that the integral is now over the approximated geometry. It is common, but not mathematically justified, to ignore the errors generated by the geometry approximation. We will also ignore the error in the geometric approximation in our analyses, just for simplicity. In the case of polygonal geometries, there is no geometric approximation, so there is at least one case where the assumption is precise. It should be noted, however, that there are often analytic results only for smooth geometries, and then before making comparisons to such analytic results, it is necessary to examine the effect of the approximated geometry.

If the original problem is a polygon



there is no geometric approximation



4.2.1 Piecewise Constant Straight Sections



Pick a set of n Points on the surface
 Define a new surface by connecting points with n lines.
 Define φ_l(x) = 1 if x is on line l, otherwise, φ_l(x) = 0

$$\Psi(\vec{x}) = \int_{\substack{\text{approx}\\\text{surface}}} G(\vec{x}, \vec{x}') \sum_{i=1}^{n} \alpha_i \varphi_i(\vec{x}') dS' = \sum_{i=1}^{n} \alpha_i \int_{line \ l_i} G(\vec{x}, \vec{x}') dS'$$

How do we determine the α_i 's?

We complete the description of using constant charge densities on straight sections as the basis. If we substitute this example basis function into the integral equation, as is done above, the result is to replace the original integration of the product of the Green's function and the density with a weighted sum of integrals over straight lines of just the Green's function. The next step is then to develop an approach for determining the weights, denoted here by α_i 's.

5 Test Points

5.1 Residuals

5.1.1 Definition and Minimization

$$R(\vec{x}) = \Psi(\vec{x}) - \int_{\substack{\text{approx} \\ \text{surface}}} G(\vec{x}, \vec{x}') \sum_{i=1}^{n} \alpha_i \varphi_i(\vec{x}') dS'$$
(10)

Pick the α_i 's to minimize $R(\vec{x})$

General Approach: Pick a set of test functions ϕ_1, \ldots, ϕ_n and force $R(\vec{x})$ to be

orthogonal to the set:

$$\int \phi_i(\vec{x}) R(\vec{x}) dS = 0 \quad \forall i \tag{11}$$

One way of assessing the accuracy of the basis function based approximation of the charge density is to examine how well the approximation satisfies the integral equation. To be more precise, we define the residual associated with the integral equation and an approximate solution, equation (10). Note that $R(\vec{x})$ is just the difference between the given potential on the surface and the potential produced by the approximated charge density. Note also that the equation is now over the approximate geometry and therefore \vec{x} and \vec{x}' are both on the approximated surface.

If the representation satisfies the integral equation exactly, then the residual $R(\vec{x})$ will be zero for all \vec{x} and the approximate solution is equal to the exact solution (provided the integral equation has a unique exact solution ... more on this later). In general, though, this is not possible, and instead we will try to pick the basis function weights, the α_i 's, to somehow minimize $R(\vec{x})$. One approach to minimizing $R(\vec{x})$ is to make it orthogonal to a collection of test functions, which may or may not be related to the basis functions, equation (11). Enforcing orthogonality in this case means ensuring that the integral of the product of $R(\vec{x})$ and $\phi(\vec{x})$ over the surface is zero.

5.1.2 Residual Minimization Using Test Functions

$$\frac{\int \phi_i(\vec{x}) R(\vec{x}) dS = 0}{\int \phi_i(\vec{x}) \Psi(\vec{x}) dS - \iint_{\substack{\text{approx} \\ \text{surface}}} \phi_i(\vec{x}) G(\vec{x}, \vec{x}') \sum_{j=1}^n \alpha_j \varphi_j(\vec{x}') dS' dS = 0$$
(12)

We will generate different methods by choosing the ϕ_1, \ldots, ϕ_n Collocation : $\phi_i(\vec{x}) = \delta(\vec{x} - \vec{x}_{t_i})$ (point matching) Galerkin Method : $\phi_i(\vec{x}) = \varphi_i(\vec{x})$ (basis = test) Weighted Residual Method : $\phi_i(\vec{x}) = 1$ if $\varphi_i(\vec{x}) \neq 0$ (averages) As noted in the equation (12), by substituting the definition of the residual into the equation (11), it is possible to generate n equations, one for each test function. The generated equation has two integrals. The first is a surface integral of the product of the given potential with the test function. The second integral is a double integral over the surface. The integrand of the double integral is a product of the test function, the Green's function, and the charge density representation.

Three different numerical techniques can be derived by altering the test functions.

5.2 Collocation

The collocation method, described in the above slide, uses shifted impulse functions as test functions, $\phi_i(\vec{x}) = \delta(\vec{x} - \vec{x}_i)$. Impulse functions, also called "delta" functions, have a *sifting* property when integrated with a smooth function $f(\vec{x})$,

$$\int f(\vec{x})\delta(\vec{x}-\vec{x}_i)dx = f(\vec{x}_i).$$

Impulse functions are also referred to as generalized functions, and they are specified only by their behavior when integrated with a smooth function. In the case of the impulse function, one can think of the function as being zero except for a very narrow interval around \vec{x}_i , and then being so large in that narrow interval that $\int \delta(\vec{x} - \vec{x}_i) dx = 1$.

As the summation equation in the middle of the above slide indicates, testing with impulse functions is equivalent to insisting that $R(\vec{x}_i) = 0$, or in words, that the potential produced by the approximated charge density should match the given potential at n test points. That the potentials match at the test points gives rise to the method's name, the point where the potential is exactly matched is "co-located" with a set of test points.

The $n \times n$ matrix equation at the bottom of the above slide has as its righthand side the potentials at the test points. The unknowns are the basis function weights. The j^{th} matrix element for the i^{th} row is the potential produced at test point x_i by a charge density equal to basis function φ_i .

5.2.1 Centroid Collocation for Piecewise Constant Bases



In the above slide, a specific collocation algorithm is described. First, the basis being used is the constant charge density on n straight sections or lines, as described above. Note that therefore the geometry is being approximated. Second, the collocation points being selected are the centroids of the basis functions, in this case just the center of each straight line. Note that the collocation point is on the approximated geometry, not the original geometry. So, one can think of the problem as having been restated to be on a polygon instead of the original circle, but then the replacement interpretation does not hold.

In collocation, or point-matching, the charge densities on each of the straight lines are selected so that the resulting potential at the line centers matches the given potential. As the equations on this slide make clear, the matrix element $A_{i,j}$ is the potential at the center of line *i* due to a unit charge density along line *j*.

It should be noted that the matrix A is dense, the charge on line j contributes to the potential everywhere. Also note that if line j is far away from line i, then

$$A_{i,j} \approx \text{length}(line_j) \times G(\vec{x}_{t_i}, \vec{x}_{t_j})$$
 (13)

▷ **Exercise 4** Suppose we are using piecewise constant centroid collocation to solve a 2-D Laplace problem, so $G(x, y, x', y') = \log \sqrt{(x - x')^2 + (y - y')^2}$. Roughly how far apart do line sections i and j have to be for equation (13) to be accurate to within one percent? Assume line j has length of one. Does your answer depend on the orientation of line j? Does your answer depend on the orientation of line j? Does your answer depend on the other, do you see why?) ■

5.2.2 Centroid Collocation Generates Nonsymmetric A

$$\Psi(\vec{x}_{t_i}) = \sum_{j=1}^n \alpha_j \overbrace{\int_{\text{line } j} G(\vec{x}_{t_i}, \vec{x}') dS'}^{A_{i,j}}$$

$$A_{1,2} = \int_{\text{line } 2} G(\vec{x}_{t_1}, \vec{x}') dS' \neq \int_{\text{line } 1} G(\vec{x}_{t_2}, \vec{x}') dS' = A_{2,1}$$
(14)

Consider the two line sections, l_1 and l_2 given in the above figure. For Laplace problems, $G(\vec{x}, \vec{x}') = G(\vec{x}', \vec{x})$, which suggests a symmetry in the underlying integral equation that is not represented in the collocation discretization. This asymmetry is shown in equation (14) by noting that $A_{1,2} \neq A_{2,1}$. That is, the potential at the center of l_2 due to a unit charge density on l_1 is not equal to the potential at the center of l_1 due to a unit charge on l_2 .

X,

It is possible to scale the variables to improve the symmetry, consider a change of variables

$$\hat{\alpha}_i = \alpha_i \times \text{length}(line_i).$$

In this change of variables, the unknowns $\hat{\alpha}_i$ are now the net line charges rather than the line charge densities. In this new system, $\hat{A}\hat{\alpha} = \Psi$, where the elements of the matrix \hat{A} are given by

$$\hat{A}_{i,j} = \frac{1}{\text{length}(line_j)} \int_{line_j} G(\vec{x}_{t_i}, \vec{x}') dS'.$$

Under the change of variables, if line j is far away from line i, then

$$\hat{A}_{i,j} \approx G(\vec{x}_{t_i}, \vec{x}_{t_j}) \approx \hat{A}_{j,i}.$$
(15)

In other words, the elements of \hat{A} corresponding to distant terms are approximately symmetric.

Is
$$\Psi(\vec{x})$$
 due to $\sigma(\vec{x}')$ the same as $\Psi(\vec{x}')$ due to $\sigma(\vec{x})$?



Green's Function is due to $\log R$

 \triangleright **Exercise 5** Give an example which shows that the scaled entries of \hat{A} can be far from symmetric. Assume we are using piecewise constant straight sections with centroid collocation and the 2-D Laplace's equation Green's function.

5.3 Galerkin



If $G(\vec{x}, \vec{x}') = G(\vec{x}', \vec{x})$ then $A_{i,j} = A_{j,i} \Rightarrow \mathbf{A}$ is symmetric

In the Galerkin method, the test functions are equal to the basis functions. In particular, one generates n equations for the basis function weights by insisting that $R(\vec{x})$ is orthogonal to each of the basis functions. Enforcing orthogonality corresponds to setting

$$\int arphi(ec{x}) R(ec{x}) dS = 0$$

and substituting the definition of $R(\vec{x})$ into the orthogonality condition yields the equation in the center of the above slide.

Note that the Galerkin method yields a system of n equations, one for each orthogonality condition, and n unknowns, one for each basis function weight. Also, the system does not have the potential explicitly as the right hand side. Instead, the i^{th} right-hand side entry is the average of the product of the potential and the i^{th} basis function.

5.3.1 Galerkin for Piecewise Constant Bases



In the Galerkin method, the basis has constant charge density on n straight sections or lines. We will think of the problem as having been restated to be on a polygon instead of the original circle. The charge densities on each of the straight lines are selected so that the resulting line averaged potential matches the line averaged given potential. As the equations on the above slide make clear, the matrix element $A_{i,j}$ is the average potential over line i, scaled by the length of line i, due to a unit charge density along line j. As with the collocation method, the matrix A is dense because the the charge on line j contributes to the averaged potentials everywhere. Also note that if line j is far away from line i, then

$$A_{i,j} \approx \text{length}(line_j) \times \text{length}(line_i) \times G(\vec{x}_{t_i}, \vec{x}_{t_j})$$
(16)

▷ **Exercise 6** Suppose we are using piecewise constant centroid collocation to solve a 2-D Laplace problem, so $G(x, y, x', y') = \log \sqrt{(x - x')^2 + (y - y')^2}$. Roughly how far apart do line sections *i* and *j* have to be for equation (16) to be accurate to within one percent? Assume line *j* has length of one. Does your answer depend on the orientation of line *j*? Does your answer depend on the orientation of line *i*? (Your answer should be different than the answer you gave for the collocation method. Do you see why?) ■

5.4 Summary

Compare the Collocation and Galerkin methods on a two-dimensional circle.



What do the test functions look like? What do the Residuals look like?



Collocation Method



Galerkin Method



6 Issues in 3D

6.1 Geometric Representation

6.1.1 Introduction

Example: Ship's Hull More errors are introduced with expansion of dimensions



Note 6

"Leaky Panels"

Many papers in the literature on solving integral equations refer to "panel methods". The name is derived from the idea of breaking a surface into flat panels. In the application area of analyzing ocean wave forces on ship hulls, panel methods are commonly used. However, it is not possible to represent a curved hull with quadrilateral flat panels. Researchers in the area often create a best fit panelled surface in which there are gaps between the edges of the panels. Such a discretization technique is often referred to as using "leaky panels", a very compelling image.



Consider solving the integral equation where the surface is the surface of the cube shown. The first step, as we have mentioned in previous lectures, is to develop a basis in which to represent the surface charge density σ .

The cube pictured in the slide has had its surface divided into panels, and a basis is derived from the panels. In particular, one can associate a basis function φ_j with each panel j by assigning $\varphi_j(\vec{x})$ the value one when \vec{x} is a point on panel j, and setting $\varphi_j(\vec{x}) = 0$ otherwise. If σ is approximated by a weighted combination of these basis functions, then the approximation is a piecewise constant representation of the charge density on the surface of the cube.

A few aspects of this basis set should be noted.

• The basis functions are orthogonal, that is if $i \neq j$,

$$\int \varphi_j(\vec{x})\varphi_i(\vec{x})dx = 0.$$

• These basis functions are normalized with respect to l_{∞} , not l_2 . That is, $\|\varphi\|_{\infty} = 1$ but

$$\|\varphi_j\|_2^2 = \int \varphi_i(\vec{x})\varphi_i(\vec{x})dx = \text{panel area.}$$

6.2 Centroid Collocation

Put collocation points at panel centroids


After one has decided on a basis with which to approximately represent the surface charge density, the next step is to develop a system of equations from which to determine the basis weights, denoted as the α_i 's. The most commonly used approach to forming such a system is to use collocation, though Galerkin methods are also quite widely used. Recall that in collocation, the basis function weights are determined by ensuring the the integral equation is exactly satisfied at a collection of "collocation" points. For panel methods, the most common choice for the position of the collocation points are the panel centroids, as shown in the cube diagram above.

The equation in the top of the above slide relates the potential at collocation point \vec{x}_{c_i} to the weights for the panel-based basis functions. To see how the equation was derived, consider evaluting the potential at the i^{th} collocation point using the original integral equation

$$\Phi(\vec{x}_{c_i}) = \int_{surface} \frac{1}{\|\vec{x}_{c_i} - \vec{x}'\|} \sigma(\vec{x}') dS', \qquad (17)$$

where Φ is the know potential on the problem surface and σ is the unknown charge density. Substituting the approximate representation for σ ,

$$\sigma(\vec{x}) \approx \sum_{j=1}^{n} \alpha_j \varphi_j(\vec{x})$$

into the integral equation results in

$$\Phi(\vec{x}_{c_i}) = \int_{surface} G(\vec{x}_{c_i}, \vec{x}') \sum_{j=1}^n \alpha_j \varphi_j(\vec{x}') dS', \qquad (18)$$

where $G(\vec{x}, \vec{x}') \equiv \frac{1}{\|\vec{x} - \vec{x}'\|}$ is used to simplify the formula. Exploiting the fact that $\varphi_j(\vec{x}) = 1$ if \vec{x} is on panel j, and zero otherwise, results in the formula at the top of the above slide.

The system of equations from which to determine the basis function weights is given in the lower corner of the slide. The right hand side of the system is the vector of known potentials at the collocation points. The i, j^{th} element of the matrix A is the potential produced at collocation point i due to a unit charge density on panel j. The vector of α 's are the unknown panel charge densities.

 \triangleright **Exercise 7** Determine a scaling of the α 's ($\hat{\alpha}_i = c_i \alpha_i$) such that the scaled matrix \hat{A} has the property

$$A_{i,j} \approx \frac{1}{\|\vec{x}_{c_i} - \vec{x}_{c_j}\|}$$

when $\|\vec{x}_{c_i} - \vec{x}_{c_j}\|$ is much larger than a panel diameter.

6.2.1 Calculating Matrix Elements



In order to calculate the matrix entries for the system of equations described in the previous slide, recall that $A_{i,j}$ is the potential produced at collocation point *i* due to a unit charge density on panel *j*. The formula for $A_{i,j}$ is given on the top right of the above slide.

The figure on the left of the above slide is a diagram of how one typically computes the panel integral given on the top right. First, consider a shift and rotation of the coordinate system so that the panel lies in the x-y plane at z = 0, with the panel's center at x = 0, y = 0. The figure in the top left shows the panel in the shifted and rotated coordinate system. Note that the collocation point must also be placed in the new coordinate system.

If panel j is reasonably well separated from collocation point i, it is possible to approximate the integral given in the top right by a single point quadrature. More specifically, one could approximates the integral of $\frac{1}{\|\vec{x}_{c_i} - \vec{x}'\|}$ by a product of $\frac{1}{\|\vec{x}_{c_i} - \vec{x}_{centroid_j}\|}$ and the panel area. As show in the middle figure, a single point quadrature is like treating the panel as if it were a point charge at the panel's centroid, where the point charge's strength is equal to the panel area. If the collocation point is close to the panel, then a single point quadrature will be insufficiently accurate. Instead, a more accurate four point quadrature scheme would be to break the panel into four subpanels, and then treat each of the subpanels as point charges at their respective centers. This simple idea is shown in the figure at the bottom of the above slide. This four point scheme is

equivalent to

$$\int_{panel_j} \frac{1}{\|\vec{x}_{c_i} - \vec{x'}\|} dS' \approx \sum_{j=1}^4 \frac{0.25 * Area}{\|\vec{x}_{c_i} - \vec{x}_{point_j}\|}$$

If the panel is a unit square in the x-y plane whose center is at the coordinate system origin, then the four \vec{x}_{point_j} 's are (x, y, z) = (0.25, 0.25, 0), (x, y, z) = (-0.25, 0.25, 0), (x, y, z) = (-0.25, -0.25, 0), and (x, y, z) = (0.25, -0.25, 0).

6.2.2 Calculating "Self Term"



The diagonal terms $A_{i,i}$ can not be computed using the quadrature approximation given on the previous slide. To see this, consider the figure at the top left of the above slide, where a panel has been shifted and rotated into the x-y plane, and the collocation point is the center of the panel. The integral that must be computed is given on the right side of the top of the above slide.

As shown in the middle of the slide, using a single point quadrature scheme will fail, because the distance between the point charge approximation to the panel and the collocation point will be zero. Therefore, the single point formula will require computing the reciprocal of zero, which is infinite. The problem is that the integrand in

$$\int_{panel_{i}} \frac{1}{\|x_{c_{i}} - x'\|} dS'$$
(19)

is singular. That is, the integrand approaches infinity at a point x' which is in the domain of integration. What is not so obvious is that (19) is an integrable singularity. Therefore, even though the integrand approaches infinity at some point, the "area under the curve" is finite.



In the above slide, we both show that

$$\int_{panel_i} \frac{1}{\|\vec{x}_{c_i} - \vec{x}'\|} dS'$$

is integrable, and also give an idea about how to compute the integral. As shown in the slide, first rotate and shift the coordinate system so that the panel is in the x-y plane at z = 0, and so that the collocation point (or equivalently the panel centroid) is at the origin. In this new coordinate system, the integral can be written as

$$A_{i,i} = \int_{panel(rs)_i} \frac{1}{\|\vec{x}'\|} dS'$$

where the notation panel(rs) is used to indicate that the integral is over the rotated and shifted panel.

On the top left of the above slide, a circular disk of radius R and center at the collocation point is inscribed in the rotated panel. In the equations that follow the figure, it is noted that the panel integral can be recast as the sum of an integral over the disk plus an integral over the rest of the panel. The integrand in the integral over the rest of the panel is no longer singular, but the integrand in the integral over the disk is still singular.

The integral over the disk can be computed analytically by using a change of variables. After rotating and shifting the panel, the disk is in the x-y plane and its center, equal to the collocation point, is at zero. Therefore,

$$\int_{disk} \frac{1}{\|\vec{x}_{c_i} - \vec{x}'\|} dS' = \int_{disk} \frac{1}{\|\vec{x}'\|} dS'$$

Apply a change of variables as transformations of two-dimensional regions. Recall this mapping from an earlier lecture in the finite difference method. Suppose that a region $\hat{\Omega}$ in the *r*- θ plane is transformed one-to-one into the region Ω by differentiable equations of the form

$$x = rcos\theta, \qquad y = rsin\theta.$$

Any function f(x, y) defined on Ω can be thought of as a function $f(x(r, \theta), y(r, \theta))$ on $\hat{\Omega}$. The integrals of these functions are related by

$$\iint_{\Omega} f(x,y) dx dy = \iint_{\hat{\Omega}} f(x(r,\theta), y(r,\theta)) |J(r,\theta)| dr d\theta$$

where $J(r,\theta)$ is the Jacobian determinant of the coordinate transformation where

$$J(r,\theta) = \left| \begin{array}{cc} \frac{\partial x}{\partial r} & \frac{\partial x}{\partial \theta} \\ \frac{\partial y}{\partial r} & \frac{\partial y}{\partial \theta} \end{array} \right| = \left| \begin{array}{cc} \cos\theta & -r\sin\theta \\ \sin\theta & r\cos\theta \end{array} \right| = r(\cos^2\theta + \sin^2\theta) = r$$

So, now the integral may be put into the transformed coordinates using this transformation $P = 2\pi$

$$\int_{disk} \frac{1}{\|\vec{x}_{c_i} - \vec{x}'\|} dS' = \int_0^R \int_0^{2\pi} \frac{1}{r} r d\theta dr$$

The integral over the disk is easily seen to be $2\pi R$.

$$\sum_{i=1}^{n} \sum_{i=1}^{n} \frac{x_{c_i}}{\sum_{i=1}^{n} A_{i,i}} = \int_{panel i} \frac{1}{\left\| x_{c_i} - x' \right\|} dS'$$

- 1. If panel is a flat polygon, analytical formulas exist.
- 2. Curved panels can be handled with projection.

7 Summary

Integral Equation Methods Exterior versus interior problems Start with using point sources Standard Solution Methods Collocation Method Galerkin Method Integrals for 3D Problems Singular Integrals

1.2 Numerical Quadrature

Numerical Methods for PDEs

Methods, Lecture 2 Numerical Quadrature

Notes by L. Proctor, S. De and J. White

November 26, 2008

1 Outline

Gaussian Quadrature

Convergence properties Essential role of orthogonal polynomials Multidimensional Integrals

Techniques for singular kernels

Adaptation and variable transformation Singular quadrature.

2 Introduction

Numerical Quadrature is employed as an approximation used to evaluate integrals. We seek an appropriate numerical procedure applied to a definite integral, $I\{f\} \equiv \int_a^b f(x)dx$, where the approximation is essentially of the form $I_n\{f\} \equiv \sum_{i=1}^n \alpha_i f(x_i)$. The *n* distinct points, x_i are the quadrature nodes we have chosen and the quadrature coefficients, or weights, are the α_j terms. In general, we would like to have the smallest possible quadrature error, $E_n\{f\} \equiv$ $I\{f\} - I_n\{f\}$.

2.1 Simple Quadrature Example



To simplify notation, consider the more generic problem of developing a good numerical technique for evaluating the integral of a function f(x) on the domain [0, 1]. We assume that the integrand is a "smooth" function, though we will examine this assumption later. First we have developed a naive approach for obtaining a good approximation of the integral, one we call a **simple quadrature scheme**.

The simplest approach is to replace the integral with the product of the interval (in this case one) and the integrand evaluated at a point inside the interval. If the selected evaluation point is the center of the interval, x = 0.5, we call the scheme **midpoint quadrature**.

A midpoint quadrature scheme replaces the area under the curve f(x) by a

rectangle whose height is the function f(x) evaluated at the midpoint x = 0.5. The scheme is exact when f(x) is a constant. However, what is less obvious is that the scheme is exact when f(x) is a line (an affine function of x) as well. The most obvious way of seeing this is by realizing that when f(x) is a straight line, the area under it is a trapezoid. This trapezoid has exactly the same area as the rectangle which this scheme uses to approximate the integral (can you see why?).



▷ **Exercise 1** Suppose endpoint quadrature (in which the area under the curve is replaced by a rectangle whose height is f(x) evaluated at x = 0 or x = 1) is used instead of midpoint quadrature. For what class of functions is endpoint quadrature exact?

2.2 Improving the Accuracy



One way of improving the midpoint quadrature scheme is to divide the interval [0, 1] into subintervals [0, 0.5] and [0.5, 1], then write the integral

$$\int_0^1 f(x)dx = \int_0^{0.5} f(x)dx + \int_{0.5}^1 f(x)dx,$$

and finally apply a midpoint rule to the integral in each subinterval. We obtain the scheme shown in the slide. The factor $\frac{1}{2}$ appearing in front of $f(\frac{1}{2})$ and $f(\frac{3}{4})$ are just the domain lengths.

 \triangleright **Exercise 1** Can you come up with an expression for the error in this case? How much does the accuracy improve?

Dividing the interval into two reduces the error, now consider using n subintervals and repeating the midpoint quadrature rule on each subinterval. We obtain the scheme

$$\int_{0}^{1} f(x)dx = \sum_{i=1}^{n} \underbrace{\frac{1}{n}}_{\text{subinterval}} f(x_{c_{i}})$$

where the centroid of the i^{th} subinterval is $x_{c_i} = \frac{1}{2}(\frac{i-1}{n} + \frac{i}{n}) = \frac{i-\frac{1}{2}}{n}$. There is no doubt that the accuracy improves, but the key question is by how much? How does this error decrease with the number of subintervals used? And finally, are there clever ways of obtaining better accuracy with less effort?

2.3 General *n*-point formula



Key questions about the method: How fast do the errors decay with n? Are there better methods?

2.4 Different Geometric Approximations

Which geometry is the most accurate?



3 General Quadrature Formula

3.1 General 1-D Formula

\int_0^1	$f(x)dx \simeq \sum_{i=1}^{n}$	$\underbrace{w_i}_{}$	$\underbrace{f(x_i)}$
	<i>i</i> _1	Weight	Evaluation
			\mathbf{Point}

Free to pick the **evaluation points**. Free to pick the **weights** for each point.

An n-point formula has 2n degrees of freedom!

After all the hard work we did dividing the domain into subintervals, we realize that we cannot even integrate a quadratic function exactly on the domain. There must be something that we can do to improve this scheme. We go back and look at the general form of the quadrature approximation scheme. All we are doing is approximating an integral by a weighted sum of function evaluations. So far we have been choosing these weights as the subinterval lengths. We have also been choosing the evaluation points as the center of the interval, in the midpoint quadrature scheme. The weights are just some normalizing factors which we have taken to be the fraction of the domain over which we are evaluating. The equality of areas of trapezoids and rectangles that we previously discussed gives us the extra polynomial accuracy of being able to obtain the area under a straight line exactly. So, what would happen if we were to choose both the integration points and the weights intelligently? For an *n*-point formula we have n weights and n evaluation points to choose. That gives us 2n degrees of freedom. Hence we must be able to exactly integrate a polynomial of degree at most (2n-1). This idea gives rise to the **Gaussian quadrature** scheme.

3.2 Evaluation Points & Weights Selection

Can make the result exact if f(x) is a polynomial

$$f(x) = a_0 + a_1 x + a_2 x^2 + \dots + a_l x^l = p_l(x)$$

Select x_i 's and w_i 's such that

$$\int_0^1 p_l(x) dx = \sum_{i=1}^n w_i p_l(x_i)$$

for ANY polynomial up to (and including) l^{th} order With 2n degrees of freedom, l = 2n - 1 Let $p_l(x)$ denote a polynomial of degree l in the variable x ($a_l \neq 0$). We want to select the weights and integration points such that the formula

$$\int_0^1 f(x) dx = \int_0^1 p_l(x) dx = \sum_{i=1}^n w_i p_l(x_i)$$

is exact for all polynomials of degree up to (and including) l. Obviously, with 2n degrees of freedom, the best we can do is l = 2n - 1.

Note 1

Example: Third Order Polynomial

As an example, consider integratin the function $f(x) = x - x^3$. from x = 0 to x = 1. The exact solution is, $I\{f\} = \int_0^1 (x - x^3) dx = \frac{1}{4}$. It is stated above that since we have a polynomial of degree three (l = 3), then we will be able to find an exact solution using two point quadrature (n = 2). But, one must note, that finding the exact solution is not simply a matter of applying midpoint quadrature, as we have done previously. The solution using midpoint quadrature, the method of which will be studied in further detail later on, will provide the exact solution.



Assuming the weights, w_i , remain bounded, and the derivatives of f(x) are bounded on [0, 1],

$$\left|\int_0^1 f(x)dx - \sum_{i=1}^n w_i f(x_i)\right| \le \frac{K}{(2n)!}$$

Gaussian quadrature converges very quickly!!

4 Error Analysis

4.1 Taylor Series Expansion

To derive the error of the midpoint quadrature scheme analytically, consider the interval [0, h], h > 0 and Taylor expand f(x) about the center of this interval, $\bar{x} = \frac{h}{2}$,

$$f(x) = f(\bar{x}) + \Delta(x)\frac{df(\bar{x})}{dx} + \frac{\Delta(x)^2}{2!}\frac{d^2f(\xi)}{dx^2} \quad for \ some \ \xi \in [0,h],$$

where $\Delta(x) = x - \bar{x}$. The last term in the expansion is the Taylor series remainder. Integrating this expansion over the interval [0, h]

$$\int_0^h f(x) dx = hf(\bar{x}) + \frac{h^3}{24} \frac{d^2 f(\xi)}{dx^2}.$$

Hence the error in the midpoint quadrature approximation is

$$E = \int_0^h f(x) dx - hf(\bar{x}) = rac{h^3}{24} rac{d^2 f(\xi)}{dx^2}.$$

A function that is first-order polynomial in x would have zero as the second derivative, and therefore the above expression tells us that the error of midpoint quadrature for such functions is identically zero. In addition, the above expression tells us that the error scales as the cube of the domain length.

 \triangleright **Exercise 2** If the midpoint quadrature scheme uses rectangular geometry to approximate the area under the curve, then why is the first derivative needed in the Taylor Series expansion? How does one represent the expansion of the trapezoidal approximation?

4.2 Example - Error vs. n





Note 2

Above is the example of integrating f(x) = sin(x) on the domain [0, 1]. We obtain progressively better answers to the integration by increasing the number of subintervals n. The error in evaluating the integral is plotted as a function of the number of subintervals (n) above. The error appears to be going down as $\mathcal{O}\left(\frac{1}{n^2}\right)$.

From what we have just seen, the error inside the i^{th} subinterval (of length $h = \frac{1}{n}$) is $\frac{h^3}{24} \frac{d^2 f(\xi_i)}{dx^2}$ for some $\xi_i \in [\frac{i-1}{n}, \frac{i}{n}]$. Hence, for the entire interval [0, 1] we can sum these errors and obtain the error, E_n for an approximation using n subintervals as

$$E_n = \frac{nh^3}{24} \underbrace{\left(\frac{1}{n} \sum_{i=1}^n \frac{d^2 f(\xi_i)}{dx^2}\right)}_M$$

It is easy to see that if f(x) is a continuous function, M (being the mean) must be bounded by the maximum and the minimum of f(x) on the interval [0,1]and hence, there must exist some $\xi \in [0,1]$ such that $M = d^2 f(\xi)/dx^2$. Hence we obtain the estimate

$$E_n = \frac{nh^3}{24} \frac{d^2 f(\xi)}{dx^2} = \frac{1}{24n^2} \frac{d^2 f(\xi)}{dx^2}$$

since h = 1/n. This error estimate tells us that the scheme is again exact for constants and linear functions on the domain (no higher order polynomials!) and, for a smooth function, the error decays **algebraically**.

4.3 The Exactness Criteria

Consider the Taylor series for f(x) expanded about x = 0

$$f(x) = f(0) + \frac{\partial f(0)}{\partial x}x + \dots + \frac{1}{l!}\frac{\partial^l f(0)}{\partial x^l}x^l + R_{l+1}$$

 R_{l+1} is the **remainder**

$$R_{l+1} = \frac{1}{(l+1)!} \frac{\partial^{l+1} f(\tilde{x})}{\partial x^{l+1}} x^{l+1}$$

where $\tilde{x} \in [0, x]$

Note 3

Of all functions, why are we interested in integrating polynomials? The reason comes from the structure of Taylor's series expansion. The Taylor expansion of a function in a local neighborhood of a point (here this point is chosen as 0 without loss of generality) is nothing but a power series expansion. The higher

the order of polynomials that our scheme can integrate means a higher order of the remainder term in the expansion. The integral of the remainder over the domain is precisely the error in numerical integration.

The exactness condition requires

$$\int_{0}^{1} p_{l}(x)dx = \int_{0}^{1} (a_{0} + a_{1}x + a_{2}x^{2} + \dots + a_{l}x^{l})dx = \sum_{i=1}^{n} w_{i}p_{l}(x_{i})$$

for any set of l + 1 coefficients a_0, a_1, \ldots, a_l

Equivalently
$$\int_0^1 a_0 dx + \int_0^1 a_1 x dx + \int_0^1 a_2 x^2 dx + \dots + \int_0^1 a_l x^l dx = \sum_{i=1}^n w_i p_l(x_i)$$

This slide needs little clarification. Our exactness criterion is

$$\int_0^1 p_l(x)dx = \int_0^1 (a_0 + a_1x + a_2x^2 + \dots + a_lx^l)dx = \sum_{i=1}^n w_i p_l(x_i)$$

which is the same as

$$a_0 \int_0^1 dx + a_1 \int_0^1 x dx + \dots + a_l \int_0^1 x^l dx = a_0 \sum_{i=1}^n w_i + a_1 \sum_{i=1}^n w_i x_i + \dots + a_l \sum_{i=1}^n w_i x_i^l$$

For this to be an identity for the (l + 1) arbitrary coefficients a_i , we must have the (l + 1) conditions

$$\sum_{i=1}^{n} w_i x_i^{\ j} = \int_0^1 x^j dx \quad for \ j = 0, 1, \dots, l$$

Using the Taylor series results, the exactness criteria, and the innate linearity of the quadrature scheme r^1

$$\int_{0}^{1} f(x)dx - \sum_{i=1}^{n} w_{i}f(x_{i}) = \underbrace{\frac{1}{(l+1)!} \frac{\partial^{l+1}f(\tilde{x})}{\partial x^{l+1}} \left(\int_{0}^{1} x^{l+1}dx - \sum_{i=1}^{n} w_{i}x_{i}^{l+1}\right)}_{Remainder}$$

Exactness condition will be satisfied if and only if

$$\int_0^1 dx = \sum_{i=1}^n w_i \cdot 1$$
$$\int_0^1 x dx = \sum_{i=1}^n w_i \cdot x_i$$
$$\vdots$$
$$\int_0^1 x^l dx = \sum_{i=1}^n w_i \cdot x_i^l$$

Reorganizing exactness equations

$$\begin{bmatrix} 1 & 1 & \cdots & 1 \\ x_1 & x_2 & \cdots & x_n \\ \vdots & \vdots & \ddots & \vdots \\ x_1^{l} & x_2^{l} & \cdots & x_n^{l} \end{bmatrix} \begin{bmatrix} w_1 \\ w_2 \\ \vdots \\ w_n \end{bmatrix} - \begin{bmatrix} 1 \\ \vdots \\ \vdots \\ 0 \\ \vdots \\ 0 \end{bmatrix} = 0$$

Nonlinear, since x_i 's and w_i 's are unknowns

What is a practical way of computing the evaluation points and weights? The system of equations is not easy to solve since x_i 's and w_i 's are unknowns.

5 Computing the Points & Weights

5.1 Newton's Method

Could use Newton's Method

$$F(y) = 0 \Rightarrow J_F(y^k)(y^{k+1} - y^k) = -F(y^k)$$

The nonlinear function for Newton is then

$$F\binom{w}{x} = \begin{bmatrix} 1 & 1 & \cdots & 1 \\ x_1 & x_2 & \cdots & x_n \\ \vdots & \vdots & \ddots & \vdots \\ x_1^l & x_2^l & \cdots & x_n^l \end{bmatrix} \begin{bmatrix} w_1 \\ w_2 \\ \vdots \\ w_n \end{bmatrix} - \begin{bmatrix} 1 \\ \vdots \\ \vdots \\ 0 \\ \end{bmatrix} = 0$$

Note 4

Newton's method is an iterative technique for finding a value y such that F(y) = 0. The method is based on linearizing the problem about a guess at y, and then

updating the value of y by solving the linearized problem. In particular, the iterate y^{k+1} is determined from y^k by solving the linear system of equations

$$F(y^{k}) + J_{F}(y^{k}) \left(y^{k+1} - y^{k} \right) = 0$$

where $J_F(y^k)$ is the Jacobian (multidimensional derivative) of the nonlinear function F(y). The iteration is continued until the updated y is sufficiently close to the exact solution, a criterion that can be difficult to verify. Newton's method does not always converge, a phenomenon that is more likely when $J_F(y)$ is nearly singular. For more about Newton's method, see the 6.336/16.910/2.096 course notes (available under open courseware).



Newton Method Jacobian reveals the problem

-				2n				
[1	1		1	0	0	••••	0	
$J_{x} \begin{pmatrix} w \\ w \end{pmatrix} = \begin{vmatrix} x_{1} \\ x_{2} \end{vmatrix}$	x_2	•••	x_n	w_1	W_2	•••	W _n	
$\left \begin{array}{c} x \\ x \end{array} \right =$	÷	۰.	÷	÷	÷	٠.	:	
x_1'	x_2'		x'_n	$lw_1x_1^{l-1}$			$lw_1 x_n^{l-1}$	
Columns become linearly dependent for high order								

Looking at the Jacobian of the problem, we realize that the first n columns become increasingly linearly dependent for large l. This is bound to happen since we are looking into the space $span\{1, x, \ldots, x^l\}$. This basis always becomes ill-conditioned with increasing l. The solution is to obtain a polynomial basis that is "normalized" in some sense so that it is properly conditioned.

5.2 Orthogonal Polynomials

5.2.1 Introduction



Exactness criteria will be satisfied if and only if

$\int_{0}^{1} c_0(x) dx = \sum_{i=1}^{n} w_i c_0(x_i)$	BUT
$\int_{0}^{1} c_{1}(x) dx = \sum_{i=1}^{n} w_{i} c_{1}(x_{i})$	Each c_i polynomial must Contain an x^i term Be linearly independent
$\int_{0}^{1} c_{i}(x) dx = \sum_{i=1}^{n} w_{i} c_{i}(x_{i})$	

The only difference from the previous set of criteria is that these polynomials have better properties than the ones we chose before.

5.2.2 Orthogonality

For the normalized integral, two polynomials are said to be **orthogonal** if

$$\int_0^1 c_i(x)c_j(x)dx = 0 \quad for \ j \neq i$$

The above integral is often referred to as an inner product and ascribed the notation

$$(c_i, c_j) = \int_0^1 c_i(x) c_j(x) dx$$

The connection between polynomial inner products and vector inner products can be seen by sampling.

5.2.3 Exactness Criteria

Consider rewriting the exactness criteria



Recall that l = 2n - 1where l = degree of polynomial & n = number of coefficients

Call the first (n-1) conditions the "lower order terms" and the last n conditions the "higher order terms."

5.2.4 Higher Order Terms Contain Lower Order Terms

Write the higher order terms differently

$$\int_{0}^{1} c_{n}(x) dx = \sum_{i=1}^{n} w_{i} c_{n}(x_{i}) \quad \Rightarrow \quad \int_{0}^{1} c_{n}(x) c_{0}(x) dx = \sum_{i=1}^{n} w_{i} c_{n}(x_{i}) c_{0}(x_{i})$$

$$\vdots$$

$$\int_{0}^{1} c_{2n-1}(x) dx = \sum_{i=1}^{n} w_{i} c_{2n-1}(x_{i}) \quad \Rightarrow \quad \int_{0}^{1} c_{n}(x) c_{n-1}(x) dx = \sum_{i=1}^{n} w_{i} c_{n}(x_{i}) c_{n-1}(x_{i})$$

The products $c_n(x)c_j(x)$ are linearly independent.

In this slide we express the "higher order terms" as conditions involving "lower order terms."

5.2.5 Using Orthogonality and Roots

Use orthogonal polynomials



Pick the x_i 's to be *n* roots of $c_n(x)$ The higher order constraints are exactly satisfied!

This elegant step relies on polynomial orthogonality. If we choose the polynomial $c_n(x)$ such that it is orthogonal to all polynomials of inferior degree (i.e. $c_0(x), c_1(x), \ldots, c_{n-1}(x)$) and the x_i 's are roots of this polynomial, then the higher order n conditions are automatically satisfied. Note that for this derivation we used polynomials which are orthogonal on the interval [0, 1]. Such polynomials are shifted and scaled versions of the classical Legendre polynomials, which are orthogonal on the interval [-1, 1].

5.2.6 Satisfying Lower Order Constraints

An abbreviated exactness equation

$$\begin{bmatrix} c_0(x_1) & \cdots & c_0(x_n) \\ \vdots & \ddots & \vdots \\ c_{n-1}(x_1) & \cdots & c_{n-1}(x_n) \end{bmatrix} \begin{bmatrix} w_1 \\ \vdots \\ w_n \end{bmatrix} = \begin{bmatrix} \int_0^1 c_0(x) dx \\ \vdots \\ \int_0^1 c_{n-1}(x) dx \end{bmatrix}$$

Now linear, x_i 's are known

Rows are sampled orthogonal polynomials.

By using the roots of $c_n(x)$ for the x_i 's, the higher order constraints are automatically satisfied. Since the x_i 's are now known, only the weights are still unknown. The lower *n* constraints can be used to determine the weights, generating a linear system.

6 Gaussian Quadrature Algorithm

1. Construct n + 1 orthogonal polynomials

$$\int_0^1 c_i(x)c_j(x)dx = 0 \quad for \ j \neq i$$

2. Compute *n* **roots**, x_i , i = 1, ..., n of the n^{th} order orthogonal polynomial such that $c_n(x_i) = 0$

3. Solve a linear system for the weights w_i

4. Approximate the integral as a sum

$$\int_{0}^{1} f(x) dx = \sum_{i=1}^{n} w_{i} f(x_{i})$$

6.1Example

Note 5

Example: Third Order Polynomial

Recall the example of the third-order polynomial, $f(x) = x - x^3$. We would like to determine w_1, w_2, x_1 , and x_2 so that the integration formula,

$$I_2\{f\} = w_1 f(x_1) + c_2 f(x_2)$$

gives the exact result.

Generalize this problem to any third-order polynomial of the form

$$f(x) = a_0 + a_1 x + a_2 x^2 + a_3 x^3.$$

Given the exactness condition, the integral can be written:

$$\int_0^1 (a_0 + a_1x + a_2x^2 + a_3x^3) dx = w_1(a_0 + a_1x_1 + a_2x_1^2 + a_3x_1^3) + w_2(a_0 + a_1x_2 + a_2x_2^2 + a_3x_2^3)$$

Gather the like terms: a_0 : $w_1 + w_2 = \int_0^1 dx = 1$

 $\begin{array}{rcl} a_{0} & & & w_{1} + w_{2} - j_{0} & dx - 1 \\ a_{1} & & & w_{1}x_{1} + w_{2}x_{2} = \int_{0}^{1} x dx = \frac{1}{2} \\ a_{2} & & & w_{1}x_{1}^{2} + w_{2}x_{2}^{2} = \int_{0}^{1} x^{2} dx = \frac{1}{3} \\ a_{3} & & & w_{1}x_{1}^{3} + w_{2}x_{2}^{3} = \int_{0}^{1} x^{3} dx = \frac{1}{4} \\ \end{array}$ There are four equations above and four unknowns, so the system can be easily solved to give, $w_{1} = \frac{1}{2} \quad w_{2} = \frac{1}{2} \quad x_{1} = \frac{3-\sqrt{3}}{6} \quad x_{2} = \frac{3+\sqrt{3}}{6}.$ Plugging in these values, for the equation above, gives the exact solution of 0.25.



6.2 Summary

This slide summarizes the technique of finding weights and integration points for Gauss quadrature.

6.2.1 Accuracy Result

$$\int_0^1 f(x)dx \simeq \sum_{i=1}^n w_i f(x_i)$$

Key properties of the method

- An *n*-point Gauss quadrature rule is **exact** for polynomials of order 2n-1
- Error is proportional to $\frac{1}{(2n)!}$ (like $\frac{1}{n^n}$)

6.2.2 Simple vs. Gauss Quadrature



Key property of the method

• Error is proportional to $\frac{1}{n^2}$

 \triangleright **Exercise 3** Do you see that the simple quadrature scheme is a special case of Gauss quadrature? \blacksquare



Notice that for a smooth function $f(x) = cos(2\pi x)$, which is infinitely differentiable, Gauss quadrature far outperforms the simple quadrature scheme

6.2.3 Evaluation Point Placement

Notice the clustering at interval ends

In the Gauss quadrature scheme the evaluation points are roots of Legendre polynomials which are clustered at the ends of the interval.

7 The Singular Kernel Problem

7.1 Calculating the "Self-Term"

Now lets go back to our problem of solving Laplace's equation on a 3D domain using boundary integral representation. We realize that we now have some sophisticated tools to handle integrals of functions that are smooth. But what about the integral on the panel where the centroid x_{c_i} is located? The Green's function blows up at the centroid. However, the function is integrable because the integrand blows up at a rate that is slower than the rate at which the surface measure goes to zero in the vicinity of the singularity. So we know that the integral exists and is finite, but is Gauss quadrature capable of performing well in the presence of this singularity?



7.2 Symmetric 1-D Example

In 1D we look at a function $f(x) = \frac{1}{\sqrt{|x|}}$ which is integrable on [-1, 1] but has a singularity at x = 0.



7.2.1 Integrable and Nonintegrable Singularities



7.2.2 Comparing Quadrature Schemes



We observe that Gauss quadrature is not very good at integrating this function. The convergence is rather poor. As a matter of fact, it is more inaccurate than the simple quadrature scheme. In the next few slides we present several techniques of handling integrals with singularities (which are integrable, of course)

- Subinterval (adaptive) quadrature
- Change of variables of integration
- Singular (Gaussian) quadrature

7.3 Improved Techniques

7.3.1 Subinterval (Adaptive) Quadrature



Subdivide the integration interval $\int_{-1}^{1} \frac{1}{\sqrt{|x|}} dx = \int_{-1}^{-0.1} \frac{1}{\sqrt{|x|}} dx + \int_{-0.1}^{0} \frac{1}{\sqrt{|x|}} dx + \int_{0}^{0.1} \frac{1}{\sqrt{|x|}} dx + \int_{0.1}^{1} \frac{1}{\sqrt{|x|}} dx$ Use Gauss quadrature in each subinterval Polynomials fit subintervals better Expensive if many subintervals used.

7.3.2 Change of Variables - for Simple Cases

Change variables to eliminate singularity

 $y^2 = x \qquad \Rightarrow \qquad 2ydy = dx$

$$\int_{-1}^{1} \frac{1}{\sqrt{|x|}} dx = 2 \int_{0}^{1} \frac{1}{\sqrt{|y^2|}} 2y dy = 2 \int_{0}^{1} 2dy$$

Apply Gauss quadrature on desingularized integrand.

7.3.3 Singular Quadrature - Complicated Cases

Basic Concept

Integrand has known singularity s(x)

$$\int_{-1}^{1} f(x)s(x)dx \text{ where } f(x) \text{ is smooth}$$

Develop a quadrature formula exact for

$$\int_{-1}^{1} p_l(x)s(x)dx$$
 where $p_l(x)$ is polynomial of order l

Calculate weights like Gauss quadrature

Note 6

It is possible to generate Gaussian quadrature schemes of the form

$$\int_{-1}^1 s(x)f(x)dx = \sum_{i=1}^n w_i f(x_i)$$

for functions which have a known singularity s(x) > 0. The quadrature formula needs to be exact when f(x) is a polynomial of order at most l. Not surprisingly, it turns out that the integration points are the n roots of a polynomial $c_n(x)$ of degree n = (l+1)/2 which is orthogonal to all polynomials of inferior degree with respect to the weight s(x), i.e.

$$\int_{-1}^{1} s(x)c_n(x)c_j(x) = 0 \qquad for \ j = 0, 1, \dots, (n-1).$$

An example is the singular integral

$$I = \int_{-1}^{1} \frac{f(x)}{\sqrt{1 - x^2}}$$

Here, $s(x) = 1/\sqrt{1-x^2}$ and the corresponding orthogonal polynomials turn out to be the Chebyshev polynomials. The integration points are given in closed form by

$$x_i = \cos\left(\pi \frac{2i-1}{2n}\right)$$

and the corresponding weights are $w_i = \pi/n$.

Singular Quadrature Weights

$$\begin{bmatrix} c_0(x_1) & \cdots & c_0(x_n) \\ \vdots & \ddots & \vdots \\ c_{n-1}(x_1) & \cdots & c_{n-1}(x_n) \end{bmatrix} \begin{bmatrix} w_1 \\ \vdots \\ w_n \end{bmatrix} = \begin{bmatrix} \int_{-1}^1 c_0(x)s(x)dx \\ \vdots \\ \int_{-1}^1 c_{n-1}(x)s(x)dx \end{bmatrix}$$

Need (analytic) formulas for integrals of c(x)s(x)

The lower order constraints can be used to compute the integration weights.

8 Summary

Easy technique for computing integrals Piecewise constant approach Gaussian quadrature Faster convergence Essential role of orthogonal polynomials **Techniques for singular kernels** Adaptation and Variable Transformation Singular quadrature **What about multiple dimensions?**

1.3 First and Second Kind Equations

Numerical Methods for PDEs

Integral Equation Methods, Lecture 3 First and Second Kind Equations

Notes by L. Proctor, S. De and J. White

December 1, 2008

1 Outline

Convergence issues in 1D

First and second kind integral equations Develop some intuition about the difficulties **Convergence for second kind equations** Consistency and stability issues **Nystrom Methods** High order convergence

2 Example Problems in 1D

2.1 First Kind Equation

 $\Psi(x) = \int_{-1}^{1} |x - x'| \sigma(x') dS' \qquad x \in [-1, 1]$ The potential is given $\Psi(x) = x^3 - x \qquad \text{The density must be computed}$ $\Psi_{x}^{*} \qquad \sigma(x) \text{ is unknown}$

In the next several slides we will investigate the convergence properties of integral equation discretization methods. How these methods converge depends on what kind of integral equation is being solving. Examining this issue will introduce one of the subtle points about integral equations.

To begin, consider the example one-dimensional first-kind integral equation on the above slide. For this equation, we assume that the potential, $\Phi(x)$, is known and that the charge density $\sigma(x)$ is unknown. Here, x is in the interval [-1, 1], and the integration is over that same interval. Note that for this example, the Green's function is given by G(x, x') = |x - x'|.

In the left plot below the equation, an example potential, $x^3 - x$, is given and plotted as a function of x. On the right is a plot of charge density as a function of x which is a possible solution to the integral equation. As we will see shortly, the solution for the charge density is not so easy to find.

2.1.1 Collocation Discretization

$$\Psi(x) = \int_{-1}^{1} |x - x'| \sigma(x') dS' \qquad x \in [-1, 1]$$
(1)

Piecewise-Constant Centroid-Collocation

$$\Psi(x_{c_i}) = \sum_{j=1}^{n} \sigma_{nj} \int_{x_{j-1}}^{x_j} |x_{c_i} - x'| dS'$$
(2)

To compute the numerical solution to this one-dimensional problem, consider solving integral equation (1) using a piecewise-constant collocation scheme. In such a scheme, we first select n + 1 points on the interval, in this case [-1, 1]. We denote those points as $\{x_0, x_1, ..., x_n\}$, as shown in the above figure. For this example, $x_0 = -1$ and $x_n = 1$. Over the subintervals define a set of basis functions, $\{\varphi_1(x), \varphi_2(x), ..., \varphi_n(x)\}$, where

$$\varphi_i(x) = \left\{ \begin{array}{cc} 1 & x \in [x_{i-1}, x_i] \\ 0 & \text{otherwise} \end{array} \right\}.$$
(3)

The charge density σ can then be represented approximately as

$$\sigma(x) \approx \sigma_n(x) \equiv \sum_{i=1}^n \sigma_{ni} \varphi_i(x), \tag{4}$$

where σ_{ni} is the weight associated with the i^{th} basis function. It may seem odd that we used the same letter to represent the density and the basis function weights, but there is a reason. The above basis set is such that only one basis function is nonzero for a given x, and basis functions only take on the value zero or one. Therefore, σ_{ni} will be equal to the approximate charge density when $x \in [x_{i-1}, x_i]$.

Charge Density Representation



Plugging the basis function representation of the charge density, equation (4), into equation (1) yields

$$\Psi(x) = \int_{-1}^{1} |x - x'| \sum_{i=1}^{n} \sigma_{ni} \varphi_i(x') dS',$$

which can be simplified by exploiting equation (3). Recall that the residual, R(x), is defined as how well the weighted combination of basis functions satisfies the integral equation. In this centroid collocation case,

$$R(x) = \Psi(x) - \sum_{j=1}^{n} \sigma_{nj} \int_{x_{j-1}}^{x_j} |x - x'| dS'.$$

If collocation is used to solve this equation, then $R(x_{c_i}) = 0$ for all *i*, where x_{c_i} is the *i*th collocation point. The collocation points shown in the figure are the subinterval center points, $x_{c_i} = \frac{1}{2}(x_{i-1} + x_i)$. Note that there are other choices for collocation points, such as $x_{c_i} = x_i$. Using the fact that $R(x_{c_i}) = 0$ leads to

$$R(x_{c_i}) = \Psi(x_{c_i}) - \sum_{j=1}^n \sigma_{nj} \int_{x_{j-1}}^{x_j} |x_{c_i} - x'| dS' = 0$$

which can be reorganized into equation (2).

The Matrix

We can now generate a system of equations that can be used to solve for the σ_{ni} 's, the piecewise constant charge densities for each of the subintervals.



One row for each collocation point

The right-hand side of this system of equations is a vector of known potentials at the interval centers (the collocation points). The i^{th} row of the matrix corresponds to unfolding the sum in the collocation equation

$$\Psi(x_{c_i}) = \sum_{j=1}^n \sigma_{nj} \int_{x_{j-1}}^{x_j} |x_{c_i} - x'| dS',$$

and the entries in the j^{th} column correspond to how much the charge on the j^{th} interval contributes to the i^{th} potential. Note that the matrix is square and dense.

 \triangleright **Exercise 1** Is the above matrix symmetric? If we used $x_{c_i} = x_i$, would the matrix still be symmetric?

2.1.2 Numerical Results with Increasing n

One usually believes that a discretization scheme should produce progressively more accurate answers as the discretization is refined. In this case, as we divide the interval into progressively finer subintervals, one might expect that the piece-

wise constant representation of the charge density given by $\sigma_n(x) \approx \sum_{i=1}^n \sigma_{ni} \varphi_i(x)$

would become more accurate as n increases.



As is clear from the above plot, the piecewise-constant centroid-collocation discretization of (1) is not converging. In the plot, which is hard to decipher without looking at a color version, the σ_{ni} 's produced using n = 10, n = 20 and n = 40 subintervals are shown. For each discretization, a point is plotted at σ_{ni} , x_i for i = 1, ..., n, so there are ten points plotted for the coarsest discretization and forty points plotted for the finest discretization, but all sets of points span the interval $x \in [-1, 1]$.

What is clear from comparing the blue points (n = 10) to the red points (n = 20) and to the green points (n = 40), is that the charge density seems to be approaching infinity as the discretization is refined. The results are certainly not converging.

Why is this happening? Is the numerical technique at fault, or is the integral equation a problem?

2.2 Second Kind Equation

We are going to postpone examining what went wrong with the discretization the first-kind integral equation, and instead examine a Second Kind integral equation example. As in the first-kind case, we are assuming the potential, $\Psi(x)$, is known and that the charge density $\sigma(x)$ is unknown. Also, x is in the interval [-1, 1], and the integration is over that same interval. Once again, the Green's function is given by G(x, x') = |x - x'|.

$$\Psi(x) = \sigma(x) + \int_{-1}^{1} |x - x'| \sigma(x') dS' \qquad x \in [-1, 1]$$



The above equation is second-kind instead of first-kind because the unknown charge density appears both inside *and* outside of the integral. In the first-kind equation, the density appeared only inside the integral. This seemingly small difference has enormous numerical ramifications.

In the left plot above, an example given potential, $\Psi(x) = x^3 - x$ is plotted as a function of x. On the right is a plot of a charge density as a function of x which satisfies this second kind integral equation. As we will see below, this equation is easily solved numerically.

2.2.1 Collocation Discretization

$$\Psi(x) = \sigma(x) + \int_{-1}^{1} |x - x'| \sigma(x') dS' \quad x \in [-1, 1]$$
(5)

Centroid Collocated Piecewise Constant Scheme

$$\Psi(x_{c_i}) = \sigma_{ni} + \sum_{j=1}^{n} \sigma_{nj} \int_{x_{j-1}}^{x_j} |x_{c_i} - x'| dS'$$
(6)

To compute the numerical solution to the one-dimensional second-kind equation (5), once again consider using a piecewise-constant collocation scheme. Again, we select n+1 points on the interval and denote those points as $\{x_0, x_1, ..., x_n\}$, as shown in the figure above. For this example, $x_0 = -1$ and $x_n = 1$. The corresponding basis functions, $\{\varphi_1(x), \varphi_2(x), ..., \varphi_n(x)\}$, are the same as in equation (3):

$$\varphi_i(x) = \left\{ \begin{array}{cc} 1 & x \in [x_{i-1}, x_i] \\ 0 & \text{otherwise} \end{array} \right\}.$$
(7)

The charge density σ is approximately represented by

$$\sigma(x) \simeq \sigma_n(x) \equiv \sum_{i=1}^n \sigma_{ni} \varphi_i(x), \tag{8}$$

where σ_{ni} is the weight associated with the i^{th} basis function. Plugging the basis function representation of the charge density, equation (8) into the second kind integral equation (5) gives:

$$\Psi(x) = \sum_{j=1}^{n} \sigma_{nj} \varphi_j(x) + \int_{-1}^{1} |x - x'| \sum_{i=1}^{n} \sigma_{ni} \varphi_i(x') dS',$$

which can be simplified by exploiting the specific basis functions, equation (7) to

$$\Psi(x) = \sum_{j=1}^{n} \sigma_{nj} \varphi_j(x) + \sum_{j=1}^{n} \sigma_{nj} \int_{x_{j-1}}^{x_j} |x - x'| dS'.$$
(9)

As shown in the above figure, the collocation points are the subinterval center points, $x_{c_i} = \frac{1}{2}(x_{i-1} + x_i)$. When collocation is used, equation (9) must be satisfied exactly at the collocation points and therefore

$$\Psi(x_{c_i}) = \sum_{i=1}^n \sigma_{ni} \varphi_j(x_{c_i}) + \sum_{j=1}^n \sigma_{nj} \int_{x_{j-1}}^{x_j} |x_{c_i} - x'| dS'.$$
(10)

Note that $\varphi_j(x_{c_i}) = 0$ when $i \neq j$, and $\varphi_i(x_{c_i}) = 1$. Using this fact yields equation (6).

The Matrix

Just as in the discretized first-kind equation, we generate a system of equations that can be used to solve for the σ_{ni} 's, the piecewise constant charge densities for each of the subintervals.

$$\begin{bmatrix} 1+ \sum_{x_0}^{x_1} |x_{c_1} - x'| dS' & \cdots & \int_{x_{n-1}}^{x_n} |x_{c_1} - x'| dS' \\ \vdots & \ddots & \vdots \\ \int_{x_0}^{x_1} |x_{c_n} - x'| dS' & \cdots & \boxed{1 + \int_{x_{n-1}}^{x_n} |x_{c_n} - x'| dS'} \end{bmatrix} \begin{bmatrix} \sigma_{n1} \\ \vdots \\ \sigma_{nn} \end{bmatrix} = \begin{bmatrix} \Psi(x_{c_1}) \\ \vdots \\ \Psi(x_{c_n}) \end{bmatrix}$$

The right-hand side of this system of equations is a vector of known potentials at interval centers (the collocation points). The i^{th} row of the matrix corresponds to unfolding the sum in the collocation equation

$$\Psi(x_{c_i}) = \sigma_{ni} + \sum_{j=1}^n \sigma_{nj} \int_{x_{j-1}}^{x_j} |x_{c_i} - x'| dS'$$

and the entries in the j^{th} column corresponds to how much the charge on the j^{th} interval contributes to the i^{th} potential.

The major difference between the matrix in this discretized second-kind example and the first-kind example is circled on the above slide. There is an additional one on the diagonal of the discretized second-kind equation that did not appear in the first-kind equation. In other words,

$$A_{second \ kind} = I + A_{first \ kind}.$$

2.2.2 Numerical Results with Increasing n

Unlike the results from discretizing the first kind equation, progressively refining the discretization of the second kind equation produces more accurate answers.



Answers Are Improving!!!

Once again, the plot is a little hard to decipher without looking at the color version. It shows the σ_{ni} 's produced using n = 10, n = 20 and n = 40 subintervals. For each discretization, a point is plotted at σ_{ni} , x_i for i = 1, ..., n, so there are ten points plotted for the coarsest discretization and forty points plotted for the finest discretization, but all sets of points span the interval $x \in [-1, 1]$. What is clear from comparing the blue points (n = 10) to the red points (n = 20) and to the green points (n = 40), is that the charge density seems to be approaching a smooth solution.

What is the essential difference between first and second kind equations? Is it some aspect of the numerical technique or are these two equations really that different? In the next slides, we will try to answer this question.

2.3 Difficulty with the First Kind Equation

We will make use of operator-function notation for much of the next sections, both for clarity and brevity. For example, the charge density, σ , and the potential, Ψ , are functions of the independent variable x. When we mean the function, we use just use the function name, such as σ or Ψ . When we give an explicit formula in terms of x for the function, or are denoting the function's value for a particular \hat{x} , we use follow the function name by the value of the independent variable in parentheses. For example, $\sigma \hat{x}$ is the value of function σ when evaluated at $x = \hat{x}$.

The operator-function notation is a little less obvious in the case of operators that map functions to functions, like the integral operator. The integral operator takes a function, in this case σ , and produces another function that we might refer to as the potential. If we denote the integral operator from (5) as K, then $K\sigma$ is a function. If we wish to evaluate the function generated by applying K to σ at some \hat{x} , then we write $(K\sigma)(x)$. Note that $K\sigma(x)$ would NOT BE CORRECT. The operator K takes functions and $\sigma(x)$ is a value.

2.3.1 Singular Integral Operator

Denote the integral operator as K

$$(K\sigma)(x) \equiv \int_{-1}^{1} |x - x'| \sigma(x') dS' = \Psi(x) \Rightarrow K\sigma = \Psi$$
(11)

The integral operator is singular : K has a null space

$$\int_{-1}^{1} \sigma_0(x) = 0, x \neq 0, \sigma_0(0) = 1$$

$$(K\sigma_0)(x) = \int_{-1}^{1} |x - x'| \sigma_0(x') dS' = 0 \text{ for all } x$$
(12)

$$\Rightarrow K\sigma_0 = 0 \tag{13}$$

If $K\sigma^a = \Psi$ then $K(\sigma^a + \sigma_0) = \Psi$

In equation (11), we introduce the abstract notion that

$$\int_{-1}^1 |x-x'|\sigma(x')dS'$$

is an operator on the function σ , which we denote with the symbol K. As shown on the top of the slide, this notation makes writing the integral equation look just like writing a matrix equation.

The key problem is that the operator K is singular. And if

$$K\sigma = \Psi$$

were a matrix equation with a singular K, one would not be surprised to discover the system of equations is hard, or impossible, to solve.

We will not try, in this lecture, to be formal about the concept of a singular operator. To do so, we would necessarily be examining details about certain types of function spaces. Instead, we will try to develop some intuition. In particular, we will draw an analogy to matrices and note that if an operator is singular, it must have a null space.

To see that K does have a null space, consider the spike function $\sigma_0(x)$ depicted in equation (13). This spike function is one at x = 0 and zero otherwise. Note, this function is *not* an impulse function. Unlike the impulse function, the spike's value at x = 0 is finite and the area under its curve is obviously zero.

As noted in equation (13), $K\sigma_0 = 0$. To see this consider that since σ_0 is nonzero only at x = 0, and therefore

$$\int_{-1}^{1} |x - x'| \sigma_0(x') dS' = |x| \int_{-1}^{1} \sigma_0(x') dS'$$

Since $\int_{-1}^{1} \sigma_0(x') dS' = 0$, as the area under σ_0 's curve is zero, then $K\sigma_0 = 0$. The statement "If $K\sigma^a = \Psi$ then $K(\sigma^a + \sigma_0) = \Psi$ " says that if K has a null space, and there exists a solution, then there exist infinitely many solutions. One last comment should be made. The spike function we generated is not unique. Simply shifting the nonzero point would generate and infinite number of spike functions which would all be in the null space of K. That is, K has an incredibly rich null space.

2.3.2 Eigenvalues

Difficulty from the Matrix



K: functions to functions, \underline{K}_n :: vectors to vectors!

As shown above, discretizing the integral equation by combining a piecewise constant charge density representation with collocation at subinterval centers results in a system of equations which relates the subinterval σ_i 's to the collocation point potentials. From this perspective, the matrix above can be thought of as a discrete representation of the operator K. We denote the matrix with \underline{K}_n with an underline to indicate that it is a matrix, and was generated using a discretization with n basis functions. We also denoted the vectors $\vec{\sigma}_n$ and $\vec{\Psi}_n$ with arrows to avoid confusing vectors with functions. Later, we will need an operator version of the discretized representation of the operator K, but for the moment, the matrix is sufficient.

Numerical Results with Increasing n
If the operator K is singular, one might expect to see that reflected in the eigenvalues of a matrix generated by discretizing K. In particular, one would expect the matrix to have eigenvalues that are near zero.



In the figure above, the eigenvalues of matrices generated by discretizing K for the 1-D problem are plotted. Discretizing using 10 subintervals generates a matrix with 10 eigenvalues plotted in blue. The blue eigenvalue closest to zero is ≈ 0.01 . As the discretization is refined to 20 subintervals, the minimum eigenvalue (plotted in red) drops to ≈ 0.003 , and with 40 subintervals the minimum eigenvalue (plotted in green) drops to ≈ 0.0009 . Examining this data suggests that as the discretization is refined, the generated matrix more accurately reflects the operator K, and therefore the matrix is becoming closer to being singular.

As the discretization is refined, the matrix is larger and has more eigenvalues. Notice that as the discretization is refined from n = 10 to n = 20 to n = 40, all the additional eigenvalues are closer to zero.

Intuition About Eigenvalues

As the discretization is refined, $\sigma_0(x)$ becomes better approximated



As the discretization is refined, K's null space can be more accurately represented.

As an alternative view of why refining the discretization for the first kind equation produces a matrix with more and more smaller eigenvalues, consider the figures above. In the top plot, one of the basis functions is plotted for a coarse discretization. In the bottom plot, one of the basis functions is plotted for a finer discretization. Notice that as the discretization is refined, these basis functions look progressively more like the spike function mentioned previously. And since the spike function is in the null space of K, one would expect that finer discretizations would generate "spikier" basis functions whose associated eigenvalues would be near zero.

2.4 Second Kind Equation Has Fewer Problems

Second Kind equation

$$((I+K)\sigma)(x) \equiv \sigma(x) + \int_{-1}^{1} |x - x'|\sigma(x')dS' = \Psi(x)$$

$$\Rightarrow (I+K)\sigma = \Psi$$
(14)

$$(I+K)\left(\begin{array}{c} 1\\ -1\\ -1\\ -1\end{array}\right) + \begin{array}{c} \sigma_{0}(x)=0, x\neq 0, \sigma_{0}(0)=1\\ -1\\ -1\\ -1\\ 1\end{array}\right)$$
$$= \begin{array}{c} 1\\ -1\\ -1\\ 1\end{array} \neq \Psi$$

 $(I+K)(\sigma_0+\sigma) \neq (I+K)\sigma$

As shown in equation (14), the abstract operator for the second-kind equation is denoted by I + K, where I here is just the identity operator and K is the integral operator.

To see why the spike function, σ_0 , is not in the null space of the operator I + K, or equivalently that

$$(I+K)(\sigma_0+\sigma) \neq (I+K)(\sigma)$$

consider the figures above. If a spike is added to a smooth σ , the (I + K) operator will preserve the spike. Another way to see this is to consider that since σ_0 is in the null space of K,

$$(I+K)\sigma_0 = (I)\sigma_0 + K\sigma_0 = \sigma_0 \neq 0.$$

2.4.1 Eigenvalues

Numerical Results with Increasing n

$$n = 10$$

 $n = 20$ $n = 40$

Eigenvalues do not get closer to zero.

As we noted before, the matrix associated with discretizing the operator I + K is identical to the sum of the identity matrix and the matrix associated with discretizing K alone. In the plot above, we once again present the eigenvalues generated by discretizing the 1-D example problem. Discretizing using 10 subintervals generates a matrix with 10 eigenvalues plotted in blue. The blue eigenvalue closest to zero is ≈ 0.2 . As the discretization is refined to 20 subintervals, the minimum eigenvalue (plotted in red) is still ≈ 0.2 , and with 40 subintervals the minimum eigenvalue (plotted in green) is still ≈ 0.2 . Examining this data suggests that as the discretization is refined, and the generated matrix more accurately reflects the operator I + K, the matrix is not becoming more singular. In fact, the eigenvalues are accumulating near one, an unsurprising result given that the eigenvalues of the discretized K operator were accumulating at zero.

▷ **Exercise 2** Estimate how many iterations will be needed for a Krylovsubspace based algorithm to converge for the 1-D discretized second-kind example. Will the number of iterations increase as the discretization is refined?

▷ **Exercise 3** Suppose the integral equation were changed to

$$\Psi(x) = \sigma(x) + \frac{1}{\lambda} \int_{-1}^{1} |x - x'| \sigma(x') dS'$$

For what value of λ would the solution no longer be unique. (you can answer this just by looking at the eigenplot above).

As the above exercise makes clear, a second-kind integral equation does not always have a unique solution. However, a first-kind equation almost never has a unique solution, the exception being when the Green's function is singular, as we will investigate in a subsequent lecture.

3 Theory of 2nd Kind Equations

The convergence theory for discretization methods applied to second-kind integral equations has an elegant simplicity, but only when examined using a carefully chosen abstraction. The theory is also surprisingly practical; the insights gained can be used to construct very high-order discretization schemes.

3.1 Comparison problem

General Second kind integral equation

$$\Psi(x) = \sigma(x) + \int G(x, x')\sigma(x')dx' \Rightarrow \Psi = (I + K)\sigma$$
(15)

Discrete matrix equivalent

$$\vec{\Psi}_n = (I + \underline{K}_n) \,\vec{\sigma}_n \tag{16}$$

How to compare function σ to vector $\vec{\sigma}_n$? How to compare operator K to matrix \underline{K}_n ?

One approach to overcoming the comparison problems is to construct representations of the discretization that are functions and operators on functions. The most obvious approach to generating the functions associated with a discretization is by interpolation, but generating the operators associated with the discretization is a little more subtle.

3.1.1 Operator-Function Notation

General Second kind integral equation

$$\Psi(x) = \sigma(x) + \int G(x, x')\sigma(x')dx' \Rightarrow \Psi = (I + K)\sigma$$
(17)

Discretized operator-function equivalent

$$\Psi_n = (I + K_n) \,\sigma_n \tag{18}$$

 σ_n , Ψ_n are functions of x (e.g. by interpolation) K_n maps functions to functions like K (How constructed)?

3.1.2 Orthogonal Galerkin

Representation $\sigma_n(x) = \sum_{i=1}^n \sigma_{ni} \varphi_i(x) \quad \int \varphi_i(x) \varphi_j(x) = \delta(i-j)$ **Projection** $\sigma_n = (P\sigma)(x)$

$$(P\sigma)(x) \equiv \sum_{i=1}^{n} \overbrace{\left(\int \sigma(x)\varphi_i(x)dx\right)}^{\sigma_{ni}} \varphi_i(x)$$



If the density sigma is to be approximated by a weighted combination of n orthogonal basis functions, then the functional representation σ_n associated with the vector of weights $\vec{\sigma}_n$ is given. If the basis functions are orthonormal, then the σ_n associated with an arbitrary σ can be constructed by simple projection. It is worth noting that the projection operator, denoted P, has no effect on σ_n . That is $P\sigma_n = \sigma_n$.

 \triangleright **Exercise 4** Why does the formula for the projection operator above require orthogonality of the basis functions?

If a Galerkin method is used to discretized the integral equation, then the associated operator is easy to construct, as shown below.

3.1.3 Ortho Galerkin Operator

$$(K\sigma_n)(x) = (KP\sigma)(x) = \sum_{i=1}^n \sigma_{ni} \int G(x, x')\varphi_i(x')dx'$$
$$(PKP\sigma)(x) = \sum_{j=1}^n \left(\int \varphi_j(x)KP\sigma(x)dx\right)\varphi_j(x)$$
$$= \sum_{j=1}^n \left(\sum_{i=1}^n \sigma_{ni} \int \int \varphi_j(x)G(x, x')\varphi_i(x')dxdx'\right)\varphi_j(x)$$
$$P(I + KP)\sigma_n = (I + PKP)\sigma_n = P\Psi$$
$$(I + K_n)\sigma_n = \Psi_n$$

The last equation on the above slide contains a subtle point. $P(I+KP)\sigma_n$ really equals $(P + PKP)\sigma_n$. However, P is equivalent to the identity operator when applied to σ_n , as projecting σ_n reproduces σ_n . So, we are free to conveniently chose to interpret (P + PKP) as (I + PKP) as a difference appears only when applying the operator to functions that are not weighted combinations of the nGalerkin basis functions.

For second-kind integral equations, one can prove a convergence theory for almost any reasonable discretization scheme, assuming that the equation has a unique solution. As noted above, second-kind integral equations do not necessarily have unique solutions, but we will restrict ourselves to the unique solution case in analyzing convergence. In particular, we will assume that the secondkind integral equation operator has a bounded inverse.

Before beginning the derivation, let's readdress the notation definitions. Let K denote the integral operator, and therefore the general form is

$$K\sigma = \int G(x,x')\sigma(x')dx'$$

Let σ_n denote a numerical approximation to σ on x based on using n basis functions. Note here that σ_n is a function of x and would typically be given by

$$\sigma_n(x) = \sum_{i=1}^n \sigma_{ni} \varphi_i(x).$$

Let K_n be the discrete representation of the integral operator. Note that K_n is *not* the matrix \underline{K}_n , but an operator that maps a function of x into another function of x. For example, if the discretization scheme uses a basis to approximate σ , and the basis weights were determined by a collocation scheme, a not necessarily unique associated K_n could be given by

$$(K_n\sigma)(x) = V\left(\int G(x,x')P\sigma(x')dx'\right)$$

where in the orthonormal basis set case

$$(P\sigma)(x) = \sum_{i=1}^{n} \left(\int \sigma(x')\varphi_i(x')dx' \right) \varphi_i(x), \tag{19}$$

and

$$(Vu)(x) = \sum_{i=1}^{n} u(x_{c_i})\varphi_i(x).$$
 (20)

where u is a arbitrary function used to define the action of operator V. Equations (19) and (20) deserve some explanation. The piecewise constant basis is orthonormal, so the formula in equation (19) is a simple projection of σ onto the basis. If centroid collocation is used, then the discrete potentials computed by evaluating the integral operator at the collocation points must be converted to a function of x by interpolation. In equation (20), the $\varphi_i(x)$'s act as interpolation functions.

With the examples of how Galerkin and centroid-collocation discretization schemes lead to function and operator representations, the second-kind integral equation convergence theory can be presented in a very transparent fashion, as will be show below. What the theorem demonstrates is that if a discretization scheme generates progressively more accurate representations of the integral operator as n increases, then the discretization method converges. That is,

$$\lim_{n \to \infty} \|\sigma - \sigma_n\| \to 0,$$

1

where the comparison between σ and σ_n is unambigious as both are functions of x.

 \triangleright **Exercise 5** Suppose a nonorthogonal basis is used to represent σ . How would the projection operator in equation (19) change? \blacksquare

Main Theorem 3.2

Given
$$(I + K)\sigma = \Psi$$
 & $||(I + K)^{-1}|| < C$
Means Equation uniquely solvable
 $(I + K_n)\sigma_n = \Psi_n$
Reminder of Discrete Equivalent

Consistency:

 $\frac{\lim_{n \to \infty} \max_{\|\sigma\|=1} \|(K - K_n)\sigma\| \to 0 \quad \text{and} \quad \lim_{n \to \infty} \|\Psi - \Psi_n\| \to 0$ $\mathbf{n} \quad \lim_{n \to \infty} \|\sigma - \sigma_n\| \to 0$ If \mathbf{Then}

Rough Proof 3.3

To derive a relationship between the errors in the computed solution and the errors in the operator representation, we write the exact equation alongside the discrete equation.

Discretized Equation **Exact Equation** $\Psi_n = (I + K_n)\sigma_n$ $\Psi = (I + K)\sigma$ $\Rightarrow \Psi - \Psi_n = (I + K)\sigma - (I + K_n)\sigma_n$ $\Rightarrow (\Psi - \Psi_n) = (\sigma - \sigma_n) + K\sigma - K_n \sigma_n$ $\Rightarrow (\Psi - \Psi_n) = (\sigma - \sigma_n) + K\sigma - K_n\sigma + K_n\sigma - K_n\sigma_n$ $\Rightarrow (\Psi - \Psi_n) = (I + K_n)(\sigma - \sigma_n) + (K - K_n)\sigma$ $\Rightarrow (\Psi - \Psi_n) - (K - K_n)\sigma = (I + K_n)(\sigma - \sigma_n)$ $\Rightarrow (I+K_n)^{-1} \left[(\Psi - \Psi_n) - (K - K_n) \sigma \right] = (\sigma - \sigma_n).$

The results on the slide below skip many of these intermediate steps but present the essential results.



The equation for the solution error (previous slide)

$$\underbrace{(\sigma_n - \sigma)}_{\text{solution error}} = (I + K_n)^{-1} (K - K_n) \sigma$$

Taking norms	$ \sigma_n - \sigma $	\leq	$\underbrace{ (I+K_n)^{-1} }_{}$	$\underbrace{ (K-K_n)\sigma }$
	Error which		Needs a	Goes to
	should go to		bound, that is	zero
	zero as n		stability	by <i>consistency</i>
	increases			

We complete deriving a relationship between the errors in the computed solution and the errors in the operator representation. In order to establish that consistency implies convergence, the inverse of the discretized operator must be bounded.

3.3.1 Stability Bound

Norm of solution error $||(\sigma_n - \sigma)|| \le ||(I + K_n)^{-1}|| ||(K - K_n)\sigma||$

Deriving the stability bound $(I + K_n)^{-1} = [I + K - (K - K_n)]^{-1} = [(I - (I + K)^{-1}(K - K_n)]^{-1}(I + K)^{-1}(K - K)^{-1}(K -$

Taking norms

$$||(I+K_n)^{-1}|| \leq \underbrace{||(I+K)^{-1}||}_{\text{Bounded by C}} || (I-(I+K)^{-1}(K-K_n))^{-1}|$$

by Assumption

Repeating from last slide

$$||(I+K_n)^{-1}|| \leq \underbrace{||(I+K)^{-1}||}_{\text{Bounded by C}} || \left(I - (I+K)^{-1}(K-K_n)\right)^{-1} ||$$

by Assumption

Bounding terms
$$||(I + K_n)^{-1}|| \le \frac{C}{1 - ||(I + K)^{-1}(K - K_n)||} < 2C$$
 for

Will be less than 0.5 for n larger than some n_0 by consistency and solvability

 $n \ge n_0$

Final result:

$$\lim_{n \to \infty} ||(K - K_n)\sigma|| = 0$$

IMPLIES

$$\lim_{n \to \infty} ||(\sigma_n - \sigma)|| = 0$$

What does this mean?

The discretization convergence of a second kind integral equation solver only depends on how well the integral is approximated.

The final result, noted on the above slide, is that the solution error is bounded by a constant multiplying the error in the integral operator representation. This suggests that any method which can accurately represent the integral operator can be used to discretize a second-kind integral equation.

4 Nystrom Method

4.1 1-D Second Kind Example

4.1.1 Collocation Discretization

Integral Equation

$$\Psi(x) = \sigma(x) + \int_{-1}^{1} G(x, x') \sigma(x') dS' \qquad x \in [-1, 1]$$

Apply quadrature to Collocation equation

$$\Psi(x_i) = \sigma(x_i) + \int_{-1}^{1} G(x_i, x') \sigma(x') dS'$$
$$\Rightarrow \Psi(x_i) = \sigma(x_i) + \sum_{j=1}^{n} w_j G(x_i, x_j) \sigma(x_j)$$

After applying quadrature to Collocation:

$$\Psi(x_i) = \sigma(x_i) + \sum_{j=1}^n w_j G(x_i, x_j) \sigma(x_j)$$

 x_i is a collocation point

 x_i 's are quadrature points

Now set quadrature points = collocation points

In Gaussian quadrature, described in the previous lecture, an integral is approximated using weighted combinations of the integrand. As a reminder, the Gaussian quadrature formula for integrating a function on the unit interval is

$$\int_0^1 f(x)dx \simeq \sum_{i=1}^n w_i f(x_i)$$

where the x_i 's are the evaluation points given by the zeros of an n^{th} -order orthogonal polynomial on the unit interval, and the w_i 's are the weights determined by solving exactness equations.

The key idea behind a Nystrom method for discretizing an integral equation is to use the Gaussian quadrature evaluation points as the test points in a collocation method for solving an integral equation. Then, the collocation method integrals can be approximated using the Gaussian quadrature scheme, resulting in a system of equations which only require evaluations of the integrand at the test=quadrature points. The second kind theory predicts that the error in such a scheme is proportional to the error in the quadrature scheme for computing the collocation integrals.

Set quadrature points = collocation points \mathbf{r}

$$\Psi(x_1) = \sigma_{n1} + \sum_{j=1}^n w_j G(x_1, x_j) \sigma_{nj}$$

$$\vdots$$

$$\Psi(x_n) = \sigma_{nn} + \sum_{j=1}^n w_j G(x_n, x_j) \sigma_{nj}$$

System of n equations in n unknowns

Collocation equation per quad/colloc point Unknown density per quad/colloc point

4.1.2 Discretization-Matrix Comparison

$\begin{array}{c} \textbf{Nystrom Matrix} \\ \begin{bmatrix} 1 + w_{l}G(x_{1}, x_{1}) & \cdots & w_{n}G(x_{1}, x_{n}) \\ \vdots & \ddots & \vdots \\ w_{l}G(x_{n}, x_{1}) & \cdots & 1 + w_{n}G(x_{n}, x_{n}) \end{bmatrix} \begin{bmatrix} \sigma_{n1} \\ \vdots \\ \sigma_{nn} \end{bmatrix} = \begin{bmatrix} \Psi(x_{1}) \\ \vdots \\ \Psi(x_{n}) \end{bmatrix}$



Nystrom Matrix

Just Green's function evaluations – No integrals Entries each have a quadrature weight Collocation points are quadrature points High order quadrature=faster convergence?

Piecewise Constant Collocation Matrix

Integrals of Green's function over line sections Distant terms equal Green's function Collocation points are basis function centroids Low order method always

4.2 K_n and Ψ_n for Nystrom Method

$$K_n \sigma = \sum_{i=1}^n \left(\sum_{j=1}^n w_j G(x_i, x_j) \sigma(x_j) \right) \varphi_i(x)$$

$$\Psi_n = \sum_{i=1}^n \Psi(x_i)\varphi_i(x)$$

4.3 Convergence

4.3.1 Theorem

In the limit as $n \to \infty$ (number of quad points $\to \infty$)

The discretization error $= \max_{\substack{||\sigma||=1}} ||(K - K_n)\sigma|| \to 0$ AT THE SAME RATE as the underlying quadrature!!

Gauss Quadrature \Rightarrow Exponential Convergence!

4.3.2 Comparison

$$\cos 2\pi x = \sigma(x) + \int_{-1}^{1} (x - x')^2 \sigma(x') dS'$$



4.3.3 Caveat

If Nystrom method can have exponential convergence, why use anything else? Gaussian quadrature has exponential convergence for **nonsingular** kernels Most physical problems of interest have **singular kernels** $(\frac{1}{r}, \frac{\exp ikr}{r}, \text{etc})$

5 Summary

Convergence Issues in 1D

1st and 2nd kind integral equations, null spaces

Convergence for second kind equations

Show consistency and stability issues

Nystrom methods

High order convergence

Did not address singular integrands

1.4 Radiation Conditions and Formulations

Numerical Methods for PDEs

Integral Equation Methods, Lecture 4 Radiation Conditions and Formulations

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1 Outline

Laplace Problems Exterior Radiation Condition Potential Representations Monopole and Dipole Densities Principle Value Ansantz and Green's Theorem Dirichlet and Neumann problems

2 3D Problems

2.1 3D Laplace Equation

Laplace's equation in 3-D

$$\nabla^2 u(\vec{x}) = \frac{\partial^2 u(\vec{x})}{\partial x^2} + \frac{\partial^2 u(\vec{x})}{\partial y^2} + \frac{\partial^2 u(\vec{x})}{\partial z^2} = 0$$

where

$$\vec{x} = (x, y, z) \in \Omega$$

and Ω is bounded by Γ .



The exterior problem is simply the region $\Re - \overline{\Omega}$ of the interior problem.



One feature of using integral equation methods is that exterior problems can be solved using the same surface discretization required to solve an interior problem (assuming a linear space-invariant problem like Laplace's equation). This is true even though the exterior domain is infinite and the interior domain is finite. Exterior problems do introduce an additional complication, one must consider the boundary condition "at infinity" (later).

2.2 Boundary Conditions

2.2.1 Dirichlet

Dirichlet Condition

 $u(\vec{x}) = u_{\Gamma}(\vec{x}) \qquad \vec{x} \in \Gamma$

Interior Problem



Can you determine the solution to Laplace's equation inside the sphere? The solution of the interior Dirichlet problem is unique.

2.2.2 Neumann



The solution of the <u>interior</u> Neumann problem is not unique. For the solution of the <u>exterior</u> Neumann problem to be unique, it is sufficient to impose a *radiation condition*. In this case, a radiation condition would be a specification of how $u(\vec{x})$ approaches zero as $\vec{x} \to \infty$.

2.2.3 Exterior

Dirichlet Boundary Condition

$$u(\vec{x}) = u_{\Gamma}(\vec{x}) \qquad \vec{x} \in \Gamma$$

Neumann Boundary Condition

$$\frac{\partial u(\vec{x})}{\partial n_{\vec{x}}} = \frac{\partial u_{\Gamma}(\vec{x})}{\partial n_{\vec{x}}} \qquad \vec{x} \in \Gamma$$

PLUS

A Radiation Condition

3 Radiation Condition

3.1 Condition at "Infinity"



3.2 Types of Conditions

A radiation condition of the form

$$\lim_{\|\vec{x}\| \to \infty} u(\vec{x}) \to 0$$

is not specific enough! Need

$$\lim_{\|\vec{x}\| \to \infty} u(\vec{x}) \to O(\|\vec{x}\|^{-1})$$
OR
$$\lim_{\|\vec{x}\| \to \infty} u(\vec{x}) \to O(\|\vec{x}\|^{-2})$$

3.3 Examples

The criteria for chosing a radiation condition are best understood by considering several physical examples.

3.3.1 Radiated Heat Problem Set-Up

$$\nabla^2 u(\vec{x}) = 0 \qquad \vec{x} \in \Omega$$
$$u(\vec{x}) = u_0(\vec{x}) \qquad \vec{x} \in \Gamma$$
Radiated Heat =
$$\int_{\Gamma} \frac{\partial u}{\partial n} dS$$



Limit of Expanding Domain

- Infinite Problem Limit as $R \to \infty$
- Heat Leaving Sphere $\int_{Sphere_R} \frac{\partial u}{\partial n} dS \text{ is Constant!}$

3.3.2 Rad Heat Case 1

$$\lim_{\|\vec{x}\|\to\infty} u(\vec{x}) \to O(\|\vec{x}\|^{-1}) \to \lim_{\|\vec{x}\|\to\infty} \frac{\partial u(\vec{x})}{\partial n} \to O(\|\vec{x}\|^{-2})$$

Since the surface of a sphere increases as R^2 :

$$\lim_{R \to \infty} \int_{Sphere_R} \frac{\partial u}{\partial n} dS \to \text{Constant}$$

Radiation condition models net heat loss.

3.3.3 Rad Heat Case 2

$$\lim_{\|\vec{x}\| \to \infty} u(\vec{x}) \to O(\|\vec{x}\|^{-2}) \to \lim_{\|\vec{x}\| \to \infty} \frac{\partial u(\vec{x})}{\partial n} \to O(\|\vec{x}\|^{-3})$$
AND

$$\int_{Sphere_R} \frac{\partial u}{\partial n} dS \to 0$$

Can NOT model heat loss!

3.3.4 Heat Transfer



Heat flows from higher temperature object to lower temperature object, but no heat radiates out.

If

$$\lim_{\|\vec{x}\| \to \infty} u(\vec{x}) \to O(\|\vec{x}\|^{-2})$$

Then

$$\lim_{\|\vec{x}\| \to \infty} \frac{\partial u(\vec{x})}{\partial n} \to O(\|\vec{x}\|^{-3})$$

And

$$\int_{Sphere_R} \frac{\partial u(\vec{x})}{\partial n} dS \to 0$$

This condition ensures all heat transferred.



The above image is supposed to represent two scenarios, each senario has two conducting bodies.

In the first senario, on the left, there are two conductors that are treated as if there is a voltage source across them. There is a positive charge on one conductor, a negative charge on the other conductor, and the net charge should be zero. If the net charge on the two conductors is zero, then the integral of the normal electric field, $\frac{\partial u(\vec{x})}{\partial n}$) over a bounding sphere, one that contains both spherical conductors, should be zero. This is just a statement of the well-known Gauss's theorem in electrostatics. In particular

$$\int_{Sphere_R} \frac{\partial u(\vec{x})}{\partial n} dS = 0$$

for any $Sphere_R$ containing the two conductors. In order for this integral to be zero regardless as the radius of $Sphere_R$ approaches ∞ ,

$$\lim_{\|\vec{x}\| \to \infty} \frac{\partial u(\vec{x})}{\partial n}$$

must decay at least as fast as $O(\|\vec{x}\|^{-3})$. Therefore, $u(\vec{x})$ must decay like $O(\|\vec{x}\|^{-2})$.

For the senario on the right, the two spherical conductors are set to two different potentials with respect to the point at infinity. In such a case, it is unlikely that the sum of the charge on the two spheres will be zero,

$$\int_{Sphere_R} \frac{\partial u(\vec{x})}{\partial n} dS \neq 0$$

and therefore $u(\vec{x})$ should not be forced to decay any faster than $O(\|\vec{x}\|^{-2})$.

3.3.6 Potential Flow



Ω

Assumptions Irrotational flow (velocity = potential gradient):

 $\mathbf{v}(\vec{x}) = \nabla u(\vec{x})$

Air is incompressible (velocity divergence free):

 $\nabla \cdot \mathbf{v}(\vec{x}) = \nabla^2 u(\vec{x}) = 0.$

Nonpenetrating wing boundary condition:

 $\nabla u(\vec{x}) \cdot \mathbf{n}(\vec{x}) = \mathbf{v}_{wing}(\vec{x}) \cdot \mathbf{n}(\vec{x}).$

What is the right radiation condition?

4 Formulations – Problem Types

4.1 Single Domain



4.2 Coupled Domain

Example: Bimetallic Electrical Conductivity



Potential and Electric Current Continuity:

$$u(\vec{x}^{+}) = u(\vec{x}^{-}) \qquad \vec{x} \in \Gamma$$
$$\alpha^{+} \frac{\partial u(\vec{x}^{+})}{\partial n_{\vec{x}}} = \alpha^{-} \frac{\partial u(\vec{x}^{-})}{\partial n_{\vec{x}}} \qquad \vec{x} \in \Gamma$$

An example of a coupled domain problem would be a conductivity problem involving multiple materials. To determine the electrical conductivity between two terminals of an object made of multiple materials, one would determine the ratio of the voltage across the object's terminals and the current flowing through the object. Electrical current density in an ideal linear conductor is a vector quantity given by the gradient of the potential, ∇u , scaled by a factor known as the conductivity of the material. In an ideal linear conductor there is no accumulation of charge at any interior point, implying that the current density has zero divergence. Therefore, the potential in an ideal linear conductor satifies Laplace's equation, $\nabla^2 u = 0$. If an object is made of multiple materials with different electrical conductivities, then the boundary between materials satisfies interface conditions. At the boundary between materials, both the potential and the current density in the surface-normal direction are continuous. Since the conductivities of the two materials are different, continuity of the current density implies a jump in the gradient of the potential across the material boundary.

4.3 Normals



Typically, the surface normal is assumed to point in the direction from the interior domain to the exterior domain. There are many situations where this typical practice is confusing or ambigious, so it is often necessary to be explicit about the direction of the normal.

5 Surface Density Integrals

5.1 Monopole & Dipole

Potential due to a monopole density (σ) :

$$u(\vec{x}) = \int_{\Gamma} \frac{1}{\|\vec{x} - \vec{x}'\|} \sigma(\vec{x}') d\Gamma'$$

Potential due to a dipole density (μ) :

$$u(\vec{x}) = \int_{\Gamma} \frac{\partial}{\partial n_{\vec{x}'}} \frac{1}{\|\vec{x} - \vec{x}'\|} \mu(\vec{x}') d\Gamma'$$

where the normal points out of the domain Ω bounded by Γ .

MONOPOLE:	I	Dipole:
++++++++++++++++++++++++++++++++++++	$ abla^2 u = 0$ satisfied $\forall \sigma, \mu$	$\begin{array}{c} + + + + + + + + + + + + + + + + + + +$

Monopole or dipole densities can be used to generate potentials that satisfy $\nabla^2 u(\vec{x}) = 0$ for all $\vec{x} \in \Omega$. The monopole and dipole potentials differ in the radiation condition they satisfy. If the surface, γ , is finite in extent, then in the limit as $\|\vec{x}\| \to \infty$, the monopole potential decays like $\|\vec{x}\|^{-1}$, and the dipole potential decays like $\|\vec{x}\|^{-1}$.

Either representation can be used to derive surface integral equations, but care must be used when evaluating the associated potentials when $\vec{x} \in \Gamma$.

5.2 Surface Potentials

The monopole potential is continuous as x passes through Γ , so

$$u_{\Gamma}(\vec{x}) = \int_{\Gamma} \frac{1}{\|\vec{x} - \vec{x}'\|} \sigma(\vec{x}') d\Gamma' \quad \vec{x} \in \Gamma$$

The dipole potential "jumps" as x passes through Γ , so the limit as $\vec{x} \to \Gamma$ of

$$u(\vec{x}) = \int_{\Gamma} \frac{\partial}{\partial n_{\vec{x}'}} \frac{1}{\|\vec{x} - \vec{x}'\|} \mu(\vec{x}') d\Gamma'$$

depends on how Γ is approached.



Don't be put-off by the graph above. The monopole potential is continuous (it does not go off to infinity, as it may seem to in the above figure), but it is not continuously differentiable, there will be a discontinuity in the derivative at x_0 .



5.2.1 Principle Value Integral

If f(y) is singular for some $y = x_0$, where $x_0 \in \Gamma$, then the principle value integral is

$$\int_{\Gamma}^{PV} f(\vec{y}) d\Gamma \equiv \lim_{\epsilon \to 0} \int_{\Gamma - B(x_0, \epsilon) \cap \Gamma} f(\vec{y}) d\Gamma$$

when $B(x_0, \epsilon)$ is the ϵ radius ball about x_0 .

The P.V. is a special kind of limit

Limit of deleting and ever shrinking portion of the integration domain. NOT EQUIVALENT TO limiting processes on f!



5.2.2 Monopole Derivative (MD)

Consider a cube geometry:



$$\frac{\partial u_{\Gamma}(\vec{x})}{\partial n_{\vec{x}}} = \lim_{\vec{x} \to \Gamma^+} \frac{\partial}{\partial n_{\vec{x}}} \int_{\Gamma} \frac{1}{\|\vec{x} - \vec{x}'\|} \sigma(\vec{x}') d\Gamma' \qquad \vec{x} \in \Gamma$$

The plus (+) in Γ^+ indicates exterior approach.

In the above slide, we consider computing the normal derivative of the monopole potential just outside the boundary γ . As will be shown in the next few slides, the derivative can be represented as the sum of a principle value integral and an extra term.

5.2.3 MD Disk Removal



Consider the entire side panel on the right of the cube in the above slide, and consider evaluating the normal derivative of the potential generated by a monopole distribution on the suface of the cube's right side. Specifically, assume that we wish to evaluate the derivative at a point \vec{x} in the center of the green disk on the cube's right side. The matter is complicated by the fact that the integrand goes to infinity when $\vec{x} = \vec{x}'$. Thus, we need to break up the integral into two pieces. One piece is the entire panel minus the green disk, and the other piece is just the green disk.

5.2.4 MD Disk Picture

$$\frac{\partial u_{\Gamma}(\vec{x})}{\partial n_{\vec{x}}} = \lim_{\vec{x} \to \Gamma^{+}, \epsilon \to 0} \left[\frac{\partial}{\partial n_{\vec{x}}} \int_{\Gamma - B(x,\epsilon)} \frac{\sigma(\vec{x}')}{\|\vec{x} - \vec{x}'\|} d\Gamma' + \frac{\partial}{\partial n_{\vec{x}}} \int_{B(x,\epsilon)} \frac{\sigma(\vec{x}')}{\|\vec{x} - \vec{x}'\|} d\Gamma' \right]$$



The first integral is the Principle Value Integral and the second integral is the integral of just the disk.

Given that the disk was extracted from right the surface of a cube, the disk is flat, and the normal is in the z-axis direction.

5.2.5 MD Disk Eval

$$\lim_{\vec{x}\to\Gamma^{+}} \frac{\partial}{\partial n_{\vec{x}}} \int_{B(x,\epsilon)} \frac{1}{\|\vec{x}-\vec{x}'\|} \sigma(\vec{x}') d\Gamma'$$

$$\approx \lim_{z\to0^{+}} \frac{\partial}{\partial z} \int_{0}^{2\pi} \int_{0}^{\epsilon} \frac{\sigma(\vec{x})}{\sqrt{r^{2}+z^{2}}} r dr d\theta$$

$$= \lim_{z\to0^{+}} \frac{\partial}{\partial z} 2\pi \sigma(\vec{x}) \left[\sqrt{\epsilon^{2}+z^{2}} - |z| \right]$$

$$= -2\pi \sigma(\vec{x}).$$

Note 1

Disk Evaluation Math

For this problem, it is quite straightforward to see how one changes from cartesian to cylindrical coordinates. But, the algebra and calculus involved in solving this integral may not be as straightforward, herein is presented one method, broken-down into bite-size pieces:

$$\lim_{z \to 0^+} \frac{\partial}{\partial z} \int_0^{2\pi} \int_0^{\epsilon} \frac{\sigma(\vec{x})}{\sqrt{r^2 + z^2}} r dr d\theta.$$

Use trigonometric substitution to solve the integral with respect to r. Substitute $r = z \tan \alpha$ and $dr = z \sec^2 \alpha d\alpha$ and simplify to get the following integral:

$$= \lim_{z \to 0^+} \frac{\partial}{\partial z} \int_0^{2\pi} \int_0^{\alpha(\epsilon)} \frac{\sigma(\vec{x}) z \sin \alpha}{\cos^2 \alpha} d\alpha d\theta.$$

This integral is easily solved using direct substitution of $u = \cos \alpha$:

$$= \lim_{z \to 0^+} \frac{\partial}{\partial z} \int_0^{2\pi} \int_0^{u[\alpha(\epsilon)]} -z \frac{\sigma(\vec{x})}{u^2} du d\theta = \lim_{z \to 0^+} \frac{\partial}{\partial z} \int_0^{2\pi} z \frac{\sigma(\vec{x})}{u} \bigg|_0^{u[\alpha(\epsilon)]} d\theta.$$

Plug back in for $u = \cos \alpha = \frac{z}{\sqrt{r^2 + z^2}}$

$$= \lim_{z \to 0^+} \frac{\partial}{\partial z} \int_0^{2\pi} \sigma(\vec{x}) \sqrt{r^2 + z^2} \bigg|_0^{r-\epsilon} d\theta$$
$$= \lim_{z \to 0^+} \frac{\partial}{\partial z} \int_0^{2\pi} \sigma(\vec{x}) \left[\sqrt{\epsilon^2 + z^2} - \sqrt{z^2} \right] d\theta$$

Integrating the last part is quite simple,

$$= \lim_{z \to 0^+} \frac{\partial}{\partial z} \sigma(\vec{x}) \left(\sqrt{\epsilon^2 + z^2} - \sqrt{z^2} \right) \theta \Big|_0^{2\pi} = 2\pi \sigma(\vec{x}) \lim_{z \to 0^+} \frac{\partial}{\partial z} \left(\sqrt{\epsilon^2 + z^2} - \sqrt{z^2} \right)$$

Finally, take the derivatives with respect to z, the normal,

$$= 2\pi\sigma(\vec{x}) \lim_{z \to 0^+} \left(\frac{z}{\sqrt{\epsilon^2 + z^2}} - \frac{z}{\sqrt{z^2}} \right).$$

It can now be seen that, since $\lim_{z \to 0^+} \frac{z}{\sqrt{\epsilon^2 + z^2}} = 0$ and $\lim_{z \to 0^+} \frac{z}{\sqrt{z^2}} = \text{sign } z$ that
 $\lim_{z \to 0^+} \frac{\partial}{\partial z} \int_0^{2\pi} \int_0^{\epsilon} \frac{\sigma(\vec{x})}{\sqrt{r^2 + z^2}} r dr d\theta = -2\pi\sigma(\vec{x}).$

5.2.6 MD Final

$$\begin{aligned} \frac{\partial u_{\Gamma}(\vec{x})}{\partial n_{\vec{x}}} &= \lim_{\vec{x} \to \Gamma^{+}} \frac{\partial}{\partial n_{\vec{x}}} \int_{\Gamma} \frac{1}{\|\vec{x} - \vec{x}'\|} \sigma(\vec{x}') d\Gamma' \\ &= \lim_{\vec{x} \to \Gamma^{+}} \lim_{\epsilon \to 0} \left[\frac{\partial}{\partial n_{\vec{x}}} \int_{\Gamma - B(x,\epsilon)} \frac{\sigma(\vec{x}')}{\|\vec{x} - \vec{x}'\|} d\Gamma' + \frac{\partial}{\partial n_{\vec{x}}} \int_{B(x,\epsilon)} \frac{\sigma(\vec{x}')}{\|\vec{x} - \vec{x}'\|} d\Gamma' \right] \\ &= \int_{\Gamma}^{PV} \frac{\partial}{\partial n_{\vec{x}}} \frac{1}{\|\vec{x} - \vec{x}'\|} \sigma(\vec{x}') d\Gamma' - 2\pi\sigma(\vec{x}') \end{aligned}$$

5.2.7 Dipole Potentials (DP)

If Γ is a flat surface

$$\int_{\Gamma}^{PV} \frac{\partial}{\partial n_{\vec{x}'}} \frac{1}{\|\vec{x} - \vec{x}'\|} d\Gamma' = 0 \qquad \vec{x} \in \Gamma.$$

Why? Rewrite using explicit form of integrand

$$\int_{\Gamma}^{PV} \frac{(\vec{x} - \vec{x}')^T n_{\vec{x}'}}{(\|\vec{x} - \vec{x}'\|)^3} d\Gamma'$$

Integrand is zero when $\vec{x} - \vec{x}'$ orthogonal to surface normal

5.2.8 DP Flat Surface

Flat result applies locally on smooth surfaces.



5.2.9 DP General Surface

If Γ is a general surface

$$u_{\Gamma}(\vec{x}) = 2\pi\mu(\vec{x}) + \int_{\Gamma}^{PV} \frac{(\vec{x} - \vec{x}')^T n_{\vec{x}'}}{(\|\vec{x} - \vec{x}'\|)^3} \mu(\vec{x}') d\Gamma'$$

when \vec{x} approaches Γ from outside Ω , and

$$u_{\Gamma}(\vec{x}) = -2\pi\mu(\vec{x}) + \int_{\Gamma}^{PV} \frac{(\vec{x} - \vec{x}')^T n_{\vec{x}'}}{(\|\vec{x} - \vec{x}'\|)^3} \mu(\vec{x}') d\Gamma'$$

when \vec{x} approaches Γ from inside.

Note 1

Derivation of the Dipole Surface Potential

The following derivation goes through the step-by-step process of deriving the dipole surface potential when \vec{x} approaches Γ from inside.

$$\begin{split} u(\vec{x}) &= \int_{\Gamma} \frac{\partial}{\partial n_{\vec{x}'}} \frac{1}{\|\vec{x} - \vec{x}'\|} \mu(\vec{x}') d\Gamma' \qquad \vec{x} \in \Gamma \\ &= \lim_{\epsilon \to 0} \int_{\Gamma - B(x_0, \epsilon)} \frac{\partial}{\partial n_{\vec{x}'}} \frac{1}{\|\vec{x} - \vec{x}'\|} \mu(\vec{x}') d\Gamma' + \lim_{\epsilon \to 0} \int_{B(x_0, \epsilon)} \frac{\partial}{\partial n_{\vec{x}'}} \frac{1}{\|\vec{x} - \vec{x}'\|} \mu(\vec{x}') d\Gamma' \\ &= \int_{\Gamma}^{PV} \frac{\partial}{\partial n_{\vec{x}'}} \frac{1}{\|\vec{x} - \vec{x}'\|} \mu(\vec{x}') d\Gamma' + \lim_{\epsilon \to 0} \int_{B(x_0, \epsilon)} \hat{n}_{x'} \cdot \nabla \frac{1}{\|\vec{x} - \vec{x}'\|} \mu(\vec{x}') d\Gamma' \\ &= \int_{\Gamma}^{PV} \frac{(\vec{x} - \vec{x}')^T n_{\vec{x}'}}{\|\vec{x} - \vec{x}'\|^3} \mu(\vec{x}') d\Gamma' + \lim_{r \to 0} \int_{0}^{\pi} \int_{0}^{\pi} \hat{n}_{x'} \cdot \nabla \left(\frac{1}{r}\right) \mu r^2 \sin \phi \ d\phi \ d\theta \\ &= 0 + \lim_{r \to 0} \int_{0}^{\pi} \int_{0}^{\pi} \mu \frac{\partial}{\partial r} \left(\frac{1}{r}\right) r^2 \sin \phi \ d\phi \ d\theta \\ &= \lim_{r \to 0} \int_{0}^{\pi} \int_{0}^{\pi} \mu \sin \phi \ d\phi \ d\theta = \left[\int_{0}^{\pi} \mu \cos \phi \ d\theta\right]_{0}^{\pi} = \int_{0}^{\pi} -2\mu \ d\theta \\ &= \left[-2\mu\theta\right]_{0}^{\pi} = -2\pi\mu \end{split}$$

6 Ansatz Formulations

6.1 Dirichlet Problem

6.1.1 Monopole Potential

For an interior or exterior problem:

$$u_{\Gamma}(ec{x}) = \int_{\Gamma} rac{1}{\|ec{x} - ec{x'}\|} \sigma(ec{x'}) d\Gamma'$$



What radiation condition?

The point of the above slide is to show that when using monopole potentials to solve Dirichlet problems on a single domain bounded by γ , the equations are the same for either the interior or the exterior problem.

6.1.2 Dipole Potential

For an exterior problem:

$$u_{\Gamma}(\vec{x}) = 2\pi\mu(\vec{x}) + \int_{\Gamma}^{PV} \frac{(\vec{x} - \vec{x}')^T n_{\vec{x}'}}{(\|\vec{x} - \vec{x}'\|)^3} \mu(\vec{x}') d\Gamma'$$

and for an interior problem:

$$u_{\Gamma}(\vec{x}) = -2\pi\mu(\vec{x}) + \int_{\Gamma}^{PV} \frac{(\vec{x} - \vec{x}')^T n_{\vec{x}'}}{(\|\vec{x} - \vec{x}'\|)^3} \mu(\vec{x}') d\Gamma'$$

Normal points from interior to exterior.



Note that the radiation condition satisfied by the monopole potential is different than the radiation condition satisfied by the dipole potential, even when used to solve the same Dirichlet problem.

6.2 Neumann Problem

6.2.1 Monopole Potential (MP)

Derivative of the monopole potential "jumps" as x passes through Γ , so

$$\frac{\partial}{\partial n_{\vec{x}}} \int_{\Gamma} \frac{1}{\|\vec{x} - \vec{x}'\|} \sigma(\vec{x}') d\Gamma'$$

takes different values just inside and just outside Γ .

6.2.2 MP Int/Ext

For an exterior problem

$$\frac{\partial u_{\Gamma}(\vec{x})}{\partial n_{\vec{x}}} = -2\pi\sigma(\vec{x}') - \int_{\Gamma}^{PV} \frac{(\vec{x}-\vec{x}')^T n_{\vec{x}}}{(\|\vec{x}-\vec{x}'\|)^3} \sigma(\vec{x}') d\Gamma'$$

and for an interior problem

$$\frac{\partial u_{\Gamma}(\vec{x})}{\partial n_{\vec{x}}} = +2\pi\sigma(\vec{x}') - \int_{\Gamma}^{PV} \frac{(\vec{x}-\vec{x}')^T n_{\vec{x}}}{(\|\vec{x}-\vec{x}'\|)^3} \sigma(\vec{x}') d\Gamma'$$

Normal points from interior to exterior.

Note that the signs for the Neumann monopole potential integral equation are different than the signs for the integral equation in the Dirichlet Dipole case. The sign changes are due to the location of the derivative evaluation. In the Neumann Monopole potential case, the derivative is taken with respect to $n_{\vec{x}}$ whereas for the Dirichlet Dipole potential case, the derivative is taken with respect to $n_{\vec{x}'}$.

6.3 Kinds of Equations

First Kind Equations

• Dirichlet Monopole potential integral equation.

Second Kind Equations

- Dirichlet Dipole potential integral equation.
- Neumann Monopole potential integral equation.

Dipole potential for Neumann?

7 Green's Theorem

Laplace's Equation Green's Function

$$\nabla^2 G(\vec{x}) = 4\pi \delta(\vec{x})$$

 $\delta(\vec{x}) \equiv$ impulse in 3-D Defined by its behavior in an integral

$$\int \delta(\vec{x}')f(\vec{x}')d\Omega' = f(0)$$

Not too hard to show

$$G(\vec{x}) = \frac{1}{\|\vec{x}\|}$$

Note 1

Just as an aside, Green's Function may be defined using a different scaling variable depending upon which source one is using. Sometimes Laplace's equation will be written:

$$\nabla^2 G(\vec{x}) = \delta(\vec{x})$$

Where you can see that the value of 4π has been left off. This will simply mean that the Green's Function is now written:

$$G(\vec{x}) = \frac{4\pi}{\|\vec{x}\|}.$$

For our purposes, we will be using the notation in the above slide, and <u>not</u> the notation given in this note.

7.1 Normal Directions

A note here about normal directions is essential. In the above section, the "normal points from interior to exterior" whereas, in the image below, the normal points from inside the domain to outside the domain. How would this impact the solution?



When we go through Green's Theorem in the following section, remember that the normal always points "out" of Ω , as it does in the above figure.

7.2 Divergence Thm

The general Divergence $\vec{T}heorem$:

For any sufficiently smooth \vec{F}

$$\int_{\Omega} \nabla \cdot \vec{F}(x) dV = \int_{\Gamma} \vec{F} \cdot \vec{n}_x dS$$

where Γ is the surface which encloses Ω . Green's theorem follows from the divergence theorem. 4π

7.3 Volume Theorem

If u satisfies Laplace's equation in Ω , then

$$4\pi u(\vec{x}) = \int_{\Gamma} \left[\frac{\frac{\partial u(\vec{x}')}{\partial n}}{\|\vec{x} - \vec{x}'\|} - \frac{\partial}{\partial n_{\vec{x}'}} \frac{1}{\|\vec{x} - \vec{x}'\|} u(\vec{x}') \right] d\Gamma' \quad \vec{x} \in \Omega$$

where the normal points out of Ω .

7.3.1 Surface

Using the Principle Value Integral:

$$2\pi u(\vec{x}) = \int_{\Gamma} \frac{\frac{\partial u(\vec{x}')}{\partial n}}{\|\vec{x} - \vec{x}'\|} d\Gamma' - \int_{\Gamma}^{PV} \frac{\partial}{\partial n_{\vec{x}'}} \frac{1}{\|\vec{x} - \vec{x}'\|} u(\vec{x}') d\Gamma' \qquad \vec{x} \in \Gamma$$

where the normal points out of Ω .

This is one of the cases where it is generally easier to define the normal as pointing out of Ω rather than having the normal point from the interior to the exterior.

7.3.2 Boundary Conditions

The boundary conditions, Dirichlet or Neumann, can be determined by using the surface form of Green's Theorem. For Dirichlet Problems, $u = u_{\Gamma}$ when $\vec{x} \in \Gamma$. So, put the known values for u into Green's Theorem for the surface, and put these known terms on the right hand side, leaving the unknown on the left hand side. Likewise, for Neumann problems, $\frac{\partial u}{\partial n} = \frac{\partial u_{\Gamma}}{\partial n}$ denotes that the derivative of u is known on the boundary when $\vec{x} \in \Gamma$. Again, put the known term on the right hand side, so that the unknown value, u is on the left hand side.

For Dirichlet Problems

$$\int_{\Gamma} \frac{\frac{\partial u(\vec{x}')}{\partial n}}{\|\vec{x} - \vec{x}'\|} d\Gamma' = 2\pi u_{\Gamma}(\vec{x}) + \int_{\Gamma}^{PV} \frac{\partial}{\partial n_{\vec{x}'}} \frac{u_{\Gamma}(\vec{x}')}{\|\vec{x} - \vec{x}'\|} d\Gamma'$$

For Neuman Problems

$$2\pi u(\vec{x}) + \int_{\Gamma}^{PV} \frac{\partial}{\partial n_{\vec{x}'}} \frac{1}{\|\vec{x} - \vec{x}'\|} u(\vec{x}') \ d\Gamma' = \int_{\Gamma} \frac{\frac{\partial u_{\Gamma}(\vec{x}')}{\partial n}}{\|\vec{x} - \vec{x}'\|} d\Gamma'$$

where normal points out of Ω (interior or exterior!)

7.4 Overview



8 Summary

Laplace Problems Exterior Radiation Condition Potential Representations Monopole and Dipole potentials Principle Value Ansatz and Green's Theorem Dirichlet and Neumann problems First and Second Kind Equations.

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1.5 First and Second Kind Theory, part 2

Numerical Methods for PDEs

Boundary Element Methods, Lecture 5 First and Second Kind Theory, part 2

Notes by L. Proctor, C. Coelho and J. White

December 8, 2008

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1 Outline

Interior Neumann Use Nystrom to Solve Look at 2-D Problems Fredholm Alternative Connection to Linear Algebra First Kind Convergence Theory

2 Exterior Formulations

2.1 Dirichlet Problem

2.1.1 Monopole Potential

For an 3D exterior problem:

$$u_{\Gamma}(\vec{x}) = \int_{\Gamma} \frac{1}{\|\vec{x} - \vec{x}'\|} \sigma(\vec{x}') d\Gamma'$$

For an 2D exterior problem:

$$u_{\Gamma}(\vec{x}) = \int_{\Gamma} \log \|\vec{x} - \vec{x}'\| \sigma(\vec{x}') d\Gamma'$$

2.1.2 Dipole Potential

For a 3-D exterior problem:

$$u_{\Gamma}(\vec{x}) = 2\pi\mu(\vec{x}) + \int_{\Gamma}^{PV} \frac{(\vec{x} - \vec{x}')^T n_{\vec{x}'}}{(\|\vec{x} - \vec{x}'\|)^3} \mu(\vec{x}') d\Gamma'$$

For a 2-D exterior problem:

$$u_{\Gamma}(\vec{x}) = \pi \mu(\vec{x}) + \int_{\Gamma}^{PV} \frac{(\vec{x} - \vec{x}')^T n_{\vec{x}'}}{(\|\vec{x} - \vec{x}'\|)^2} \mu(\vec{x}') d\Gamma'$$

Normal points from interior to exterior.

2.2 Neumann Problem

2.2.1 Monopole

For an exterior 3D problem

$$\frac{\partial u_{\Gamma}(\vec{x})}{\partial n_{\vec{x}}} = -2\pi\sigma(\vec{x}') - \int_{\Gamma}^{PV} \frac{(\vec{x} - \vec{x}')^T n_{\vec{x}}}{(\|\vec{x} - \vec{x}'\|)^3} \sigma(\vec{x}') d\Gamma'$$

For an exterior 2D problem

$$\frac{\partial u_{\Gamma}(\vec{x})}{\partial n_{\vec{x}}} = -\pi\sigma(\vec{x}') - \int_{\Gamma}^{PV} \frac{(\vec{x} - \vec{x}')^T n_{\vec{x}}}{(\|\vec{x} - \vec{x}'\|)^2} \sigma(\vec{x}') d\Gamma'$$

Normal points from interior to exterior.

3 Interior Example

3.1 3D Case

3.1.1 Monopole Potential

Surface Potential

$$u_{\Gamma}(\vec{x}) = \int_{\Gamma} \frac{1}{\|\vec{x} - \vec{x}'\|} \sigma(\vec{x}') d\Gamma' \qquad \vec{x} \in \Gamma$$

Surface Normal Derivative

$$\frac{\partial u_{\Gamma}(\vec{x})}{\partial n_{\vec{x}}} = \frac{\partial}{\partial n_{\vec{x}}} \int_{\Gamma} \frac{1}{\|\vec{x} - \vec{x}'\|} \sigma(\vec{x}') d\Gamma' \qquad \vec{x} \in \Gamma$$

Normal points to exterior.

3.1.2 Interior Neumann

Monopole Potential Using the P.V. Integral

$$\frac{\partial u_{\Gamma}(\vec{x})}{\partial n_{\vec{x}}} = 2\pi\sigma(\vec{x}) + \int_{\Gamma}^{PV} \frac{\partial}{\partial n_{\vec{x}}} \frac{1}{\|\vec{x} - \vec{x}'\|} \sigma(\vec{x}') d\Gamma'$$

looks 2^{nd} Kind Equation, Try Nystrom.



3.1.3 Nystrom Method

Set quadrature points = collocation points

$$\frac{\partial u_{\Gamma}(\vec{x}_1)}{\partial n_{\vec{x}}} = 2\pi\sigma_{n1} + \sum_{10\vec{j}=1}^n w_{1,j} \frac{\partial}{\partial n_{\vec{x}}} \frac{1}{\|\vec{x}_1 - \vec{x}_j\|} \sigma_{nj}$$

$$\frac{\partial u_{\Gamma}(\vec{x}_n)}{\partial n_{\vec{x}}} = 2\pi\sigma_{nn} + \sum_{j=1}^n w_{n,j} \frac{\partial}{\partial n_{\vec{x}}} \frac{1}{\|\vec{x}_n - \vec{x}_j\|} \sigma_{nj}$$

.

n equations in n unknowns

$$j = i$$
 case (self-term)?

3.1.4 i = j

For the monopole 3-D Neumann Formulation,

$$G(\vec{x}, \vec{x}') = \frac{\partial}{\partial n_{\vec{x}}} \frac{1}{\|\vec{x} - \vec{x}'\|} = -\frac{(\vec{x} - \vec{x}')^T n_{\vec{x}}}{(\|\vec{x} - \vec{x}'\|)^3}$$

PROBLEM: $G(\vec{x}, \vec{x}')$ blows up as $\vec{x} \to \vec{x}'$.

3.2 2-D Case

Monopole Neumann Formulation

$$\frac{\partial u_{\Gamma}(\vec{x})}{\partial n_{\vec{x}}} = \pi \sigma(\vec{x}) + \int_{\Gamma}^{PV} \frac{\partial}{\partial n_{\vec{x}}} \log \|\vec{x} - \vec{x}'\| \sigma(\vec{x}') d\Gamma'$$

Simplifying the Green's function,

$$G(\vec{x}, \vec{x}') = \frac{\partial}{\partial n_{\vec{x}}} \log \|\vec{x} - \vec{x}'\| = -\frac{(\vec{x} - \vec{x}')^T n_{\vec{x}}}{(\|\vec{x} - \vec{x}'\|)^2}$$

3.2.1 Smooth Γ

 $G(\vec{x}, \vec{x}')$ finite as $\vec{x} \to \vec{x}'$ if Γ smooth.



3.2.2 Nonsmooth Γ

 $G(\vec{x}, \vec{x}')$ not finite as $\vec{x} \to \vec{x}'$ if x' is on a corner.


3.2.3 Disk Example



Note uniform quadrature weights on the circle. Resulting matrix is singular! Why?

4 Second Kind Theorem

4.1 Theorem

Given

$$(I+K)\sigma = \Psi$$
 (Integral Eqn.)
 $(I+K_n)\sigma_n = \Psi_n$ (Discretized Eqn.)
AND
 $\|(I+K)^{-1}\| < C$ Unique solvability

If

$$\lim_{n\to\infty} \|(K-K_n)\| \to 0 \text{ and } \|\Psi-\Psi_n\| \to 0$$

Then

$$\lim_{n \to \infty} \|\sigma - \sigma_n\| \to 0$$

4.1.1 Scaled Example

Define Scaled Variables

$$\Psi \equiv \frac{1}{\pi} \frac{\partial u_{\Gamma}(\vec{x})}{\partial n_{\vec{x}}}$$
$$K \equiv \frac{1}{\pi} \int_{\Gamma}^{PV} \frac{\partial}{\partial n_{\vec{x}}} \log \|\vec{x} - \vec{x}'\| \sigma(\vec{x}') d\Gamma'$$

The 2-D Neumann problem becomes

$$(I+K)\sigma = \Psi$$

4.1.2 Key Property

Main assumption of second kind theory: $(I+K)^{-1} \text{ is bounded}.$

Is $(I + K)^{-1}$ bounded for Interior Neumann Problem?

4.2 Linear Algebra

Given Ax = b, $A \in \Re^{n \times n}$, $x, b \in \Re^n$ A^{-1} exists and is bounded iff Ay = 0 implies y = 0 (no null space) If Ay = 0 for $y \neq 0$ then either Ax = b has an infinite # of solutions Ax = b then $A(x + \alpha y) = b$ OR

> Ax = b does not have a solution b is not in the column space of A

4.3 3-D Null Space

Consider $\tilde{\sigma}$ defined by

$$u_{\Gamma}(\vec{x}) = 1 = \int_{\Gamma} \frac{1}{\|\vec{x} - \vec{x}'\|} \tilde{\sigma}(\vec{x}') d\Gamma' \qquad \vec{x} \in \Gamma$$

Then

$$\frac{\partial u_{\Gamma}(\vec{x})}{\partial n_{\vec{x}}} = 0 = 2\pi\tilde{\sigma}(\vec{x}) + \int_{\Gamma}^{PV} \frac{\partial}{\partial n_{\vec{x}}} \frac{1}{\|\vec{x} - \vec{x'}\|} \tilde{\sigma}(\vec{x'}) d\Gamma'$$

 $\tilde{\sigma}$ is in the Null space of I + K

 $(I+K)^{-1}$ is not bounded!!

4.4 Fredholm Alternative

General Theorem

For I + K either

$$(I+K)\sigma = \Psi$$
 has an infinite $\#$ of solutions
OR

 $(I+K)\sigma = \Psi$ has no solution

4.4.1 2D Example

Scaled Equations:

$$\frac{1}{\pi} \frac{\partial u_{\Gamma}(\vec{x})}{\partial n_{\vec{x}}} = \sigma(\vec{x}) + \frac{1}{\pi} \int_{\Gamma}^{PV} \frac{\partial}{\partial n_{\vec{x}}} \log \|\vec{x} - \vec{x}'\| \sigma(\vec{x}') d\Gamma'$$

For a solution to exist

$$\int_{\Gamma} \frac{\partial u_{\Gamma}(\vec{x})}{\partial n_{\vec{x}}} d\Gamma = 0$$

4.4.2 Fixes

Add a point constraint Fix u at some point Force σ orthogonal to null space Need the null space May need to solve 1^{st} kind equation Use SVD to solve singular system Can be computationally expensive

5 1st Kind Convergence

Three-dimensional Laplace's equation

• Unknowns might be physically meaningful.

$$u(x) = \int \frac{1}{\|\vec{x} - \vec{x}'\|} \quad \underbrace{\sigma(\vec{x}')}_{\text{charge density}} \quad dS'$$

• Might match boundary conditions

– Dirichlet and $\frac{1}{R}$ radiation condition

5.1 Nonsingular Green's Function

Denote the integral operator as K

$$K\sigma \equiv \int_{-1}^{1} |x - x'|\sigma(x')dS' \Rightarrow K\sigma = \Psi$$

The integral operator is singular : K has a null space

$$K\sigma_{0} = \int_{-1}^{1} |x - x'| \sigma_{0}(x') dS' = 0$$
(2)

 $K\sigma^a = \Psi$ then $K(\sigma^a + \sigma_0) = \Psi$ If

The Singular Kernel 5.2



5.3**Convergence** Analysis

Partial Differential Equation form:

 $\nabla^2 u = f$ in Ω u = 0on Γ

 Ω is the volume domain Γ is the problem surface

"Nearly" Equivalent weak form

$$\underbrace{\int_{\Omega} \nabla u \nabla v dx}_{a(u,v)} = \underbrace{\int_{\Omega} f v dx}_{l(v)} \qquad \forall v \in H^{1}(\Omega)$$

Introduced an abstract notation for the equation, u must satisfy:

$$a(u, v) = l(v) \qquad \forall v \in H^1(\Omega)$$

Introduce an approximate solution $u^n = \sum_{i=1}^n \alpha_i \varphi_i$ $\Rightarrow u^n$ is a weighted sum of basis functions

The basis functions define a space

$$X_n = \left\{ v \in X_n | v = \sum_{i=1}^n \beta_i \varphi_i \quad \text{for some } \beta_i' \mathbf{s} \right\}$$



Key Idea

a(u, u) defines a norm on $H_0^1(\Omega)$ $a(u, u) \equiv |||u|||$ u is restricted to be 0 at 0 & 1! Using the norm properties, it is possible to show

If $a(u^n, \varphi_i) = l(\varphi_i)$	$\forall \varphi_i \in \{\varphi_1, \varphi_2, \dots, \varphi_n\}$
Then $ u - u^n =$	$\min_{w \in X} \ u - w^n \ $
Solution	Projection
Error	Error

5.3.1 Optimality Result



How well can you fit the exact solution with a member of X_n ? You must measure the error in the ||| ||| norm

5.3.2 Sobolov Space

"Weak" Form for the integral equation $\underbrace{\iint_{\Gamma} v(x) \frac{1}{\|x - x'\|} \sigma(x') dS' dS}_{a(\sigma, v)} = \underbrace{\int_{\Gamma} v(x) \Psi(x) dS}_{l(v)} \quad \forall v \in \bar{H}(\Gamma)$ The difficulty is defining $\bar{H}(\Gamma)$ with right properties Must exclude $\sigma(x)$'s where $\int \frac{1}{\|x - x'\|} \sigma(x') dS' = 0$ $\bar{H}(\Gamma)$ is a fractional Sobolev Space

We won't say more about this!

5.3.3 Use FEM Key Idea

$$a(\sigma, \sigma) \text{ defines a norm on } \bar{H}(\Gamma) \qquad a(\sigma, \sigma) = \||\sigma|\|$$

$$\sigma^n = \sum_{i=1}^n \alpha_i \quad \underbrace{\varphi_i(x)}_{\text{Basis}} \qquad X_n = \left\{ v \in X_n | v = \sum_{i=1}^n \beta_i \varphi_i \quad \text{for some } \beta'_i s \right\}$$

Functions

5.3.4 FEM Idea Cont.

Using the norm properties, it is possible to show

If

$$a(u^n,\varphi_i) = l(\varphi_i) \qquad \forall \varphi_i \in \{\varphi_1,\varphi_2,\ldots,\varphi_n\}$$

Then	$\underbrace{\ u-u^n\ }$	$= \min_{w_n \in X_n}$	$\underbrace{\ u-w^n\ }$
	Solution		Projection
	Error		Error

6 Summary

Interior Neumann Use Nystrom to Solve Look at 2-D Problems Fredholm Alternative Connection to Linear Algebra First Kind Convergence Theory

Mostly Waved hands.

1.6 Fast Algorithms for Integral Equation Methods

Numerical Methods for PDEs

Boundary Element Methods, Lecture 6 Fast Algorithms for Integral Equations

L. Proctor, S. De, K. Nabors, J. Phillips, B. Buchmann, & J. White

December 10, 2008

1 Outline

Reasons for Fast Solvers Collocation System Reminder Fast Solver General Approach Using Iterative methods Fast matrix-vector products Fast Multipole Algorithms Precorrected-FFT Algorithms

2 Background

2.1 Discretize Surface Into Panels

2.1.1 Piecewise Constant Basis

Integral Equation : $\Psi(x) = \int_{surface} \frac{\sigma(x')}{\|x-x'\|} dS'$



2.1.2 Centroid Collocation

Put collocation points at panel centroids



2.2 Dense Matrix

2.2.1 Resultant Dense Matrix

Matrix Entries Are Never Zero

$$A_{i,j} = \int_{panel_j} \frac{1}{\|x - x'\|} dS'$$

Distant Elements Decay Slowly

$$\propto \frac{1}{\|x - x'\|}$$

Too Slow To Ignore.

2.2.2 Complicated Examples



Need More than 100,000 unknowns!!

Need 100 Gigabytes to Store Matrix.

2.2.3 Gaussian Elimination



 n^3 – Too Expensive!

3 Iterative Methods

3.1 Electrostatics Application

General Iterative "Algorithm"

0 : Guess at panel charges $\vec{\alpha}$

1 : Compute the centroid potentials from the charges ${\bf A}\vec{\alpha}$

- 2 : Compare the computed to known potentials $\mathbf{R}=\boldsymbol{\Psi}-\mathbf{A}\vec{\alpha}$
- 3 : Fix the panel charges, go to Step 1. $\vec{\alpha}$

3.2 Conjugate Gradient (CG)

Conjugate Gradient (CG) Methods are iterative methods useful for solving systems of equations involving symmetric matrices $\mathbf{A} = \mathbf{A}^T$. The rate of iteration convergence for CG can be related to the ratio of the maximum to the minimum eigenvalue of A.

3.2.1 CG for 2nd Kind



Eigenvalues for 2^{nd} Kind Integral Equation

3.2.2 CG for 2nd Kind Cont.

Conjugate-Gradient convergence rate

$$\left\| r^{k} \right\| \leq 2 \left(\sqrt{\frac{\lambda_{\max}}{\lambda_{\min}}} - 1 \left/ \sqrt{\frac{\lambda_{\max}}{\lambda_{\min}}} + 1 \right)^{k} \left\| r^{0} \right\|$$

For discretized Second Kind equations

$$\frac{\lambda_{max}}{\lambda_{min}}$$
 is bounded independent of n

Number of CG iterations independent of n!!

3.2.3 Steps of CG

The k^{th} step of the Conjugate Gradient Algorithm



3.2.4 Cost of CG

Complexity of the Conjugate Gradient Method

compute Ap_k	Dense Matrix-vector product costs O(n ²)
$\alpha_{k} = \frac{(r^{*}) (Ap_{k})}{(Ap_{k})^{T} (Ap_{k})} \qquad \qquad \text{Vect}$	or inner products, O(n)
$x^{k-1} = x^k + \alpha_k p_k$ $r^{k+1} = r^k - \alpha_k A p_k$	Vector Adds, O(n)
$p_{k+1} = r^{k+1} - \frac{\left(Ar^{k+1}\right)^T \left(Ap_k\right)}{\left(Ap_k\right)^T \left(Ap_k\right)} p_k$	Inner products, total cost O(n)

Algorithm is $O(n^2)$ for integral equations even though # of iterations, k, is small!

3.2.5 Accelerate CG?

Accelerate the Conjugate Gradient Method Exactly compute Ap_k

Dense matrix-vector (M-V) product costs $O(n^2)$ Approximately compute $\mathbf{A}p_k$

Reduces M-V product costs to O(n) or $O(n \log n)$ Need a fast approximation for matrix-vector products

4 Fast Solvers

4.1 Direct Computation



Physical interpretation:
 Ap = N "potentials" due to N charges.

• $O(N^2)$ if done naively

1D Strip of Charge in 3D Space 4.2

Simplification of the A Matrix 4.2.1

-

1-D Strip of Charge in 3-D Space

•	•	•	•	•	•	•	•
$\begin{bmatrix} A_{11} \\ A_{21} \\ \vdots \\ A_{81} \end{bmatrix}$	$\begin{array}{c} A_{12} \\ A_{22} \\ \vdots \\ A_{82} \end{array}$	 	$\begin{array}{c} A_{18} \\ A_{28} \\ \vdots \\ A_{88} \end{array} \right]$		$\begin{bmatrix} \alpha_1 \\ \alpha_2 \\ \vdots \\ \alpha_8 \end{bmatrix}$	=	$\left[\begin{array}{c} \Psi_1\\ \Psi_2\\ \vdots\\ \Psi_8\end{array}\right]$

What can one say about the A matrix?

4.2.2 Properties of A.

The A matrix is:

- Symmetric Panel i exerts the exact same charge on j that j exerts on i
- All the Diagonal Values are the Same

$$A_{ii} = \int_{\text{panel}_i} \frac{1}{\|\vec{x}' - \vec{x}_{c_i}\|}$$

• Each Superdiagonal & Subdiagonal Element is Equal along Its Own Diagonal as Well

4.2.3More Properties of A.

How many unique entry values are there in A?



Note 1

The above matrix is one of a class of matrices called Toeplitz Matrices. In a Toeplitz matrix, the matrix entries along any diagonal have the same value, but the different diagonals can have different values. For this reason, an $n \times n$ Toeplitz matrix has only 2n - 1 distinct values. A particular important special case of Toeplitz matrices are Circulant Matrices. Circulant matrices are "periodized" Toeplitz matrices in that the first super-diagonal has the same values as the n - 1 sub-diagonal, the second super-diagonal has the same values as the n - 2 sub-diagonal, etc. Circulatant matrices are diagonalized by the discrete Fourier transform, a property we will use in the section covering the Precorrected-FFT methods.

4.2.4 Geometric Simplification

Approximate (by grouping) the elements that are a "reasonable distance" away from the element which you are evaluating



4.3 Fast Potential Concept



- Decompose potential into short- and long- range.
- Approximate long-range part of potential.
- Sum short-range part in normal manner
- Multilevel decomposition for "O(N)" algorithm

4.4 Computational Costs



- Gaussian Elimination: $O(n^3)$ time, $O(n^2)$ memory
- Iterative with direct M-V: $O(n^2)$ time, $O(n^2)$ memory
- Fast Methods: O(n) time, O(n) memory

5 Multipole Algorithms

5.1 Direct Potential Evaluation



• Complete evaluation at d points costs d^2 operations.

5.2 Multipole Representation

5.2.1 1D Strip in 3D Space



The cost of forming clusters is, in general, $O(n + \frac{n}{2} + \frac{n}{4} + \frac{n}{8} + ...) \approx O(n)$

What is the cost of estimating the evaluation point potential?



The cost of gathering clusters is $O(n \log n)$

5.2.2 Computational Example

A few multipoles (monopoles, dipoles, quadrapoles, etc) can accurately represent the potential due to a cluster of charges, with the accuracy improving with increasing distance from the cluster increases. For example, if one is very far from the cluster, the potential due to the cluster will be nearly identical to the potential of a point charge whose location is at the center of the cluster and whose value is the sum of the cluster charges. The accuracy of such a monopole representation can be improved by adding dipole, quadrapole and higher order multipoles. Note, however, that higher order multipoles generate potentials that depend on the multipole's orientation, and that must be considered.



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5.2.3 General Case



• Approximate potential at point *i*:

$$v_i(r_i, \phi_i, \theta_i) \approx \sum_{j=0}^{order} \sum_{k=-j}^j \frac{M_j^k}{r_i^{j+1}} Y_j^k(\phi_i, \theta_i)$$

• Multipole coefficients function of panel charges:

$$M_j^k \stackrel{\triangle}{=} \sum_{i=1}^d \frac{q_i}{A_i} \int_{\text{panel } i} \rho^j Y_j^{-k}(\alpha, \beta) dA.$$

- Computing Multipole expansions costs order d operations.
- Each approximate potential evaluation costs order 1 operations.

d potential evaluation due to d panels in order d operations

5.3 Error Scale Invariance



 $\operatorname{Error} \le K \left(\frac{R}{r}\right)^{order+1}$

Error $\leq K \left(\frac{3R}{3r}\right)^{order+1}$

5.4 Multipole Algorithm Hierarchy



Hierarchy guarantees: • Bounded error:

Error
$$\leq K \left(\frac{R}{r}\right)^{order+1}$$

 $\leq K \left(\frac{1}{2}\right)^{order+1}$

order = 2 yields one percent accuracy.

• Order n ops for n potentials.

5.5 Local Representations

5.5.1 Cost Reduction



Construct a local expansion to represent distant charge potentials.
Evaluate a single local expansion, rather than many multipole expansions, at each evaluation point.

5.5.2 Clustered Evaluations



- Local expansion summarizes the influence of distant charge for clusters of evaluation points.
- Gives O(n) potential evaluation when combined with coalescing of charge done by multipole expansions.
- Approximate potential at point *i*: $v_i(r_i, \phi_i, \theta_i) \approx \sum_{j=0}^{order} \sum_{k=-j}^{j} L_j^k Y_j^k(\phi_i, \theta_i) r_i^j$.

5.5.3 Summary of Operations



• Multipole and local expansions are built using complementary hierarchies.

- Complete calculation consists of:
 - 1. Build multipoles (Upward Pass).
 - 2. Build locals (Downward Pass).
 - 3. Evaluate local expansions and nearby charge potential (Evaluation Pass).

5.5.4 Hierarchy Construction



- First build the multipole expansions moving upward from child to parent.
- Then build the local expansions by moving downward from parent to child.
- Computation has a tree structure.

5.5.5 Construction Details

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- Conversion of multipole expansions to local expansions.
- A child's local expansion is its parents local expansion plus conversions of multipole expansions in child's interaction range.

5.6 Adaptive Algorithm

5.6.1 Multipole Inefficiency



$$v_4(x, y, z) = q_1 P_{41} + q_2 P_{42} + q_3 P_{43}$$

Multipole Evaluation



$$v_4(x, y, z) \approx \bar{M}_0^0 \frac{1}{r} + \bar{M}_1^0 \frac{z}{r^3} - \bar{M}_1^1 \frac{x}{2r^3} - \tilde{M}_1^1 \frac{y}{2r^3}$$

Using Multipole MORE expensive than Direct.

5.6.2 Simple Adaptive Scheme

If there are fewer panels than multipole coefficients, calculate the panels' influence directly.

- Similarly, local expansions are not used if there are fewer evaluation points than local expansion coefficients.
- Retains O(mn) complexity for nonuniform panel distributions.

- 5.7 Computational Examples
- 5.7.1 Sphere Potential Distribution



- Potential given by $\psi(x) = -\frac{x_3}{2||x||^3}$.
- Charge given by $\sigma(x) = \frac{-3}{8\pi}x_3$.



- Error should decay like $\frac{1}{n}$.
- Multipole approximations eventually interfere.
- Higher-order multipole expansions needed for higher accuracy.

5.7.2 Two Sphere Example



- Potential on each sphere: $\psi(x) = -\frac{x_3}{2||x||^3}$.
- Does not correspond to a simple physical problem.



- Direct matrix-vector product cost increases like n^2 .
- Multipole matrix-vector product cost increases like n.
- The slope for the multipole algorithm depends on accuracy.
- For order 2 expansions, breakpoint is about n = 400.

5.8 Complexity Summary

For an integral equation discretized with n panels:

- Gaussian elimination: $O(n^3)$.
- Iterative Matrix Solution, direct M-V $O(n^2)$.
- Multipole accelerated Iterative method O(n).

6 Precorrected-FFT

6.1 Introduction

Strip of Charge in Space

Bring the ends of the strip of charge together to form a ring.



Flattening the ring leads to the figure shown below on the left. Forming a ring from the strip of charges results in a system of equations with even more structure that the Toeplitz matrix system described above. The matrix in the ring case will be *circulent*.



The above circulant matrix is the matrix representation of periodic convolution. This convolutional structure is partly due to the homogenity of the geometry, and partly because the Green's function is translation invariant. The Green's function for Laplace's equation, G(x, x'), is translationally invariant because G(x, x') only depends on the difference, x - x'. As mentioned above, the discrete Fourier transform diagonalizes circulent matrices, and therefore circulent matrices can be inverted in $n \log n$ time using the fast Fourier transform.

6.2 Algorithm Outline



• Finer Discretizations Usually Yield Finer Grids.

- 1. Project panel charges on grid
- 2. Calculate grid-charge potentials o
- 3. Interpolate grid potentials onto **p**
- 4. Local corrections [compute nearby interactions dire

6.3 Algorithm Analysis

6.3.1 Interpolation and Projection



Approximate potential Ψ at x due to charge at y by *interpolating* potential using points and weights x_i, w_i

Interpolate: potential at x due to unit charge at y

$$\Psi(x|y) \simeq \hat{\Psi}(x|y) = \sum w_i g(x_i, y)$$

Anterpolate: potential at y due to unit charge at x

$$\Psi(y|x) \simeq \hat{\Psi}(y|x) = \sum w_i g(y, x_i)$$

So

$$\hat{\Psi}(y|x) = \hat{\Psi}(x|y)$$

Same as representing charge at x with w_i and evaluating at y



Equivalent conditions:

- Approx Potential in cell due to charge at large R.
- Approx Potential at large R due to charge in cell.
- Cost is O(N)



• Let H be grid charge-potential mapping

$$H:q_g\to\Psi_g$$

- H is Toeplitz
- Embed H in circulant matrix

$$\begin{bmatrix} \psi_g \\ x \end{bmatrix} = \begin{bmatrix} H & X \\ X & X \end{bmatrix} \begin{bmatrix} q_g \\ 0 \end{bmatrix}$$

- Use FFT for matrix multiply Must Have Translation Invariance
- Cost $O(M \log_2 M), M = \#$ cells

6.3.3 Nearby Interactions



6.3.4 Inhomogeneity Problem

K			4.3
			1.

• Empty Grid due to FFT - Inefficient

6.4 Examples





6.5 PFFT vs. Multipole

• Comparisons: PFFT p = 3 to Multi l = 2

		SLIDE 49			
Example	CPU	Memory	Product	Error	
via	0.61	0.37	0.23	0.18	
woven5x5	0.45	0.48	0.22	0.09	
cube	0.38	0.32	0.12	0.12	
bus3x8	0.27	0.27	0.07	0.01	
SRAM	0.39	0.43	0.17	0.07	
mean	0.42	0.37	0.16	0.09	

• Faster with $10 \times$ better accuracy !

6.6 PFFT vs. direct

6.6.1 Memory

	Example	Memory Usage		
Name	Panels[conductors]	P/FFT	Direct	
via	6120[4]	21 Mb	(286 Mb)	
woven5x5	9360[10]	$50 { m ~Mb}$	(668 Mb)	
woven15	82080[30]	$246 { m ~Mb}$	(50.2 Gb)	
cube	126150[1]	$225 { m ~Mb}$	(119 Gb)	

6.6.2 Time

Example	CPU Usage				
Name	P/FFT	Dir. Iter.	Gauss. Elim.		
via	1.1 min	(5.6 min)	(1.9 hrs)		
woven5x5	5.2 min	(42 min)	(6.9 hrs)		
woven15	1.7 hrs	(11.5 days)	(194 days)		
cube	3.3 min	$(8.4 \mathrm{hrs})$	(2.7 yrs)		

7 Summary

Reasons for Fast Solvers Collocation System Reminder Fast Solver General Approach Using Iterative methods Fast matrix-vector products Two Fast Methods Fast Multipole - Multiresolution Precorrected-FFT - Translation Invariance

Chapter 2

Foundations of Algorithms and Computational Techniques in Systems Biology

Engineering has always played a role in biology, specifically in the past couple of decades the field of computational biology has emerged and contributed greatly. Foundations of Algorithms and Computational Techniques in Systems Biology is a course that gives an overview of topics of interest to a computational biologist. The course covers protein modeling, modeling networks, and image processing. These are the top three areas in computational biology, and this course shows how one may use computational techniques to solve various problems with a biological application. This is very interesting to both the biologist and the computer scientist.

2.1 Motivation/Overview



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PHYSICAL, CHEMICAL, & BIOLOGICAL MODELING OF PROTEINS

Proteins:

· biological polymers of about 20 amino acids

polymers are any kind of large molecules made of repeating identical or similar subunits called monomers

- · "perfect" homogeneous, pure synthesis
- around 10k copies in a cell
- linear, unbranched chains of a <u>unique sequence</u>
- · generally fold to characteristic structure with no additional information

sequence <i>folding</i> (1D)	structure $(3D) \longrightarrow$	chemical functions	$\xrightarrow{\text{biological}} $	network functions
protein ↑ mRNA ↑ genome (DNA)	x-ray crystallography NMR	binding catalysis	synthesis/ degradation energy storage/ utilization gene expression development immune surveillance	control points – decision "robustness" time keepers oscillators important area of growth

Why Model?

· Understanding : model facilitates development of understanding reason for properties

- mechanistic basis for function
- disease
- Prediction
 - experiment planning
 - validate a model or select among models

• Design

- perturbation : improve properties
- intervention : repair

2

Models of Proteins 2.2

MIT 6.581/BE.482

9 February 2006 Thursday

FOUNDATIONS OF ALGORITHMS AND COMPUTATIONAL TECHNIQUES IN SYSTEMS BIOLOGY Spring 2006





(1) size: small - large $R_{Gly}: -H \longrightarrow R_{Trp}$ hydrophobic polar (2) polarity: charged ⊕_NH₂ R_{Leu}: R_{Asn} : R_{Arg} : -NH2 NH

(3) uniformity of character

(4) local backbone flexibility Gly Pro (flexible) (rigidity)

Coordinate systems:

1) Absolute Cartesian Coordinates



Think of the molecules as graphs where



- · atoms appear to have a fixed spherical size & approach to contact neighbors
- · complementary electrostatics





 $U_{\text{NON-COVALENT}} = \sum_{i>j} \left(\frac{B_{ij}}{r_{ij\downarrow}^{12}} - \frac{C_{ij}}{r_{ij\downarrow}^{6}} \right)$ $\sum_{i > j} \frac{q_i q_j}{\varepsilon r_{ij}}$ Electrostatics →Coulombic van der Waals →Lennard-Jones Uvdw 1 electron cloud repulsion dipole dipole interaction $> r_{ij}$

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2.3 Discrete Conformational Search



This method is provably correct, but can't prove in mining time Approach:

Aproan. 1. Eliminate iteratively singles until no more possible 2. Develop higher-order eliminations (pairs) _____ 3. when all eliminations done, enumerate remaining space Application:

Insulin: 76 positions: 2.7×10^{36} conformations Iterative approach (9 iterations) $2.7 \times 10^{36} \longrightarrow 7200 \longrightarrow$ search by enumeration 93% of buried positions "correct"

2.4 Binding and Docking



Simulation







2.6 Molecular Dynamics and Electrostatics

Q(1) operations to compute



tar each of N evaluation atoms Fi=∑ Fij

2.7 Continuum Electrostatic Modeling I





Polsson Eqn: 024=-420 plug in the Faurier Representation on bothsides omatch Fourier terms....

$$\overline{\underline{U}}[m_s, m_y, m_z] = \underbrace{\underbrace{-\frac{4\pi}{2\pi y^* m_z^*}}}_{(scalar)} \overline{\overline{e}}[m_s, m_y, m_z]$$

-(sume Fourier components) to solve periodic problem: - take charge density - complete Fourier Series - easily complete Fourier Series for periodic potential







sign



+ 2 R $\Delta G_{binding}^{elec} = \tilde{q}_{1}^{T} L \tilde{q}_{1} + 2 \tilde{q}_{1}^{T} C \tilde{q}_{R} + \tilde{q}_{R}^{T} R \tilde{q}_{R}$ ligand interactions receptor desolvation recovered desolvation

 $\overline{\nabla}_{q_{L}}^{e} \left(\Delta G_{\text{binding}}^{\text{elec}} \right) = 2L \overline{q}_{L} + 2G \overline{q}_{R} = \overline{0}$ $\Rightarrow L \overline{q}_{L} = C \overline{q}_{R}$ ⇒ gr = -E'Cq

JPhys B 105: 880-888 (2001). Chorismate mutase

JACS 125: 55913-9 (2003). 10-fold improvement in binding





2.11 Statistical Mechanics







2.12 Statistical Mechanics

2.13 Formulating Models

6.581/BE-482	
LECTURE 14: FORMALATING MODELS 4 APRIL 2006	
DIFFERENTIAL EQUATION MODEL	Dimerization case: k.[A][A]= k=[A]
	$[A]^2 = \frac{k_4}{k_4} [A_2]$
CELL	$[A] = \sqrt{\frac{1}{5}} \left[A_2 \right]^{\frac{1}{2}}$
[X;]"" [X] SURFACE	Also given:
(Exj1 Nucleus)	$A(0) = A_0 A_2(0) = 0$
([X _k]//	Steady-State:
	[A]+2[A2]=A0
	CHEMICAL KINETICS FORM
VERY GENERAL SETTING	$\frac{d}{dt}[x_i] = \sum_{j=1}^{t} \gamma_j^{t} [x_j] destination$
	(Typical Case vi <0)
	(degradation process)
$\chi = 1$	
time-variation of '	$f_{\mathbf{x}}^{\mathbf{z}}[\mathbf{x}_{i}] = \sum_{i=1}^{N} \mathbf{x}_{i}^{\mathbf{z}}[\mathbf{x}_{i}] + \sum_{i=1}^{N} \mathbf{x}_{i}^{\mathbf{z}}[\mathbf{x}_{i}]\mathbf{x}_{i}$
きなる くう ないしょう	jai J jai kei k
rate of rate of	Sile<0
production consumption	Typical Production Process
EX: A+B	
4[c]=x [1][8] x [c]	$ = \sum_{i=1}^{n} \gamma_i^{(X_i)} + \sum_{j=1}^{n} \gamma_j^{(X_i)} + \sum_{j=1}^{n} \gamma$
4[1]= K.[2]-K.[2]-K.[2]	4 5 5 AF TU- TU. 1
4[B]=K.[C]-K.[A][B]	j=1 K=1 PjikLujJLukJ
+ within conditions	MWITPLE INPUTS
A(0), B(0), C(0)	$\mathcal{X}_{1}^{\prime} < 0 \Rightarrow$ models degradation \mathcal{Y}_{1}^{\prime}
Ex: Dimerization	i ≠ 0 ⇒ models conversion x; γ ∧j
$A \stackrel{\mathbb{N}}{\underset{K_2}{\longrightarrow}} A_2$	χ_{jk} $\chi_{j} + \chi_{k} = \frac{\partial_{jk}}{\partial d} \chi_{k}$
$d_{t}^{d}[A_{2}] = K_{1}[A][A] - K_{2}[A_{2}]$	It/[xi] is involved only in
$\frac{1}{4}[A] = 2K_2[A_2] - 2K_1[A][A]$	reactions that consume $[x_i] \Rightarrow \gamma_i \leq 0, \gamma_{i,k} \leq 0, \alpha_{j,i} \leq 0$
STEADY-STATE (Equilibrium)	Exiles at $t=0$
Keactions are in Dalance	Then All $[X_i]$'s ≥ 0 for all t $[X_i]$ \uparrow
dt L~1-0 Continuum	Proof: Suppose EX; J <o at="" for="" i="" some="" t<="" th=""></o>
approximation bic molecules are	(#IX.1<0 - X;<0 +
reacting, but overall the net vate is balanced	CONTRADICTION Twinen = 0
Lots of copies of x	C

$$\frac{d}{dt}\vec{X} = A^{(1)}\vec{X} + A^{(2)} \begin{bmatrix} [x,][x,]] \\ [x,][x_{n}] \\ \vdots \\ [x,]] \\ \vdots \\ [x,]] \end{bmatrix} \xrightarrow{terms} terms \\ \vec{X}^{(1)}[x_{n}] \\ \vec{X}^{(2)} \\ \vec{X$$







2.15 Steady-State Problems











2.17 Parameter Estimation; Robustness, Fragility,

Control



struction



2.19 Deconvolution



H=UZVT



-> Remove terms with small o

2.20 Deconvolution II





2.21 Blind Deconvolution



2.22 Optical Flow





2.23 High-Throughput Data and Analysis


70 of _ reaction - 1 M particles 3 - 100 particles 100 particles ODE time

Chapter 3

Introduction to Numerical Simulation

Introduction to Numerical Simulation teaches an immense amount of material, it's a general overview for graduate level numerical analysis and covers everything from how to formulate equations based on a model to procedures one should use when solving linear equations. This chapter offers in depth analysis as to why these procedures work and allows the reader to look at problems from several different perspectives. The material is presented in an easy to follow manner with several examples to elucidate each point.

INTRODUCTION TO NUMERICAL SIMULATION

LECTURE 1. Example Problems and Basic Equations

COURSE OUTLINE:

RELATION TO OTHER COURSES

6.337J / 18.335J

Numerical Linear Algebra

Ax=b

PDE

2.097J / 6.339J / 16.920J

PDE solvers

•	Assembling system of equations automatically	[3 lec]	
•	Steady state solution		
	 Linear Problems: Ax=b 	[7 lec]	
	 Non-Linear Problems: F(x)=0 	[3 lec]	
•	Dynamics (ODE solvers)		
	 Time domain integration 	[3 lec]	
	 Periodic steady state 	[2 lec]	
•	PDE solvers	. ,	
	 Integral Equation Methods 	[2 lec]	
	 Finite Element Methods (FEM) 	[1 lec]	
	 Finite Difference Methods (FD) 	[2 lec]	
	 Preconditioners for PDE solvers 	[1 lec]	
•	Model Order Reduction	[2 lec]	

2.096J / 6.336J / 16.910J

Numerical Simulation

Breadth course

6.255J / 2.098J / 15.093J Optimization

Breadth course

6.581 / 20.482

Computational Biology

6.242 Model Order Reduction

(biannual)

F(x)=0

ODE



- Recent Developments
- Fast Solvers for Integral Formulations
- · Structural Analysis of Automobiles Equations
 - Force-displacement relationships for mechanical elements (plates, beams, shells) and sum of forces = 0.

Automatic Error Control

- Partial Differential Equations of Continuum Mechanics Recent Developments
- Meshless Methods, Iterative Methods,



now Coventor

2

ಲು

Drag Force Analysis of Aircraft

- Equations Navier-Stokes Partial
- Differential Equations Recent Developments
- Multigrid Methods for Unstructured Grids
- · Analysis of Cell Traps for Sorting
- Cytometry Equations
 - Navier-Stokes Partial
 - Differential Equations
- Recent Developments Multigrid Methods for
- Unstructured Grids · Engine Thermal Analysis
 - Equations
 - Poisson Partial Differential Equations
 - Recent Developments
 - Fast Integral Equation Solvers, Monte-Carlo Methods
- Micromachined Device Performance Analysis Equations
 - Elastomechanics, Electrostatics, Stokes Flow
 - Recent Developments
 - Fast Integral Equation Solvers, Matrix-Implicit Multi-level Newton Methods for coupled domain
- problems · Stock Option Pricing for Hedge Funds
- Equations
 - Black-Scholes Partial Differential Equation
- Recent Developments
- Financial Service Companies are hiring engineers,
- mathematicians, and
 - physicists



- · Virtual Environments for Computer Games
 - Equations
 Multibody Dynamics, Elastic
 - Collision Equations
 - o Recent Developments Multirate
 - integration
 - methods, parallel simulation
- Virtual Surgery
 - Equations
 - Partial Differential Equations of
 - Elastomechanics o Recent Developments · Parallel Computing, Fast



- Methods Biomolecule Electrostatic Optimization
- Equations The Poisson Partial Differential Equation
- o Recent Developments
 - · Matrix-Implicit Iterative Methods, Fast Integral Equation Solvers





Ligand Receptor (drug molecule) (protein molecule)

Ecm protein

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3

Stock Price

THE COMPUTER SIMULATION SCENARIO



COURSE PHILOSOPHY

Examine Several Modern Techniques Understand, practically and theoretically, how the techniques perform on

- representative, but real, applications Why Prove Theorems?
- Guarantees, given assumptions, that the method will always work.
- Can help debug programs.

The theorem proof can tell you what to do in practice.

EXAMPLE PROBLEMS

Power Distribution on an Integrated Circuit



Is there at least 3V across the ALU? 0 Design Objectives for the VLSI Problem Select topology and metal widths and lengths so that

a) Voltage across every function block > 3V b) Minimize the area used for the metal wires First Step – Analysis Tools Given the topology and metal widths and lengths determine the voltage across the ALU, Cache and Decoder. o Who uses VLSI Tools? Several big companies IBM, Motorola, TI, Intel, Compaq, Sony, Hitachi Nonfunctional prototype costs: · Increases time to market · Design rework costs millions 1000's of small companies · Small companies make application circuits disk drives, graphics accelerators, CD players, cell phones · What is the cost of nonfunctional prototypes? Out of business o Who makes VLSI Tools? Company Market capital Employees Sales Cadence 4,000 1.3 billion 3.8 billion 1.5 billion Synopsis/Avanti 5,000 6.9 billion Mentor Graphics 2,600 0.6 billion 1.4 billion (Data from 2003) Companies compete by improving analysis efficiency. • Modeling VLSI Circuit Power Distribution

- Power Supply provide current at a certain voltage
- Functional blocks draw current
- The wire resistance generates losses
- Modeling the Circuit
- · Supply becomes a Voltage Source







Element current is related to voltage across the element

- Two Types of Equations
- Sum of currents at each node = 0 • Conservation/Balance Law Equation

$$\begin{array}{c}
R_{1} \quad i_{1} \quad i_{3} \quad R_{3} \quad (i_{1} - i_{2} + i_{3} = 0) \\
\hline
R_{2} \\
\hline
i_{2} \quad OR \\
(-i_{1} + i_{2} - i_{3} = 0)
\end{array}$$

Note: here $i_1 = -3$ means it goes to the right i.e. opposite to the arrow





7

Modeling the Space Frame













 $f_y^* = \frac{y^* - y}{L} \epsilon (L_0 - L) = \frac{6 - 3}{5} 10 (3 - 5) = -12$







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incoming heat

per unit length

 T_{N-1} T_N

Temperature analogous to Voltage Heat Flow analogous to Current



v = T(0) $i = h \Delta x$ T(1)

SUMMARY OF KEY POINTS:

- Many Applications of simulation
 - Picked Three Representative Examples Circuits, Struts and Joints, Heat Flow in Bar
- Two Types of Unknowns
 Circuit Node Voltages, Element Currents



 $\vec{r} = (x, y)$

Conducting Bar – Temperature, Section Heat Flows

$$T_i \bigcirc T_{i-1}$$

Two Types of Equations

- Conservation/Balance Laws
 - Circuit Sum of Currents at each node = 0

$$\overset{V_{i-1}}{\underset{i_{s}}{\overset{W_{i}}{\bigoplus}}} \overset{R_{\theta}}{\underset{i_{n}}{\overset{V_{i-1}}{\bigoplus}}} \overset{R_{\theta}}{\underset{i_{n}}{\overset{V_{i-1}}{\bigoplus}}} \quad i_{A} - i_{B} + i_{s} = 0$$

Struts – Sum of Forces at each joint = 0

• A **B ***
$$\vec{f}_A^* - \vec{f}_B^* + \vec{f}_L = 0$$

• Bar – Sum of heat flows into control volume = 0

Incoming Heat
$$h_s$$

 T_{i-1} $h_{i,j-1}$ T_i $h_{i-1,i}$ T_{i+1} $h_{i-1,j}$ h_{i-

 $h_{i,i-1} - h_{i+1,i} - \tilde{h}_s \Delta x = 0$

Constitutive Equation

 (x_1, y_1)

-0-A

Circuit - current-voltage relationship

• Struts - force-displacement relationship

(x2,

$$f_{A,y}^* = \frac{x_2 - x_1}{L_A} \varepsilon (L_{A,0} - L_A)$$

$$f_{A,y}^* = \frac{y_2 - y_1}{L_A} \varepsilon (L_{A,0} - L_A)$$



INTRODUCTION TO NUMERICAL SIMULATION

LECTURE 2. Equation Formulation & Node-Branch Stamping

TODAY'S OUTLINE:

•

- · Formulating Equations
 - Circuit Example
 - Struts and Joints Example
 - Matrix Construction From Schematics
 - Node-Branch "Stamping Procedure"
 - Circuits
 - Struts and Joints

FORMULATING EQUATIONS FROM SCHEMATICS Circuit Example



• Step 1: Identifying Unknowns

(0)









 Step 3: Constitutive Equations
 Use Constitutive Equations to relate branch currents to node voltages (Currents flow from plus node to minus node)



Struts Example



3.2Т Node Branch Stamp-

ing



Will the solutions be the same? Will the set of conservation law equations be different?

Conservation laws for the two examples will be exactly the same. The perceived force "direction" is inconsequential, it is the adjacent forces that matter.





4



Use Constitutive Equations to relate strut forces to joint positions.



Struts: force-displacement relationship



 f_x^* is a positive number – force is to the right, this makes sense for a stretched strut.

· Bar: temperature drop-heat flow relationship

$$T_{i} \bigoplus_{T_{i-1}} T_{i-1} \qquad h_{i+1,i} = \kappa \frac{T_{i+1} - T_{i}}{\Delta x}$$

GENERATING MATRICES FROM SCHEMATICS

hi+1,i



Assume linear constitutive equations... B N branch node currents voltages N KCL Eqns l_b V_{i} B Constitutive Eqns • One matrix column for each unknown - N columns for the Node voltage - B columns for the Branch currents · One matrix row for each equation - N rows for KCL - B rows for element constitutive equations (linear and square system!) o Conservation Equation $=-i_{sA}$ $\bullet - i_A + i_B$ 0 $-i_E = -i_{sB} - i_{sC}$ $-i_B$ 0 $=i_{sC}$ ic 0 1) $=i_{sA}+i_{sB}$ $-i_C + i_D$ Matrix Form for the Equations The matrix A is usually not square 0 $-i_{sA}$ one row 0 $i_{sB} - i_{sC}$ for each ic KCL 6 isc equation 4 -1 $i_{sA} + i_{sB}$

B

Circuit Example



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How does each resistor contribute to the matrix?



How does each current source contribute to the Conservation Law Equation?



Conservation Matrix Equation Generation Algorithm

For each resistor if $(n_1 \neq 0)$ then $A(n_1,k) \leftarrow 1$ if $(n_2 \neq 0)$ then $A(n_2,k) \leftarrow -1$





Constitutive Equation



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First relate branch currents to branch voltages Second determine voltages across resistors (Branch Voltages)





 $\begin{vmatrix} i_B \\ i_C \end{vmatrix} - \alpha A^T \begin{vmatrix} V_1 \\ V_2 \\ V_3 \end{vmatrix}$

i_D

i_E

The node voltages can be related to branch currents

- A^T relates node voltages to branch voltages.

- α relates branch voltages to branch currents.

a is square and diagonal

• Node-Branch Form $I_b - \alpha A^T V_N = 0$ Constitutive Relation

 $A I_b = I_s$ Conservation Law

$$\begin{array}{ccc} B \updownarrow \begin{bmatrix} I & -\alpha A^T \\ A & 0 \\ B & \overleftarrow{N} \end{array} \end{bmatrix} \begin{bmatrix} I_h \\ V_N \end{bmatrix} = \begin{bmatrix} 0 \\ I_s \end{bmatrix}$$

N = number of <u>N</u>odes with unknown voltages B = number of <u>B</u>ranches with unknown currents

Struts Example

o In 2-D

- One pair of columns for each unknown
 - J pairs of columns for the joint positions
 - S pairs of columns for the strut forces
- One pair of matrix rows for each equation
- J pairs of rows for the force equilibrium equations
 S pairs of rows for the *linearized* constitutive relations
- Follow Approach Parallel to Circuits
 - (1) Form an "Incidence Matrix," A, from Conservation Law.
 - (2) Determine strut deformation using A^{T} .
 - (3) Use linearized constitutive equations to relate strut deformation.
- (4) Combine (1), (2), and (3) to generate a node-branch form.
- Conservation Laws



Note that struts A & B only contribute one pair of entries into the A matrix and strut C contributes two pairs of entries into the A matrix. This is because struts A & B are connected to the wall (ground) and strut C has two free ends.

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 $V_{4} = 0$







Linearization is only a valid estimate if the strut is very close to its original configuration -x and y are very close to x_0 and y_0 .

What if the strut is *rotated*, is this linearization still valid? It seems like it should be because the strut is not stretched (only rotated), but it isn't. Because the non-linearity is *not* in the relationship between how much the strut is stretched and the force going through the strut but the non-linearity is in the projection of that force onto the *x* and *y* axis. Referred to as the "geometric" nonlinearity in finite element literature. Example.



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We will learn more about linearizing equations later on when we study Newton's Method. The a(s,s) block



Note that the arguments in the α matrix are the difference in the original position between node 1 and node 2. This is because the force of the strut is dependent upon the relative position of the two joints on either side.



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• Node-Branch Form $f_S - \alpha A^T u = 0$ $A f_S = f_L$

Constitutive Equation Conservation Law

$$\begin{array}{c} 2 \cdot S \stackrel{\uparrow}{\downarrow} \begin{bmatrix} I & -\alpha A^T \\ A & 0 \\ 2 \cdot J \stackrel{\uparrow}{\downarrow} \begin{bmatrix} A & 0 \\ 2 \cdot S & \overleftarrow{2} \cdot J \end{bmatrix} \begin{bmatrix} f_s \\ u \end{bmatrix} = \begin{bmatrix} 0 \\ f_L \end{bmatrix}$$

S = Number of Struts J = Number of Unfixed Joints

Comparison

struts
$$2 \cdot S \ddagger \begin{bmatrix} I & -\alpha A^T \\ A & 0 \end{bmatrix} \begin{bmatrix} f_s \\ u \end{bmatrix} = \begin{bmatrix} 0 \\ f_L \end{bmatrix}$$
$$2 \cdot J \ddagger \begin{bmatrix} 0 \\ 2 \cdot S & \overleftarrow{2} \cdot J \end{bmatrix}$$
circuit
$$B \ddagger \begin{bmatrix} I & -\alpha A^T \\ N \ddagger \begin{bmatrix} I & -\alpha A^T \\ A & 0 \\ \overrightarrow{B} & \overrightarrow{N} \end{bmatrix} \begin{bmatrix} I_b \\ V_N \end{bmatrix} = \begin{bmatrix} 0 \\ I_s \end{bmatrix}$$

 Summary of Key points...

 O
 Developed algorithms for automatically constructing matrix equations from schematics using

 E
 Conservation law

 Constitutive equations

 Looked at one formulation: node-branch
 Next time: nodal formulation

INTRODUCTION TO NUMERICAL SIMULATION

LECTURE 3. Equation Formulation - Nodal Analysis

TODAY'S OUTLINE:

- Matrix Construction from Schematics
 - Nodal "Stamping Procedure"
 - Struts and Joints
 - Circuits
 - Comparing Node-Branch vs. Nodal
- Solution of Linear Systems

 Existence and Uniqueness

MATRIX CONSTRUCTION FROM SCHEMATICS Nodal "Stamping Procedure"

Struts and Joints

Two struts aligned with the X axis









(1) Number the nodes with one node as 0.

Two struts aligned with the X axis

(2) Write a conservation law at each node except (0) in terms of the node voltages!



Notice that the contributions are positive on the diagonal and negative on the off-diagonal. *G* is square.





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3















- Node-branch
- Node-branch
- General constitutive equations
- Large sparser system
- No diagonal dominance
- Nodal
- · Conserved quantity must be a function of node variables
- Smaller denser system
- · Diagonally dominant & symmetric

SOLUTION OF LINEAR SYSTEMS

$$\begin{bmatrix} \uparrow & \uparrow & \dots & \uparrow \\ \vec{M}_1 & \vec{M}_2 & \dots & \vec{M}_N \\ \downarrow & \downarrow & \dots & \downarrow \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ x_N \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ \vdots \\ b_N \end{bmatrix} \implies x_1 \vec{M}_1 + x_2 \vec{M}_2 + \dots + x_N \vec{M}_N = \vec{b}$$

Find a set of weights, x, so that the weighted sum of the columns of the matrix M is equal to the right hand side b.

11







★ Is the solution unique? Suppose there exist weights, y₁, ..., y_N, not all zero y₁M₁ + y₂M₂ + ... + y_NM_N = 0 → [nullspace] Then if **M**x = b, then **M**(x + y) = b A solution is unique only if the columns of **M** are linearly independent. Linearly independent columns means: $\sum_{i \neq k} \alpha_i \vec{M}_i \neq \vec{M}_k$

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$\mathbf{M}\vec{x} = \vec{b}$ $\mathbf{M}\vec{y} = 0$	\vec{x} is a solution linearly dependent columns in M	
$\mathbf{M}(\vec{x} + \vec{y}) = \vec{b}$	$\vec{x} + \vec{y}$ is a solution	

If $\mathbf{M}\vec{y} = \mathbf{0}$ then $\alpha \mathbf{M}\vec{y} = \mathbf{0} \Rightarrow \mathbf{M}(\vec{x} + \alpha \vec{y}) = \vec{b} \Rightarrow$ infinite solutions







Struts and Joints Example.)

$$\begin{aligned} x_{1} & f_{x}^{*} x_{2} \\ x_{01} & y_{2} \\ x_{1} & x_{1} \\ x_{01} & u_{x_{2}} \\ x_{2} \\ x_{1} \\ x_{2} \\ x_{2} \\ x_{2} \\ x_{2} \\ x_{2} \\ x_{1} \\ x_{1} \\ x_{2} \\ x_{1} \\ x_{1} \\ x_{2} \\ x_{1} \\ x_{2} \\ x_{1} \\ x_{1} \\ x_{2} \\ x_{1} \\ x_{1} \\ x_{2} \\ x_{1} \\ x_{1} \\ x_{1} \\ x_{2} \\ x_{1} \\ x_{1} \\ x_{1} \\ x_{2} \\ x_{1} \\ x_$$

Singular System.



Heat Conducting Bar Example.)



Boundary conditions

if heat in = heat out \rightarrow no unique solution

Could displace solution by a factor of 100, and it would still solve the system









Square Matrices

Given Mx = b, where M is square If a solution exists for all b, then the solution for a specific b is unique. For a solution to exist for any b, the columns of M must span all N-length vectors. Since there are only N columns of the matrix M to span this space, these vectors must be linearly independent.

A square matrix with linearly independent columns is said to be nonsingular.

INTRODUCTION TO NUMERICAL SIMULATION

LECTURE 4. Linear Systems - LU Decomposition

TODAY'S OUTLINE:

· Solution of Linear Systems

- Gaussian Elimination Basics
 - LU factorization
 - · Computational Complexity
 - · Pivoting for Growth Control

SOLUTION OF LINEAR SYSTEMS

Gaussian Elimination Basics

LU Factorization

- Important Properties
 - Gaussian Elimination Method for solving $\mathbf{M}\vec{x} = \vec{b}$
 - A "Direct" Method
 - Finite Termination for exact result (ignoring roundoff error)
 - Produces accurate results for a broad range of matrices
 - Computationally expensive
 - o Reminder by Example

3×3 Example

$$\begin{bmatrix} M_{11} & M_{12} & M_{13} \end{bmatrix} \begin{bmatrix} x_1 \\ y_1 \end{bmatrix} \begin{bmatrix} b_1 \\ y_2 \end{bmatrix} = M_{11}x_1 + M_{12}x_2 + M_{13}x_3 = b_1$$

 M_{21} M_{22} M_{23} $x_2 = b_2$ $M_{21}x_1 + M_{22}x_2 + M_{23}x_3 = b_2$

 M_{31} M_{32} M_{33} x_3 b_3 $M_{31}x_1 + M_{32}x_2 + M_{33}x_3 = b_3$

Use equation 1 to eliminate x_1 from equation 2 and 3 $M_{11}x_1 + M_{12}x_2 + M_{13}x_3 = b_1$

$$M_{21} - \frac{M_{21}}{M_{11}} M_{11} \Big] x_1 + \Big(M_{22} - \frac{M_{21}}{M_{11}} M_{12} \Big) x_2 + \Big(M_{23} - \frac{M_{21}}{M_{11}} M_{13} \Big) x_3 = b_2 - \frac{M_{21}}{M_{11}} b_1 \\ M_{31} - \frac{M_{31}}{M_{11}} M_{11} \Big] x_1 + \Big(M_{32} - \frac{M_{31}}{M_{11}} M_{12} \Big) x_2 + \Big(M_{33} - \frac{M_{31}}{M_{11}} M_{13} \Big) x_3 = b_3 - \frac{M_{31}}{M_{11}} b_1 \\ \end{pmatrix}$$

м









b

х =

M

Fitting the pieces together



Store all the data in one matrix - no extra memory needed.



3



The L matrix will contain all subdiagonal elements with a diagonal of all 1's

• Factoring
An "in place" implementation
$$\begin{bmatrix} M_{11} & M_{12} & M_{13} \\ M_{21} & M_{22} & M_{23} \\ M_{31} & M_{32} & M_{33} \end{bmatrix} \Longrightarrow \begin{bmatrix} M_{11} & M_{12} & M_{13} \\ \frac{M_{21}}{M_{11}} & \widetilde{M}_{22} & \widetilde{M}_{23} \\ \frac{M_{21}}{M_{11}} & \widetilde{M}_{22} & \widetilde{M}_{23} \\ \frac{M_{31}}{M_{11}} & \frac{\widetilde{M}_{32}}{\widetilde{M}_{22}} & \widetilde{M}_{33} \end{bmatrix}$$

 Example – Heat Flow Temperature analogous to Voltage Heat Flow analogous to Current





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5










 Wo important i neorems
 Partial pivoting (swapping rows) <u>always</u> succeeds if M is nonsingular.





(N-1) Linearly Independent Columns

 LU factorization applied to a strictly diagonally dominant matrix will <u>never</u> produce a zero pivot.



 $-1 \times 10^{-7} + \pi \cdot 10^{-7} = 2.141592653589793 \cdot 10^{-7}$ All 15 decimal digits are correct in this case Ex #2.) $1 + \pi \cdot 10^{-7} = 1.000000314159265$ lost 7 digits of π -1.0000001+1.000000314159265 $= 2.141592652105118 \cdot 10^{-7}$ junk Only the first 8 decimal digits are correct in this case, the remainder are garbage. Lose digits of precision Key Issue Avoid additions and subtractions between large and small numbers! EVEN BETTER: AVOID GENERATING LARGE NUMBERS AT ALL IF POSSIBLE !!! Back to the contrived Example Back to the contrived Example $\begin{bmatrix} 10^{-17} & 1 \\ 1 & 2 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 1 \\ 3 \end{bmatrix}$ $LU_{Exact} = \begin{bmatrix} 1 & 0 \\ 10^{17} & 1 \end{bmatrix} \begin{bmatrix} 10^{-17} & 1 \\ 0 & 2 - 10^{17} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 1 \\ 3 \end{bmatrix}$ $LU_{Rounded} = \begin{bmatrix} 1 & 0 \\ 10^{17} & 1 \end{bmatrix} \begin{bmatrix} 10^{-17} & 1 \\ 0 & -10^{17} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 1 \\ 3 \end{bmatrix}$ $\begin{bmatrix} x_1 \\ x_2 \end{bmatrix}_{Exact} = \begin{bmatrix} 1 \\ 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix}_{Rounded} = \begin{bmatrix} 0 \\ 1 \end{bmatrix} \bigstar$ Original Problem: $\begin{bmatrix} 10^{-17} & 1 \\ 1 & 2 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 1 \\ 3 \end{bmatrix}$ Swap the rows $\begin{bmatrix} 1 & 2 \\ 10^{-17} & 1 \end{bmatrix}$ Partial Pivoting for Roundoff Reduction If $|M_{ji}| < \max_{j>i} |M_{ji}|$ swap row *i* with $\arg\left(\max_{j>i} |M_{ij}|\right)$

$$LU_{reordered} = \begin{bmatrix} 1 & 0 \\ 10^{-17} & 1 \end{bmatrix} \begin{bmatrix} 1 & 2 \\ 0 & 1-2 \cdot 10^{-17} \end{bmatrix}$$

This multiplier
is small
This term still gets rounded,
but this time the multiplier
does not 'overpower' the row.

$$L = \begin{bmatrix} 1 & 0 \\ 10^{-17} & 1 \end{bmatrix} U = \begin{bmatrix} 1 & 2 \\ 0 & 1 \end{bmatrix}$$

Let's solve $LUx = b$.
First, solve $Ly = b$:

$$Ly = b \rightarrow \begin{bmatrix} 1 & 0 \\ 10^{-17} & 1 \end{bmatrix} \begin{bmatrix} y_1 \\ y_2 \end{bmatrix} = \begin{bmatrix} 1 \\ 3 \end{bmatrix}$$

Next, solve $Ux = y$:

$$Ux = y \rightarrow \begin{bmatrix} 1 & 2 \\ 0 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 1 \\ 3 \end{bmatrix} \Rightarrow x = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} -5 \\ 3 \end{bmatrix} \bigstar$$
 STILL
WRONG
PROBLEM: Need to swap the elements of b as well!
Solve $Ly = b$:

$$Ly = b \rightarrow \begin{bmatrix} 1 & 0 \\ 10^{-17} & 1 \end{bmatrix} \begin{bmatrix} y_1 \\ y_2 \end{bmatrix} = \begin{bmatrix} 3 \\ 1 \end{bmatrix}$$

$$\Rightarrow y = \begin{bmatrix} y_1 \\ y_2 \end{bmatrix} = \begin{bmatrix} 3 \\ 1 \end{bmatrix}$$

Next, solve $Ux = y$:

$$Ux = y \rightarrow \begin{bmatrix} 1 & 0 \\ 10^{-17} & 1 \end{bmatrix} \begin{bmatrix} y_1 \\ y_2 \end{bmatrix} = \begin{bmatrix} 3 \\ 1 \end{bmatrix}$$

Next, solve $Ux = y$:

$$Ux = y \rightarrow \begin{bmatrix} 1 & 0 \\ 10^{-17} & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 3 \\ 1 \end{bmatrix}$$

Next, solve $Ux = y$:

$$Ux = y \rightarrow \begin{bmatrix} 1 & 0 \\ 1 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 3 \\ 1 \end{bmatrix} \Rightarrow x = \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 1 \\ 1 \end{bmatrix}$$
 Correct Answer.

-

e.	Exan	nple.										
	[1	0	0	•••	0	17-	7 1					
	10	1	0		0	0	×1	01				
	0	10	1		0	0 2	^x 2	02				
	:	3	٠.,	٠.,			v ₃ =	b3				
	0	0		10	1	0	:	:				
	0	0		0	10	1 -1	N	b_n				
	L		[1	0	0		1]	$\begin{bmatrix} x_1 \end{bmatrix}$	Ιſ	b_1		
			10	1	0		-10	X2		b_2		
		⇒		10	1		0	X3	=	b_3		
					٠.	·.,	:	:		:		
						10	1	XN		b _n		
			Ē1	0	0		1	T_{x_1}	1	$\begin{bmatrix} b_1 \end{bmatrix}$		
			10	1	0		-10	x2		b2		
		\Rightarrow	0	10	1		-100) X3	=	ba		
					٠.	٠.	:	:		:		
						10	1	XA		b_n		
			-	[]	0	0		1	-	Tx	1	[]
				10	0 1	0]	0	x2		1
		\Rightarrow	=	⇒ 0	10) [- 1	00	.X3	=	
						۰.	÷.,	-		:		
							10	10	N-1			

 $\begin{bmatrix} 10 & -10^{N-1} \end{bmatrix} x_N \end{bmatrix} \begin{bmatrix} b_n \end{bmatrix}$ The last entry is very large! Use partial pivoting – swap first and second rows...

- If the matrix is strictly diagonally dominant
- if the matrix is strictly diagonally dominant
 or if use partial pivoting for round-off reduction:
 The multipliers will <u>always</u> be smaller than one in magnitude.
 The maximum magnitude entry in the *LU* factors will never be larger than 2⁽ⁿ⁻¹⁾ times the maximum magnitude entry in the original matrix.

For 1000 nodes, what is 2^{1000} , Know that $2^{10} \approx 10^3$. So, $2^{1000} = (2^{10})^{100} = 10^{100} \rightarrow$ Very large number! Might not be a very useful theorem \rightarrow generally not this large.



INTRODUCTION TO NUMERICAL SIMULATION

LECTURE 5.

Linear Systems - Conditioning

TODAY'S OUTLINE:

- · Solution of Dense Linear Systems
 - Hard to Solve Problems
- Perturbation Analysis and Conditioning
- · Solution of Sparse Linear Systems
 - LU Factorization Reminder.
 - Example of Problems with Sparse Matrices
 - · Struts and joints, resistor grids, 3-D heat flow
 - Tridiagonal Matrix Factorization
 - General Sparse Factorization
 Fill-in and Reordering
 - Graph Based Approach
 - Sparse Matrix Data Structures
 - Scattering



 $1 t_N t_N^2$



 $t_N^N \alpha_N$

 $f(t_N)$

IN











orthogonal columns

 $M = \begin{bmatrix} \vec{M}_1, \vec{M}_2 \end{bmatrix}$ Solving, $\mathbf{M}x = b$ is finding $x_1\vec{M}_1 + x_2\vec{M}_2 = \vec{b}$



Columns nearly aligned

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When vectors are nearly aligned, difficult to determine how much of \vec{M}_1 versus how much of $~\vec{M}_2$ $M = \begin{bmatrix} 1 & 0 \\ 0 & 10^{-6} \end{bmatrix}$

 $M = \begin{bmatrix} 1 & 1 - 10^{-6} \\ 1 - 10^{-6} & 1 \\ \text{columns nearly aligned} \end{bmatrix}$









Columns are very close to being linearly dependent (closely aligned) – have a larger condition number / columns are not close to orthogonal









Perturbation Analysis and Conditioning Induced Norms

o Matrix Magnification Question Suppose $y = \mathbf{M}x$ How much larger is y than x? OR How much does **M** magnify x?









Matrix Norms [Standard Induced *l*-norms]
 Definition:

$$||M||_{l} = \max_{x} \frac{||Mx||_{l}}{||x||_{l}} = \max_{x|_{l}=1} ||Mx||_{l}$$





|A| is a measure of the largest gain or magnification power





★ Hard to Solve Systems
○ Perturbation Analysis
$$\begin{bmatrix} \left\| \delta x \right\| & 2 \\ \|x + \delta x \| \end{bmatrix} & \left\| M \right\| & \left[M + \delta M \right] & \left[x + \delta x \right] &= b \\ M & x &= b \\ Perturbation Equation$$

$$\frac{\left(M + \delta M \right) & \left(x + \delta x \right) \\ models LU models solution \\ roundoff perturbation$$
Since Mx + b = 0
$$M\delta x = -\delta M (x + \delta x) \Rightarrow \delta x = -M^{-1}\delta M (x + \delta x)$$
Taking Norms
$$\| \delta x \| \leq \|M^{-1}\| \|M\| \frac{1}{\|M\|} \|\delta M\| \|x + \delta x\|$$
Relative Error Relation
$$\frac{\| \delta x \|}{\|x + \delta x\|} \leq \left\| M^{-1} \| \|M\| \frac{\| \delta M \|}{\|M\|}$$
Relative The Relation
$$\frac{\| \delta x \|}{\|x + \delta x\|} \leq \left\| M^{-1} \| M \| \frac{\| \delta M \|}{\|M\|}$$

$$\begin{split} \mathbf{M}\vec{x} &= \vec{b} \\ \text{Assume } \frac{\|\delta M\|}{\|M\|} = 10^{-16} \quad \text{What is} \quad \left\|\frac{\delta x}{\|x\|}\right| = \frac{\|\delta x\|}{\|x + \delta x\|} = ? \\ \text{e.g.} \quad \left\|M^{-1}\right\| \|M\| = 100 \qquad \left\|\frac{\|\delta x\|}{\|x + \delta x\|} \le 100 \cdot 10^{-16} = 10^{-14} \\ \text{e.g.} \quad \left\|M^{-1}\right\| \|M\| = 10^{13} \\ \frac{\|\delta x\|}{\|x + \delta x\|} \le 10^{13} \cdot 10^{-16} = 10^{-3} = 0.1\% \text{ error} \\ \text{e.g.} \quad \left\|M^{-1}\right\| \|M\| = 10^{20} \\ \frac{\|\delta x\|}{\|x + \delta x\|} \le 10^{20} \cdot 10^{-16} = 10^{4} = 10^{6}\% \text{ error!!} \\ \text{Large condition number} \Rightarrow \text{Big error} \end{split}$$

o Geometric Analysis / Polynomial Interpolation

















node)









Where can fill-in occur?









INTRODUCTION TO NUMERICAL SIMULATION

LECTURE 6.A.

Solution of Sparse Linear Systems

TODAY'S OUTLINE:

- Solution of Sparse Linear Systems
 - General Sparse Factorization
 - · Graph Based Approach
 - Sparse Matrix Data Structures
 - Scattering

SOLUTION OF SPARSE LINEAR SYSTEMS General Sparse Factorization

- Graph Based Approach
- Oraph Dased Appro
 Construction

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- 0 Construction
 - Structurally Symmetric Matrices and Graphs Note that *structurally* symmetric does not imply that the values of the matrix are symmetric (i.e., $a_{12} \neq a_{21}$)



One Node Per Matrix Row
One Edge Per Off-Diagonal Pair

Can one apply these graph-based techniques to the following matrix?

1	3	0	0	
0	1	2	0	
0	-		-	

- 0 2 1 5
- 0 0 5 1

This matrix is not "structurally symmetric" in that there is a zero in a_{21} and a non-zero value in a_{12} .

- Can still use this approach, just treat it as if there is a nonzero in the a_{21} place and use the graphs to do the analysis; there will be some efficiency loss, but the methods will still work.
- Thus, this technique may be applied to "mildly" structurally symmetric matrices by assuming there is structural symmetry with some loss of efficiency in treating some of the zeros as if they are non-zeros.



When node 1 is removed, in the matrix, the non-zero entries in the first row create fill-ins that connect up the other nodes that can be seen in the graph.















Factoring an $m \times m$ grid.

Dense LU: cost is $O(m^2)^3 = O(m^6)$

Sparse LU: each row has approximately 5 nonzeros matrix has approximately $5m^2$ nonzeros cost is at least $O(m^3)$

Separator is of size 2m

Factoring: cost is $O(2m)^3 = O(8m^3)$ {dense matrix} Subgraph: size is m/2Create 4 separators: size is $2 \times m/2 = m$

Cost is $O(m^3)$ per separator $\rightarrow O(4m^3)$

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As one continues to subdivide, the cost keeps halving. Adding together these costs, the final cost will be roughly $O(m^3)$



Structure & Dimension



 $\dot{\gamma}_{1}$



- or to Avoid Fill-ins?
- A.) LU factorization applied to a strictly diagonally dominant matrix will never produce a zero pivot.
- B.) The matrix entries produced by LU factorization applied to a strictly diagonally dominant matrix will never increase by more than a factor of $2^{(n-1)}$. [which is the best you can do by pivoting for growth control] Bottom Line:
- If your matrix is strictly diagonally dominant no need for numerical pivot for growth control so just pivot for sparsity control! 9 Sparse Factorization Approach 1. Assume matrix requires NO numerical pivoting.
- - Diagonally dominant or symmetric positive definite.
 Use graphs to determine matrix ordering.
 - Many graph manipulation tricks used.
 - 3. Form data structures for storing filled-in matrix. Lots of additional nonzeros added.
 - 4. Put numerical values in data structure and factor
 - Computation must be organized carefully!

Sparse Matrix Data Structures











Why store sparse matrix information in a data structure array?

- · Too much storage space storing all the zero entries
- Avoid floating point computation on all the zeros (minimal in comparison to the memory cost)

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Memory reference cost

, i i i i i i i i i i i i i i i i i i i	ng Sour	ce Row	i from	Farget F	low j:		
Pow i	$M_{i,i+1}$	$M_{i,i+7}$	$M_{i,i+15}$				
Rowr	<i>i</i> +1	<i>i</i> + 7	<i>i</i> +15				
Pow i	M _{i,i+1}	$M_{i,i+4}$	$M_{i,i+5}$	$M_{i,i+7}$	$M_{i,i+9}$	$M_{i,i+12}$	$M_{i,i+15}$
Row J	i+1	<i>i</i> + 4	i+5	<i>i</i> +7	i+9	<i>i</i> +12	i+15
	Rows	Ops		Misses			
Res	Rows 300	Ops 904,3	887	Misses 248,96	7		
Res RAM	Rows 300 2,806	Ops 904,3 1,017	887 289	Misses 248,96′ 3,817,5	7 87	More miss	cs
Res RAM Grid	Rows 300 2,806 4,356	Ops 904,3 1,017 3,180	887 ,289),726	Misses 248,96 3,817,5 3,597,7	7 87 46	More misse than operat	cs tions!

Scattering



Read all the elements in Row j, and scatter them in an n-length vector
 Access only the needed elements using array indexing!



- Useful for estimating Sparse GE complexity
- Sparse Data Structures
 - Scattering

INTRODUCTION TO NUMERICAL SIMULATION

LECTURE 6.B.

QR Factorization

TODAY'S OUTLINE:

Singular Problems

- Projection Formulas
- Modified Gram-Schmidt Algorithm

LU decomposition fails

· But a solution exists! Actually, many





LU will fail



2

 $\dot{\omega}$

-1

Linear Systems

I.

QR Factorization



Recall weighted sum of columns view of systems of equations

actually lots of them. How do we find them?

Systems of Linear Equations – Summary Table

	$b \in range\{M\}$	b∉ range{M}
L.I. columns	Solution exists and is unique Use LU	No solutions Find the "closest"
L.D. columns	Infinite solutions exist Find one or all	No solutions Find the "closest"



Definition of orthonormal matrix Q: $\vec{Q}_i \cdot \vec{Q}_j = 0$ if $i \neq j$ and $\vec{Q}_i \cdot \vec{Q}_i = 1$ Picture for the two-dimensional case



Q is orthonormal

$$\mathbf{Q}^{T}\mathbf{Q} = \begin{bmatrix} \overline{Q}_{1}^{T}\overline{Q}_{1} & \overline{Q}_{1}^{T}\overline{Q}_{2} & \overline{Q}_{1}^{T}\overline{Q}_{3} & \cdots \\ \overline{Q}_{2}^{T}\overline{Q}_{2} & \overline{Q}_{2}^{T}\overline{Q}_{2} & \overline{Q}_{2}^{T}\overline{Q}_{3} & \cdots \\ \overline{Q}_{1}^{T}\overline{Q}_{1} & \overline{Q}_{1}^{T}\overline{Q}_{2} & \overline{Q}_{2}^{T}\overline{Q}_{3} & \cdots \\ \overline{Q}_{1}^{T}\overline{Q}_{2}\overline{N} & \overline{Q}_{1}^{T}\overline{Q}_{2} & \overline{Q}_{1}^{T}\overline{Q}_{3} & \cdots \\ \overline{Q}_{1}^{T}\overline{Q}_{2}\overline{N} & \overline{Q}_{1}^{T}\overline{Q}_{2} & \overline{Q}_{1}^{T}\overline{Q}_{2} \\ \bullet & \overline{Q}^{T}\mathbf{Q} = \mathbf{Q}^{T}\mathbf{Q} = \mathbf{I} \quad \Rightarrow \quad \mathbf{Q}^{T} = \mathbf{Q}^{-1} \\ \bullet \text{ QR Algorithm Key Idea} \\ \begin{bmatrix} \overrightarrow{M}_{1} & \overrightarrow{M}_{2} & \cdots & \overrightarrow{M}_{N} \\ \vdots \\ \downarrow & \downarrow & \downarrow & \downarrow \\ \overline{P}_{1} & \overline{P}_{2} & \cdots & \overline{Q}_{N} \end{bmatrix} \begin{bmatrix} \overrightarrow{P}_{1} & \overrightarrow{P}_{1} & \overrightarrow{P}_{1} \\ \vdots \\ \overrightarrow{P}_{1} & \overrightarrow{P}_{2} & \cdots & \overrightarrow{P}_{N} \end{bmatrix} \begin{bmatrix} \overrightarrow{P}_{1} & \overrightarrow{P}_{1} & \overrightarrow{P}_{1} \\ \vdots \\ \overrightarrow{P}_{1} & \overrightarrow{P}_{2} & \cdots & \overrightarrow{P}_{N} \end{bmatrix} \begin{bmatrix} \overrightarrow{P}_{1} & \overrightarrow{P}_{1} & \overrightarrow{P}_{1} \\ \vdots \\ \overrightarrow{P}_{N} \end{bmatrix} = \begin{bmatrix} \overrightarrow{P}_{1} \\ \overrightarrow{P}_{2} \\ \vdots \\ \overrightarrow{P}_{N} \end{bmatrix} = \begin{bmatrix} \overrightarrow{P}_{1} \\ \vdots \\ \overrightarrow{P}_{N} \end{bmatrix} \\ \text{ We have a set of orthonormal columns that span the same subspace as the columns of M, then solving the system is very easy: \\ \hline{\mathbf{Q}^{T}} & \overrightarrow{P} & \overrightarrow{P} & \overrightarrow{P} \\ \hline{\mathbf{Q}}_{1} & \overrightarrow{P} & \overrightarrow{P} & \overrightarrow{P} \\ \hline{\mathbf{Q}}_{2} & \overrightarrow{P} & \overrightarrow{P} & \overrightarrow{P} \\ \hline{\mathbf{Q}}_{1} & \overrightarrow{P} & \overrightarrow{P} & \overrightarrow{P} & \overrightarrow{P} \\ \hline{\mathbf{Q}}_{1} & \overrightarrow{P} & \overrightarrow{P} & \overrightarrow{P} \\ \hline{\mathbf{Q}}_{1} & \overrightarrow{P} & \overrightarrow{P} & \overrightarrow{P} \\ \hline{\mathbf{Q}}_{2} & \overrightarrow{P} & \overrightarrow{P} & \overrightarrow{P} \\ \hline{\mathbf{Q}}_{1} & \overrightarrow{P} & \overrightarrow{P} & \overrightarrow{P} \\ \hline{\mathbf{Q}}_{2} & \overrightarrow{P} & \overrightarrow{P} \\ \hline{\mathbf{Q}}_{1} & \overrightarrow{P} & \overrightarrow{P} \\ \hline{\mathbf{Q}}_{2} & \overrightarrow{P} & \overrightarrow{P} \\ \hline{\mathbf{Q}}_{2} & \overrightarrow{P} & \overrightarrow{P} \\ \hline{\mathbf{Q}}_{1} & \overrightarrow{P} & \overrightarrow{P} \\ \hline{\mathbf{Q}}_{2} & \overrightarrow{P} & \overrightarrow{P} \\ \hline{\mathbf{Q}}_{1} & \overrightarrow{P} \\ \overrightarrow{P} & \overrightarrow{P} \\ \hline{\mathbf{Q}}_{2} & \overrightarrow{P} \\ \overrightarrow{P} & \overrightarrow{P} \\ \hline{\mathbf{Q}}_{2} & \overrightarrow{P} \\ \overrightarrow{P} \\ \overrightarrow{P} & \overrightarrow{P} \\ \overrightarrow{P} & \overrightarrow{P} \\ \overrightarrow{P} & \overrightarrow{P} \\ \overrightarrow{P} & \overrightarrow{P} \\ \overrightarrow{P} \\ \overrightarrow{P} & \overrightarrow{P} \\ \overrightarrow{P} \\ \overrightarrow{P} & \overrightarrow{P} \\ \overrightarrow{P} \\ \overrightarrow{P} \\ \overrightarrow{P} & \overrightarrow{P} \\ \overrightarrow{P} \\ \overrightarrow{P} \\ \overrightarrow{P} \\ \overrightarrow{P} \\ \overrightarrow{P} & \overrightarrow{P} \\ \overrightarrow{P} \end{array}$$



Given \vec{M}_1, \vec{M}_2 , find $\vec{Q}_2 = \vec{M}_2 - r_{12}\vec{M}_1$ so that

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Another way to write this is: $\,M$ is factored into Q and R

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$ \frac{\vec{M}_1 \cdot \left(\vec{M}_3 - r_{13}\vec{M}_1 - r_{23}\vec{M}_2\right) = 0}{\vec{M}_2 \cdot \left(\vec{M}_2 - r_{12}\vec{M}_1 - r_{23}\vec{M}_2\right) = 0} = 0 $	$\Rightarrow \begin{bmatrix} \vec{M}_1 \cdot \vec{M}_1 & \vec{M}_1 \cdot \vec{M}_2 \\ \vec{M}_2 \cdot \vec{M}_1 & \vec{M}_2 \cdot \vec{M}_2 \end{bmatrix} \begin{bmatrix} r_{13} \\ r_{23} \end{bmatrix} = \begin{bmatrix} \vec{M}_1 \cdot \vec{M}_3 \\ \vec{M}_2 \cdot \vec{M}_3 \end{bmatrix}$
To Orthogonalize the Nth Vector	n a state a sta
$\begin{bmatrix} \vec{M}_1 \cdot \vec{M}_1 & \cdots & \vec{M}_1 \cdot \vec{M}_N \end{bmatrix}$	$[r_{1,N}] \begin{bmatrix} \vec{M}_1 \cdot \vec{M}_N \end{bmatrix}$
1 2 1	: . :
$\vec{M}_{N-1} \cdot \vec{M}_1 \cdots \vec{M}_{N-1} \cdot \vec{M}_N$	$\vec{r}_{N-1} = \vec{r}_{N-1,N} = \vec{M}_N \cdot \vec{M}_N$
N^2 inner products or N^3	work.

$\begin{split} \vec{M}_3 &= \eta_3 \vec{Q}_1 + r_{23} \vec{Q}_2 + r_{33} \vec{Q}_3 \\ \vec{\tilde{M}}_3 &= \vec{M}_3 - (\vec{Q}_1 \cdot \vec{M}_3) \vec{Q}_1 - (\vec{Q}_2 \cdot \vec{M}_3) \\ \Rightarrow \vec{\tilde{M}}_3 &= \vec{M}_3 - \eta_3 \vec{Q}_1 - r_{23} \vec{Q}_2 \\ \vec{Q}_3 &= \frac{\vec{\tilde{M}}_3}{\ \vec{\tilde{M}}_3\ } \end{split}$	$\vec{\mathcal{Q}}_{13} = \vec{\mathcal{Q}}_1 \cdot \vec{\mathcal{M}}_3$ $\vec{\mathcal{Q}}_{2} \qquad \vec{\mathcal{P}}_{23} = \vec{\mathcal{Q}}_2 \cdot \vec{\mathcal{M}}_3$ $\vec{\mathcal{P}}_{33} = \vec{\mathcal{Q}}_3 \cdot \vec{\mathcal{M}}_3 = \left\ \vec{\tilde{\mathcal{M}}}_3 \right\ $
$\begin{split} \vec{M}_1 & \vec{Q}_1 \leftarrow \frac{\vec{M}_1}{\left \vec{M}_1\right } \checkmark^{P_{11}} \\ \vec{M}_2 & \vec{M}_2 \leftarrow \vec{M}_2 - (\underbrace{\vec{Q}_1 \cdot \vec{M}_2}_{r_2}) \underbrace{\vec{Q}_1}_{r_2} \\ \vec{M}_3 & \vec{M}_3 \leftarrow \vec{M}_3 - (\underbrace{\vec{Q}_1 \cdot \vec{M}_3}_{r_1}) \underbrace{\vec{Q}_1}_{r_2} - (\underbrace{\vec{Q}_2 \cdot \vec{M}_2}_{r_2}) \\ \vec{M}_4 & \vec{M}_4 \leftarrow \vec{M}_4 - (\underbrace{\vec{Q}_1 \cdot \vec{M}_4}_{r_4}) \underbrace{\vec{Q}_1}_{r_4} - (\underbrace{\vec{Q}_2 \cdot \vec{M}_4}_{r_4}) \\ \hline \end{split}$	$\begin{array}{c} Q_{2} \leftarrow \frac{\tilde{M}_{2}}{[\tilde{M}_{2}]} \checkmark r_{22} \\ \hline t_{3} [\tilde{Q}_{2} & \bar{Q}_{3} \leftarrow \frac{\tilde{M}_{3}}{[\tilde{M}_{3}]} \checkmark r_{33} \\ \hline t_{4} [\tilde{Q}_{2} - (\underline{\tilde{Q}}_{3} \cdot \underline{\tilde{M}}_{4}) \overline{\tilde{Q}}_{3} & \bar{Q}_{4} \leftarrow \frac{\tilde{M}_{4}}{[\tilde{M}_{4}]} \checkmark r_{44} \end{array}$
: $\vec{M}_N \vec{M}_N \leftarrow \vec{M}_N - \underbrace{(\vec{\mathcal{O}}_1 \cdot \vec{M}_N)}_{\eta_N} \overleftarrow{\mathcal{O}}_1 - \underbrace{(\vec{\mathcal{O}}_2 \cdot \vec{M}_N)}_{\eta_N} = \underbrace{\vec{\mathcal{O}}_2 \cdot \vec{\mathcal{O}}_1}_{\eta_N} = \underbrace{\vec{\mathcal{O}}_2 \cdot \vec{\mathcal{O}}_1}_{\eta_N} = \underbrace{\vec{\mathcal{O}}_2 \cdot \vec{\mathcal{O}}_1}_{\eta_N} = \underbrace{\vec{\mathcal{O}}_2 \cdot \vec{\mathcal{O}}_2}_{\eta_N} = \underbrace{\vec{\mathcal{O}}_2 \cdot \vec{\mathcal{O}}_2}_{\eta$	(\widetilde{M}_N) $(\widetilde{Q} - \dots - (\widetilde{Q}_{N-1} \cdot \widetilde{M}_N))$ (\widetilde{Q}_{N-1}) $(\widetilde{r}_{N-1,N})$
$\begin{bmatrix} \vec{M}_1 & \vec{M}_2 & \vec{M}_3 \end{bmatrix} = \begin{bmatrix} \mathbf{X} \\ \mathbf{Q} \end{bmatrix}$	$\vec{\tilde{M}}_{2} \vec{\tilde{M}}_{3}$
$\begin{bmatrix} \vec{M}_1 & \vec{M}_2 & \vec{M}_3 \end{bmatrix} = \begin{bmatrix} \vec{Q}_1 & \vec{Q}_1 \end{bmatrix}$	$\tilde{\mathbf{X}}_{2} \tilde{\mathbf{X}}_{3} \begin{bmatrix} r_{11} & r_{12} & r_{13} \\ r_{22} & r_{23} \end{bmatrix}$



				$\vec{Q}_1 \leftarrow$	$\frac{\vec{M}_1}{\vec{M}_1}$								
\vec{M}_2	\leftarrow	\vec{M}_2	-	$\left(\vec{Q}_1\cdot\vec{M}\right)$	$_{2})\vec{Q}_{1}$		\vec{Q}_2	$\leftarrow \frac{\vec{M}}{\ \vec{M}\ }$	2				-
\vec{M}_3	\leftarrow	\vec{M}_3	-	$\left(\vec{Q}_1\cdot\vec{M}\right)$	$_{3})\vec{\mathcal{Q}}_{1}$	-	$(\bar{Q}_2$	(\vec{M}_3)	$\vec{2}_2$		$\vec{Q}_3 \leftarrow$	$-\frac{\vec{M}_3}{\vec{M}_3}$	
\vec{M}_4	\leftarrow	\vec{M}_4	-	$\left(\vec{Q}_1\cdot\vec{M}\right)$	$_{4})\vec{Q}_{1}$	-	$(\vec{Q}_2 \cdot$	(\vec{M}_4)	$\vec{2}_2$	-	$(\vec{Q}_3 \cdot$	$(\vec{M}_4)\vec{Q}_3$	
		$ \begin{array}{c} \stackrel{\uparrow}{\underset{\scriptstyle i}{\underset{\scriptstyle j}{\underset{\scriptstyle j}{\atop\scriptstyle j}{\underset{\scriptstyle j}{\underset{\scriptstyle j}{\atop\scriptstyle j}{\underset{\scriptstyle j}{\underset{\scriptstyle j}{\atop\scriptstyle j}{\underset{\scriptstyle j}{\atop\scriptstyle j}{\underset{\scriptstyle j}{\atopi}{\atopi}{\atop_{\scriptstyle j}{\atopi}{\atop_{\scriptstyle j}{\atop_{\scriptstyle j}{\atop_{\scriptstyle j}{\atopi}{\atop_{\scriptstyle j}{\atop_{\scriptstyle j}{\atop_{l}{i}{\atop_{\scriptstyle j}{\atop_{\scriptstyle j}{\atop_{\scriptstyle j}{\atop_{\scriptstyle j}{\atop_{l}}{\atop_{\scriptstyle j}{\atop_{l}}{}}{i}}}}}}}}}}}}}}}}}}}}}}}}}}}}}}$		$\left(\frac{\gamma_1}{\sqrt{M_1+M_2}}\right)$ (6.3)	$\begin{bmatrix} \uparrow & & \\ 0 & 0 \\ 0 & 0 \\ 0 & 0 \\ 0 & 0 \\ 0 & 0 \end{bmatrix}$) (0) (0) (1) (1) (1) (1) (1) (1) (1) (1) (1) (1	$\begin{pmatrix} \mathbf{x}_{i} \\ \mathbf{x}_{i} \end{pmatrix}$		13 14 12 12 12 12 12 12 12 12] M₄ ↓ rd	(ni na ras step		

★ Basic Algorithm
o Source Row Approach
For i = 1 to N { "For each Source Column" $\vec{r}_{ii} = \sqrt{\vec{M}_i \cdot \vec{M}_i}$ Normalize $\sum_{i=1}^{N} 2N \approx 2N^2$ operations $\vec{Q}_i = \frac{1}{r_{ii}} \vec{M}_i$ For j = i + 1 to N { "For each target column right of source" $\vec{r}_{ij} \leftarrow \vec{M}_j \cdot \vec{Q}_i$ $\vec{M}_j \leftarrow \vec{M}_j - r_{ij}\vec{Q}_i$ $\sum_{i=1}^{N} (N - i)2N \approx N^3$ operations







		Complexity
Step 1)	Factor $\mathbf{M} = \mathbf{Q}\mathbf{R}$	$O(N^3)$
0. 0.	(Q orthonormal, R upper triangular)	O(N)
Step 2)	Solve $\mathbf{Q}\vec{\mathbf{v}} = b$.	0.012
	$\left(\text{Very easy}: \ \vec{y} = \mathbf{Q}^T \vec{b} \right)$	$O(N^{-})$
Step 3)	Backsolve the triangular system $R\vec{x} = \vec{y}$	$O(N^2)$
Step 3)	Solve $\mathbf{Q}^{T} = \mathbf{Q}^{T} \vec{b}$ (Very easy : $\vec{y} = \mathbf{Q}^{T} \vec{b}$) Backsolve the triangular system $R\vec{x} = \vec{y}$	$O(N^2)$ $O(N^2)$

 $\begin{array}{l} \mathbf{M}\vec{x} = \vec{b} \\ \text{Step 1} \end{array} \qquad \begin{array}{l} \text{Factor } \mathbf{M} = \mathbf{Q}\mathbf{R} \\ \mathbf{M}\vec{x} = \vec{b} \Leftrightarrow \mathbf{Q}\mathbf{R}\vec{x} = \vec{b} \\ \vec{y} \end{array}$ $\begin{array}{l} \text{Step 2} \end{array} \qquad \begin{array}{l} \text{Solve for } \vec{y} : \mathbf{Q}\vec{y} = \vec{b} \end{array}$



Step 3)



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Computational

INTRODUCTION TO NUMERICAL SIMULATION

LECTURES 7 & 8.

QR and Krylov-Subspace Matrix Solution Methods

TODAY'S OUTLINE:

* Minimization View of QR

- Singular Matrix
- Basic Minimization Approach
- Orthogonalized Search Directions
- QR and Length Minimization Produce Identical Results
- * Arbitrary Subspace Algorithm
- Orthogonalization of Search Directions
- * Generalized Conjugate Residual Algorithm
 - Krylov-subspace
 - Simplification in the symmetric case
 - Leaky and insulating examples

QR FACTORIZATION

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Example.	Г	2 -1	0	1		
	A = -	-1 2	-1			
	L	0 -1	2			
	[] (0 0][2	-1	0]		
	$-\frac{1}{2}$ 1		$\frac{3}{2}$	-1		
		$\frac{1}{3}$ 1 0	0	$\frac{4}{3}$		
	L		U			
	-0.89	-0.35	0.27]	-2.2	1.8	-0.45
	0.44	-0.72	0.53	0	-1.7	1.9
	0	0.60	0.80	0	0	1.1
		Q			R	

Matrix is Singular, column of Q is zero * Zero Column If a column is zero $\vec{M}_i = \sum_{j=1}^{i+1} w_j \vec{M}_j$ j=1

 $\{\vec{M}_1, ..., \vec{M}_i\}$ not linearly independent

What if a column becomes zero? [n.

N

$$\begin{bmatrix} \uparrow & \uparrow & \uparrow & \uparrow \\ \vec{Q}_1 & 0 & \vec{M}_3 & \cdots & \vec{M}_N \\ \downarrow & \downarrow & \downarrow & \downarrow \\ \end{bmatrix} \begin{bmatrix} r_{11} & r_{12} & r_{13} & \cdots & r_{1N} \\ 0 & 0 & 0 & \cdots & 0 \\ 0 & 0 & 0 & \cdots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \cdots & 0 \end{bmatrix}$$
$$\begin{bmatrix} \uparrow & \uparrow & \uparrow & \uparrow \\ \vec{Q}_1 & 0 & \vec{Q}_3 & \cdots & \vec{Q}_N \\ \downarrow & \downarrow & \downarrow & \downarrow \\ \end{bmatrix} \begin{bmatrix} r_{11} & r_{12} & r_{13} & \cdots & r_{1N} \\ 0 & 0 & 0 & \cdots & 0 \\ 0 & 0 & r_{33} & \cdots & r_{3N} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & 0 & \cdots & r_{NN} \end{bmatrix}$$
Matrix MUST be Singular!!
1. Do not try to normalize the column.
2. Do not use the column as a source for orthogonalization.
3. Perform backward substitution as well as possible.

$$\mathbf{QR}\vec{x} = \vec{b} \Rightarrow \mathbf{R}\vec{x} = \mathbf{Q}^T \vec{b}$$
$$\Rightarrow \begin{bmatrix} r_{11} & r_{12} & r_{13} & \cdots & r_{1N} \\ 0 & 0 & 0 & \cdots & 0 \\ r_{33} & \cdots & r_{3N} \\ & \ddots & \vdots \\ r_{NN} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ \vdots \\ x_N \end{bmatrix} = \begin{bmatrix} \leftarrow & \vec{Q}_1 & \rightarrow \\ 0 & \cdots & 0 \\ \leftarrow & \vec{Q}_3 & \rightarrow \\ \vdots \\ \leftarrow & \vec{Q}_N & \rightarrow \end{bmatrix}$$
QR Factorization

Problem $\mathbf{M}\vec{x} = \vec{b}$: (1) $\mathbf{Q} \cdot \underbrace{\mathbf{R}}_{\vec{y}} = \vec{b}$ (2) Solve $\mathbf{Q}\vec{v} = \vec{b} \implies \vec{y} = \mathbf{Q}^T \vec{b}$ (3) Solve $\mathbf{P}\vec{z} = \vec{v} \implies \vec{z}$

(3) Solve
$$\mathbf{R}x = y \rightarrow x$$

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 \uparrow \overline{b} \downarrow Singular Example









If $\tilde{b} \notin span\{M\} \rightarrow$ there is no solution



choose this \vec{b} outside of span {M} knowing that we cannot get a solution). What is going on?

We find that \vec{b} is the projection of \vec{b} onto the subspace of the columns of M.

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Minimization View of QR










One-dimensional minimization yields same result as orthogonal projection on the column!



 $\|R(\vec{x})\|^{2} = R(\vec{x})^{T} R(\vec{x}) = (\vec{b} - x_{1}\mathbf{M}\vec{e}_{1} - x_{2}\mathbf{M}\vec{e}_{2})^{T} (\vec{b} - x_{1}\mathbf{M}\vec{e}_{1} - x_{2}\mathbf{M}\vec{e}_{2})$ $=\vec{b}^{T}\vec{b}-2x_{1}\vec{b}^{T}\mathbf{M}\vec{e}_{1}+x_{1}^{2}(\mathbf{M}\vec{e}_{1})^{T}(\mathbf{M}\vec{e}_{1})-2x_{2}\vec{b}^{T}\mathbf{M}\vec{e}_{2}+x_{2}^{2}(\mathbf{M}\vec{e}_{2})^{T}(\mathbf{M}\vec{e}_{2})$ $+\frac{2x_1x_2(\mathbf{M}\vec{e}_1)^T(\mathbf{M}\vec{e}_2)}{(\mathbf{M}\vec{e}_1)^T(\mathbf{M}\vec{e}_2)} > coupling term.$ More general search directions $x = y_1 \overline{p}_1 + y_2 \overline{p}_2$ such that: $\mathbf{M}x = y_1 \mathbf{M} \vec{p}_1 + y_2 \mathbf{M} \vec{p}_2 \quad span{\{\vec{p}_1, \vec{p}_2\}} = span{\{\vec{e}_1, \vec{e}_2\}}$ $\mathbf{R}(\vec{x})^T \mathbf{R}(\vec{x}) = \vec{b}^T \vec{b} - 2y_1 \vec{b}^T \mathbf{M} \vec{p}_1 + y_1^2 (\mathbf{M} \vec{p}_1)^T (\mathbf{M} \vec{p}_1)$ $-2y_{2}\bar{b}^{T}\mathbf{M}\bar{p}_{2}+y_{2}^{2}(\mathbf{M}\bar{p}_{2})^{T}(\mathbf{M}\bar{p}_{2})+2y_{1}y_{2}(\mathbf{M}\bar{p}_{1})^{T}(\mathbf{M}\bar{p}_{2})$ Jung term If $\vec{p}_1^T \mathbf{M}^T \mathbf{M} \vec{p}_2 = 0$ Minimizations Decouple!!! If search directions are M-orthonormal $\underbrace{\mathbf{M}\vec{p}_i}_{\vec{Q}_i} \cdot \underbrace{\mathbf{M}\vec{p}_j}_{\vec{Q}_j} = \begin{cases} 0 & i \neq j \\ 1 & i = j \end{cases}$ Decoupled minimizations can be done individually Minimize: $-2v_i(\vec{b}^T \mathbf{M} \vec{p}_i) + y_i^2 (\mathbf{M} \vec{p}_i)^T (\mathbf{M} \vec{p}_i)$ $-2y_i(\vec{b}\cdot\vec{Q}_i)+y_i^2(\vec{Q}_i\cdot\vec{Q}_i)$ Differentiating: $-2(\vec{b}\cdot\vec{Q}_i)+2y_i=0 \Rightarrow y_i=\vec{b}\cdot\vec{Q}_i$ Minimization yields same result as orthogonal projection!

Now $x = x_1e_1 + x_2e_2$ and $Mx = x_1Me_1 + x_2e_3$ <u>Residual Minimization</u>



$$\begin{split} & \{\hat{\boldsymbol{e}}_{1}, \hat{\boldsymbol{e}}_{2}, ..., \hat{\boldsymbol{e}}_{i}\} \\ & \{\vec{p}_{1}, \vec{p}_{2}, ..., \vec{p}_{i}\} \\ & \vec{x}^{i} = \underbrace{(\vec{b} \cdot \mathbf{M}\vec{p}_{1})}_{y_{1}} \vec{p}_{1} + ... + \underbrace{(\vec{b} \cdot \mathbf{M}\vec{p}_{i})}_{y_{i}} \vec{p}_{i} \\ & \vec{x}^{i+1} = \underbrace{(\vec{b} \cdot \mathbf{M}\vec{p}_{1})}_{x^{i}} \vec{p}_{1} + ... + \underbrace{(\vec{b} \cdot \mathbf{M}\vec{p}_{i})}_{x^{i}} \vec{p}_{i} + \underbrace{(\vec{b} \cdot \mathbf{M}\vec{p}_{i})}_{x^{i}} \vec{p}_{i+1} \end{split}$$

Forming M-orthonormal minimization directions The *i*th search direction equals M-orthonormalized unit vector











Example. $\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \begin{bmatrix} 1 \\ 2 \\ 3 \end{bmatrix}$ What is the best initial search direction? Mp $\vec{b} = 2$ 3 $\vec{p}_1 = \begin{bmatrix} 1 \\ 2 \\ 3 \end{bmatrix}_{1}$ \mathbf{M} is the identity, so $\mathbf{M}\vec{b} = \vec{b}$ $\Rightarrow \vec{p}_1 = \frac{\vec{b}}{\|\vec{b}\|}$ $y_1 = \vec{b} \cdot \mathbf{M} \vec{p}_1 = \left\| \vec{b} \right\|$ x-space $y_1 \cdot \mathbf{M}\vec{p}_1 = \left|\vec{b}\right| \cdot \frac{\vec{b}}{\left|\vec{b}\right|} = \vec{b}$ Try { \vec{b} ,...} $\vec{b} \xrightarrow{\mathbf{M}} \mathbf{M} \vec{b} = \vec{b}$ Normalize $\vec{p}_{1} = \frac{\vec{b}}{\|\mathbf{M}\vec{b}\|} = \frac{\vec{b}}{\|\vec{b}\|} \xrightarrow{\mathbf{M}} \frac{\mathbf{M}\vec{b}}{\|\mathbf{M}\vec{b}\|} = \frac{\vec{b}}{\|\vec{b}\|}$ $\vec{p}_{1} = y_{1}\mathbf{M}\vec{p}_{1}$ $y_{1} = \vec{b} \cdot \mathbf{M}\vec{p}_{1} = \|\vec{b}\|$ $\vec{p}_{1} = \|\vec{b}\| \frac{I \cdot \vec{b}}{\|\vec{b}\|} = \vec{b}$ One Step Use QR to solve



x minimizes $(b - \mathbf{M}x)^T (b - \mathbf{M}x)$ $\mathbf{M}z = 0$ z nonzero exists if **M** is singular

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ARBITRARY SUBSPACE ALGORITHM



Computational Approach

$$\begin{split} & \text{Minimizing } \left\| \vec{r}^{k} \right\|_{2}^{2} = \left\| \vec{b} - \sum_{i=0}^{k-1} \alpha_{i} \mathbf{M} \vec{w}_{i} \right\|_{2}^{2} \text{ is easy if } \mathbf{M} \vec{w}_{i} \text{ were orthonormal!} \\ & \text{Create a set of vectors } \{ \vec{p}_{0}, \vec{p}_{1}, ..., \vec{p}_{k-1} \} \text{ such that} \\ & span\{ \vec{p}_{0}, ..., \vec{p}_{k-1} \} = span\{ \vec{w}_{0}, ..., \vec{w}_{k-1} \} \\ & \mathbf{M} \vec{p}_{i} \cdot \mathbf{M} \vec{p}_{j} = \begin{cases} 0 & i \neq j \\ 1 & i = j \end{cases} \end{split}$$

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Gram-Schnidt on
$$\mathbf{M}\tilde{w}_{i}$$
's
 $\overline{w}_{0} \rightarrow \overline{p}_{0}$
 $\overline{w}_{1} \rightarrow \overline{p}_{1}$ want $(\mathbf{M}\tilde{p}_{1}) \cdot (\mathbf{M}\tilde{p}_{0}) = 0$
 $= (\mathbf{M}\tilde{w}_{1} - \mathbf{M}\tilde{a}p_{0})^{T} (\mathbf{M}\tilde{p}_{0}) = 0$
 $\Rightarrow (\mathbf{M}\tilde{w}_{1} - \mathbf{M}\tilde{a}p_{0})^{T} (\mathbf{M}\tilde{p}_{0}) = 0$
 $= [\overline{\mathcal{Q}}_{1} \cdots \mathbf{M}\tilde{\mathcal{M}}_{N}] \Rightarrow [\mathbf{M}\hat{e}_{1} \cdots \mathbf{M}\hat{e}_{N}]$
columns orthogonal
 $= [\overline{\mathcal{Q}}_{1} \cdots \overline{\mathcal{Q}}_{N}][\mathbf{R}]$
M-space
 $\mathbf{M}\tilde{p}_{i} \leftarrow \mathbf{M}\tilde{w}_{i} - \sum_{j=0}^{i-1} (\underbrace{\mathbf{M}\tilde{p}_{j} \cdot \mathbf{M}\tilde{w}_{i}}_{v_{j}}) \underbrace{\mathbf{M}\tilde{p}_{j}}_{v_{j}}$
x-space
 $\tilde{p}_{i} \leftarrow \tilde{w}_{i} - \sum_{j=0}^{i-1} (\mathbf{M}\tilde{p}_{j} \cdot \mathbf{M}\tilde{w}_{i})]\tilde{p}_{j}$
Arbitrary Subspace Solution Algorithm
Given M, δ and a set of search directions $\{\tilde{w}_{0}, \tilde{w}_{1}, ..., \tilde{w}_{k}\}$
() orthogonalize the $\overline{W}\tilde{v}$;
 $\hat{v}_{i} \cdot i = 1 \text{ to } k - \tilde{p}_{k} = \tilde{w}_{k} - \sum_{i=0}^{k-1} (\mathbf{M}\tilde{p}_{i})^{T} (\mathbf{M}\tilde{p}_{i}) \tilde{p}_{i}$
(2) compute the \tilde{r} minimizing solution x^{k+1}
 $\tilde{x}^{k+1} = \sum_{i=0}^{k} (\frac{r^{0}T}{(\mathbf{M}\tilde{p}_{i})} \tilde{p}_{i}$









GENERALIZED CONJUGATE RESIDUAL ALGORITHM





Idea: search along the gradient i.e. the steepest descent directions for f(x) i.e. along

 $\text{Pick } span\{\vec{w}_{0},...,\vec{w}_{k-1}\} = span\{\nabla_{x} f(\vec{x}^{0}),...,\nabla_{x} f(\vec{x}^{k-1})\} = span\{\vec{r}^{0}, \vec{r}^{1},...,\vec{r}^{k-1}\}$

Gradient of f(x) = Residual at x

Krylov Subspace

the current residual direction.

Does not extend to non-symmetric, non positive definite case.

Note: $span\{\nabla_x f(\vec{x}^0),...,\nabla_x f(\vec{x}^{k-1})\} = span\{\vec{y}^{(0)},...,\vec{r}^{k-1}\}$





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Gaussian Elimination cost = O(n)

Algorithm Cost for iteration k For j = 1 to k - 1 O(k) inner products: O(kn) mult. $\bar{p}_k \leftarrow \bar{p}_k - (\mathbf{M}\bar{p}_j) \cdot (\mathbf{M}\bar{p}_k)\bar{p}_j$ If M is sparse the inner products are the dominant cost: total cost = $O(n) + O(2n) + ... + O(kn) = O(k^2n)$ As k (# of iters) approaches N, total cost = $O(n^3)$ Better Converge Fast!! Symmetric Case – An amazing fact that will not be derived If $\mathbf{M} = \mathbf{M}^T$ then $\mathbf{M}^{Tk} \perp \mathbf{M}\bar{p}_j$ for j < k - 2For \mathbf{M} to k - 1 Only ONE inner products: O(n) mult. $\bar{p}_k \leftarrow \bar{p}_k - (\mathbf{M}\bar{p}_{k-1}) \cdot (\mathbf{M}\bar{p}_k)\bar{p}_{k-1}$ If k (# of iters) $\rightarrow n$, then symmetric, sparse, GCR is worst case $O(n^2)$ Better Converge Fast!!

			SP	ARSE	DENSE
$\mathbf{M}\vec{r}^{k}$	Matrix	c-Vector oduct		O(n)	$O(n^2)$
$\sum (M \bar{r}^{k})^{T} (M$	(\bar{p}_i) M-Orth	M-Orthogonalize		Q(kn) ((n)	$O(\kappa n) O(n)$
$\left(\bar{r}^{k}\right)^{T}(\mathbf{M}\bar{p}_{k})$) Optin Size	Optimal Step Size Update		O(n)	<i>O</i> (<i>n</i>)
	Т	Total		$O(Kn) + O(K^2n)$ inner products cost	$\underbrace{O(Kn^2)}_{\substack{\text{matrix-vector}\\ \text{product cost}}} + \underbrace{O(Kn}_{\substack{\text{inner}\\ \text{products}}}$
	Wor As	Worst Case As $K \Rightarrow n$		$\frac{1}{O(n^2)} + O(n^2)$ to reduce k!!	$\frac{O(n^3)}{O(n^3)} + O(n^3)$ Need to reduce k AND matrix- vector cost
	Bes For (5-10 i	t Case small <i>K</i> iterations)	$\underbrace{O(n)+O(n)}_{O(n)}$		$\underbrace{O(n^2) + O(n)}_{O(n^2)}$
				Note: If ve	"fast" $\{=O(n)\}$ matr ctor product then $O($
mparison -	- GCR & LU			tot	tal also for delise cas
mparison -	- GCR & LU Sparse Symmetric	Sparse	metric	Dense	an also for delise cas
GCR	- GCR & LU Sparse Symmetric	Sparse 2 Asymr	$\frac{1}{2n}$	Dense $O(Kn^2)$	an also for dense cas
mparison - GCR LU	- GCR & LU Sparse Symmetric O(Kn) $O(n^{1.2-1.8})$	Sparse Sparse O(K) $O(n^{1.2})$	$\frac{metric}{(2^n)}$	Dense $O(\kappa n^2)$ $O(n^3)$	ar also for delise cas
GCR LU • Makir • If the LU fo	-GCR & LU Sparse Symmetric O(Kn) $O(n^{1.2-1.8})$ g GCR conv number of itu r dense prob	Sparse c Asymm O(K) $O(n^{1.2})$ rerge fast cc erations k is lems.	metric $\binom{2n}{2-1.8}$ an be ves about	Dense $O(Kn^2)$ $O(n^3)$ ry problem s constant and	pecific small GCR is much :
GCR UU • Makir • If the LU fo	-GCR & LU Sparse Symmetric O(Kn) $O(n^{1.2-1.8})$ og GCR conv number of ite r dense prob Sparse O(n)	Sparse c Asymm O(K) $O(n^{1.2})$ rerege fast ca erations k is lems. Dense $O(n^2)$	metric $\binom{2^{2}n}{2^{-1.8}}$ an be ve	Dense $O(\kappa n^2)$ $O(n^3)$ ry problem s constant and	pecific small GCR is much :













INTRODUCTION TO NUMERICAL SIMULATION



ω 0 Linear Systems I. GCR Convergence

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$$\begin{split} \mathbf{M} &= \mathbf{U} \Lambda \mathbf{U}^{-1} & \text{Real Matrix} \\ \lambda \text{ eigenvalue } &\Leftrightarrow & (\mathbf{M} \cdot \lambda \mathbf{I}) \text{ is singular} \\ & \mathbf{U} \Lambda \mathbf{U}^{-1} - \lambda \mathbf{U} \mathbf{I} \mathbf{U}^{-1} \rightarrow \text{ singular} \end{split}$$



Symmetric Matrices $\longrightarrow \mathbf{U}^T \mathbf{U} = \mathbf{I}$ Orthonormal Eigenvectors $\mathbf{M} = \mathbf{U} \wedge \mathbf{U}^{-1}$ $\mathbf{M}^T = (\mathbf{U}^{-1})^T \wedge \mathbf{U}^T$ Symmetric $\longrightarrow \mathbf{M} = \mathbf{M}^T$ $\mathbf{U} \wedge \mathbf{U}^{-1} = (\mathbf{U}^{-1})^T \wedge \mathbf{U}^T \Leftrightarrow \mathbf{U}^{-1} = \mathbf{U}^T$ $\longrightarrow \mathbf{U}$ orthonormal

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Spectral Radius



The spectral Radius of M is the radius of the smallest circle, centered at the origin, which encloses all of M's eigenvalues.

Suppose $\vec{y} = \mathbf{M}\vec{x}$

How much larger is \vec{y} than \vec{x} ? OR How much does M magnify \vec{x} ?







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Unit Length Rod







 $\begin{aligned} & \frac{\operatorname{Krylov}\operatorname{Subspace}}{\operatorname{Note}: \operatorname{span}\left\{\nabla_{x}f\left(x^{0}\right)...,\nabla_{x}f\left(x^{k-1}\right)\right\} = \operatorname{span}\left\{v^{0},...,r^{k-1}\right\} \\ & \text{If} \quad \operatorname{span}\left\{\vec{w}_{0},...,\vec{w}_{k-1}\right\} = \operatorname{span}\left\{v^{0},...,r^{k-1}\right\} \\ & \text{then} \quad r^{k} = r^{0} - \sum_{i=0}^{k-1} \alpha_{i}Mr^{i} \\ & \text{and} \quad \operatorname{span}\left\{v^{0},...,r^{k-1}\right\} = \underbrace{\operatorname{span}\left\{v^{0},\mathbf{M}r^{0},...,\mathbf{M}^{k-1}r^{0}\right\}}_{\operatorname{Krylov}\operatorname{Subspace}} \end{aligned}$

270

$$\begin{split} & \begin{array}{l} k = 1 \\ & \text{Step 1.} \\ & \overline{p}_{0} \leftarrow \frac{\overline{r}^{0}}{\left\|\mathbf{M}\overline{r}^{0}\right\|} \in span\left\{\overline{r}^{0}\right\} \\ & y_{0} \leftarrow \overline{r}^{0}^{T}\mathbf{M}\overline{p}_{0} \quad \text{scalar} \\ & \overline{x}^{1} \leftarrow 0 + y_{0}\overline{p}_{0} \quad \in span\left\{\overline{p}_{0}\right\} = span\left\{\overline{r}^{0}\right\} \\ & \overline{r}^{1} \leftarrow \overline{r}^{0} - y_{0}\mathbf{M}\overline{p}_{0} \quad \in span\left\{\overline{r}^{0}, \mathbf{M}\overline{p}_{0}\right\} = span\left\{\overline{r}^{0}, \mathbf{M}\overline{r}^{0}\right\} \\ & k = 2 \\ & \text{Step 2.} \\ & \overline{p}_{1} \leftarrow \overline{r}^{1} - \frac{scalar}{(\mathbf{M}\overline{p}_{0})^{T}(\mathbf{M}\overline{r}^{1})}\overline{p}_{0} \quad \in span\left\{\overline{r}^{1}, \overline{p}_{0}\right\} = span\left\{\overline{r}^{0}, \mathbf{M}\overline{r}^{0}\right\} \\ & p_{1} \leftarrow \frac{\overline{p}_{1}}{\left\|\mathbf{M}\overline{p}_{1}\right\|} \quad \in span\left\{\overline{r}^{0}, \mathbf{M}\overline{r}^{0}\right\} \\ & y_{1} \leftarrow \overline{r}^{1}^{T}\mathbf{M}\overline{p}_{0} \quad \text{scalar} \\ & \overline{x}^{2} \leftarrow \overline{x}^{1} + y_{1}\overline{p}_{1} \quad \in span\left\{\overline{x}^{1}, \mathbf{M}\overline{p}_{1}\right\} = span\left[\overline{r}^{0}, \mathbf{M}\overline{r}^{0}\right\} \\ & \overline{r}^{2} \leftarrow \overline{r}^{1} - y_{1}\mathbf{M}\overline{p}_{1} \quad \in span\left\{\overline{x}^{1}, \mathbf{M}\overline{p}_{1}\right\} = span\left[\overline{r}^{0}, \mathbf{M}\overline{r}^{0}, \mathbf{M}^{2}\overline{r}^{0}\right\} \\ & \overline{x}^{k} \leftarrow \alpha_{0}\mathbf{I}\overline{r}^{0} + \alpha_{1}\mathbf{M}\overline{r}^{0} + \ldots + \alpha_{k-1}\mathbf{M}^{k-1}\overline{r}^{0} = \xi_{k-1}(\mathbf{M})\overline{r}^{0} \\ & \overline{r}^{k} \leftarrow \beta_{0}\overline{r}^{0} + \beta_{1}\mathbf{M}\overline{r}^{0} + \ldots + \beta_{k}\mathbf{M}^{k}\overline{r}^{0} = sp_{k}(\mathbf{M})\overline{r}^{0} \end{split}$$

$$\begin{split} & \frac{|y_{10}| (|\mathbf{n}|\mathbf{n}| + \mathbf{v}|\mathbf{k}|\mathbf{m}|}{\tilde{p}_{k-1}} \leftarrow r^{k-1} - (\mathbf{M}\tilde{\rho}_{k-2})^T \left(\tilde{r}^{k-1}\right) \tilde{p}_{k-2} \\ & \bar{p}_{k-1} \in span \left[\tilde{r}^0, \mathbf{M}\tilde{r}^0, \dots, \mathbf{M}^{k-1}\tilde{r}^0 \right] \\ & \bar{x}^k \leftarrow \bar{x}^{k-1} + y_{k-1} \tilde{p}_{k-1} \\ & \bar{x}^k \in span \left[\tilde{r}^0, \mathbf{M}\tilde{r}^0, \dots, \mathbf{M}^{k-1}\tilde{r}^0 \right] \\ & \bar{x}^k = \xi_{k-1} (\mathbf{M})\tilde{r}^0 \rightarrow (k-1)^{\text{th}} \text{ order polynomial in } \mathbf{M} \\ & \bar{r}^k \leftarrow \bar{r}^{k-1} - y_{k-1} \mathbf{M}\tilde{p}_{k-1} \\ & \bar{r}^k \in span \left[\tilde{r}^0, \mathbf{M}\tilde{r}^0, \dots, \mathbf{M}^{k-1}\tilde{r}^0, \mathbf{M}^k \tilde{r}^0 \right] \\ & \bar{r}^k = \varphi_k (\mathbf{M})\tilde{r}^0 \rightarrow k^{\text{th}} \text{ order polynomial in } \mathbf{M} \end{split}$$

$$\begin{split} & \varphi_k \left(0 \right) \quad \mathbf{M} = \mathbf{0} \qquad \mathbf{M} \vec{\mathbf{x}} = \vec{b} \qquad \mathbf{0} \vec{\mathbf{x}} = \vec{b} \qquad \vec{r}^k = \vec{b} \\ & \varphi_k \left(0 \right) = \beta_0 \mathbf{I} + \beta_k \mathbf{M} \mathbf{I} + \beta_k \mathbf{M}^2 + \dots + \beta_k \mathbf{M}^k = \beta_0 \mathbf{I} \\ & \vec{b} = \vec{r}^k = \varphi_k \left(0 \right) \vec{r}^0 \qquad \Rightarrow \varphi_k \left(0 \right) = \mathbf{I} \end{split}$$

Residual Minimization

If $\vec{x}^{k} \in span\{\vec{r}^{0}, \mathbf{M}\vec{r}^{0}, ..., \mathbf{M}^{k-1}\vec{r}^{0}\}$ minimizing $\left|\vec{r}^{k}\right|_{2}^{2}$

 $1.) \quad \vec{x}^k = \xi_{k-1}(\mathbf{M})\vec{r}^0$

 $\xi_{k-1}(\mathbf{M})$ is the k^{th} order poly minimizing $\left|\vec{r}^{k}\right|_{2}^{2}$ 2.) $\vec{r}^{k} = \varphi_{k}(\mathbf{M})\vec{r}^{0}$

 $= \vec{b} - \mathbf{M}\vec{x}^{k} = (\mathbf{I} - \mathbf{M}\boldsymbol{\xi}_{k}(\mathbf{M}))\vec{r}^{0}$

where $\varphi_k(\mathbf{M})r^0$ is the k^{th} order poly minimizing $\left|\vec{r}^k\right|_2^2$ subject to $\varphi_k(0) = \mathbf{I}$ Polynomial Property only a function of solution space and residual minimization

Explicitly construct Krylov subspace: $\vec{r}^0 = \vec{b} - \mathbf{M} \frac{\mathbf{g}^0}{\mathbf{0}}$ $\mathbf{M} \vec{r}^0 = \mathbf{M} \cdot \vec{r}^0$ $\mathbf{M}^2 \vec{r}^0 = \mathbf{M} \cdot \mathbf{M} \vec{r}^0$: $\mathbf{M} \vec{r}^k = \mathbf{M} \cdot \mathbf{M}^{k-1} \vec{r}^0$ Orthogonalize & Project Residual Minimizing Optimality Property:

 $\begin{aligned} \left| \overline{r}^{k} \right| &= \left| \varphi_{k} \left(\mathbf{M} \right) \overline{r}^{0} \right| \leq \left| \overline{\varphi}_{k} \left(\mathbf{M} \right) \overline{r}^{0} \right| \leq \left| \overline{\varphi}_{k} \left(\mathbf{M} \right) \right| \left| \overline{r}^{0} \right| \\ \overline{\varphi}_{k} \text{ is any } k^{\text{th}} \text{ order poly such that } \overline{\varphi}_{k} (0) = \mathbf{I} \\ \hline \begin{array}{c} \hline \text{Interefor} \\ \text{Any polynomial which satisfies the constraint can be used to get} \\ \text{an upper bound on} \left| \overline{r^{k}} \right| \leq \left| \overline{\varphi}_{k} \left(\mathbf{M} \right) \right| \end{aligned}$





Spectral Mapping Theorem. Given a polynomial $\widetilde{\wp}_p(x) = a_0 + a_1x + ... + a_px^p$ Apply the polynomial to a matrix $\widetilde{\wp}_p(\mathbf{M}) = a_0I + a_1\mathbf{M} + ... + a_p\mathbf{M}^p$ Then *spectrum* $(\widetilde{\wp}_p(\mathbf{M})) = \widetilde{\wp}_p(spectrum(\mathbf{M}))$ Proof. Note a property of matrix powers $\mathbf{M}\mathbf{M} = U\lambda U^{-1}U\lambda U^{-1} = U\lambda^2 U^{-1} \implies \mathbf{M}^p = U\lambda^p U^{-1}$ Apply to the polynomial of the matrix

Apply to the polynomial of the matrix $\widetilde{\wp}_{p}(\mathbf{M}) = a_{0}\mathbf{U}\mathbf{U}^{-1} + a_{1}\mathbf{U}\lambda\mathbf{U}^{-1} + ... + a_{p}\mathbf{U}\lambda^{p}\mathbf{U}^{-1}$ Factoring $\widetilde{\wp}_{p}(\mathbf{M}) = \mathbf{U}\underbrace{\left(a_{0}I + a_{1}\lambda + ... + a_{p}\lambda^{p}\right)}_{diagonal}\mathbf{U}^{-1}$

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,



$$\begin{split} & \| \mathbf{f} \mathbf{M} = \mathbf{M}^{\mathsf{T}} \text{ then} \\ & 1. \quad \mathbf{M} \text{ has orthonormal eigenvectors} \\ & \quad cond(\mathbf{U}) = \begin{bmatrix} \uparrow & & \uparrow \\ \vec{u}_1 & \cdots & \vec{u}_N \\ \downarrow & & \downarrow \end{bmatrix} \begin{bmatrix} \uparrow & & \uparrow \\ \vec{u}_1 & \cdots & \vec{u}_N \\ \downarrow & & \downarrow \end{bmatrix}^{-1} = 1 \\ & \implies \| \widetilde{\wp}_k(\mathbf{M}) \| = \max_i \widetilde{\wp}_k(\lambda_i) \\ & 2. \quad \mathbf{M} \text{ has real eigenvalues} \\ & \text{ If } \mathbf{M} \text{ is positive definite, then } \lambda(\mathbf{M}) > 0 \end{split}$$

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Bound on quantity of residual reduction

$$\begin{vmatrix} \vec{r}^{k+1} \\ < \varepsilon & \longrightarrow & \mathbf{M}\vec{x}^{k+1} \cong \vec{b} \\ \\ \parallel \vec{x}^{k+1} & \longleftarrow & \vec{x}^* = \mathbf{M}^{-1}\vec{b} \\ \\ \vec{b} - \mathbf{M}\vec{x}^{k+1} = 0 \\ \Rightarrow \left| \vec{x}^{k+1} - \vec{x}^* \right| = \left| \mathbf{M}^{-1}\vec{r}^{k+1} \right| \le \left| \mathbf{M}^{-1} \right| \left| \vec{r}^{k+1} \right| = \left| \mathbf{M}^{-1} \right| \varepsilon$$

Polynomial Error Bound

r^{k}	=	$\wp_k(\mathbf{M})$	$\wp_k(\mathbf{M})$ is the GCR polynomial of order k
	\leq	$ \widetilde{\wp}_k(\mathbf{M}) $	$\widetilde{\varphi}_k(\mathbf{M})$ is any k^{th} order polynomial such that $\widetilde{\varphi}_k(0) = \mathbf{I}$
	VI VI	$cond(\mathbf{U}) \cdot \max_{i} \widetilde{\wp}_{k}(\lambda_{i}) $ $\max_{i} \widetilde{\wp}_{k}(\lambda_{i}) $	U are the eigenvectors of M λ_i are the eigenvalues of M if M is symmetric then cond(U) = 1



 If M has only q distinct eigenvalues, the residual minimizing Krylov subspace algorithm converges in at most q steps Proof;









0.5

275

1.5



Polynomial Min-Max Problem [Convergence for $\mathbf{M} = \mathbf{M}^{T}$] Consider $\lambda(M) \in [\lambda_{\min}, \lambda_{\max}]$. $\lambda_{\min} > 0$ Then a good polynomial (i.e. $\|\vec{\wp}_{k}(\mathbf{M})\|$ is small) can be found by solving the min - max problem

$$\min_{\substack{k^{th} \text{ order } x \in [\lambda_{\min}, \lambda_{\max}]}} \max_{\substack{\beta \neq k \\ \beta \neq k}} \widetilde{\varphi}_k(x)$$

 $\overline{\delta_k(0)}=1$ The min-max problem is exactly solved by Chebyshev Polynomials













Heat Conducting Bar example

Sometimes GCR can do much better than Chebyshev bound





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INTRODUCTION TO NUMERICAL SIMULATION

LECTURE 10.

Preconditioners

TODAY'S OUTLINE:

* Preconditioners

- Diagonal Preconditioner
- Blockdiagonal Preconditioner
- Incomplete Factorization Preconditioner
- ✤ GCR for different Right Hand Sides
 - Recycling the Krylov subspace

PRECONDITIONERS

Diagonal Preconditioner

* Diagonal Example

1	0 0	0	[1	0	0		0		
)	1	0		0	0	2	0		0
)	0	٠.	٠.	:	0	0	÷.,	÷.,	:
	÷	۰.	1	0	1	÷	÷.,	N-1	0
)		0	0	1	0		0	0	N





3.10Linear Systems I. GCR Preconditioners





Diagonal Preconditioner Try as preconditioner the inverse of the diagonal [] * * * * * * * $\left| \frac{r^{1}}{r^{0}} \right| \approx \text{small}$ **1******







 $(\mathbf{D}^{-1}\mathbf{A})\vec{\mathbf{x}} = (l + \mathbf{D}^{-1}\mathbf{A}_{nd})\vec{\mathbf{x}} = \mathbf{D}^{-1}\vec{b}$

The inverse of a diagonal is cheap to compute

Usually improves convergence



 Δv $\overline{100}$ One small Δv



 $\frac{\lambda_{max}}{\lambda_{min}} > 100$

For the heat conducting bar, which convergence curve is GCR?





-1.5

-20

0.5

1.5

them.

2

5



Blockdiagonal Preconditioner Line Schemes Grid Matrix 6 6 6 6 . do 0 6.6.6 . 4 * * * -6 6 . 0 0 636 8 di. ÷ Φ 4 0 19-. . 45 3 - 183

8

Tridiagonal Matrices factor quickly











Key Idea

Pick M such that :

a.) Pick $\widetilde{\mathbf{M}} \approx \mathbf{M}$

b.) $\widetilde{\mathbf{M}}$ is easy to factor : $\widetilde{\mathbf{M}} = \widetilde{\mathbf{L}}\widetilde{\mathbf{U}}$

Use as preconditioner : $\mathbf{P} = \widetilde{\mathbf{M}}^{-1}$

 $\mathbf{M}\vec{x} = \vec{b}$

 $\widetilde{\mathbf{M}}^{-1}\mathbf{M}\vec{x} = \widetilde{\mathbf{M}}^{-1}\vec{b}$

But do NOT calculate $\mathbf{P} = \widetilde{\mathbf{M}}^{-1}$

 $\widetilde{\mathbf{M}}^{-1}\mathbf{M}\vec{x} = \underline{\widetilde{\mathbf{M}}}^{-1}\underline{\vec{b}}$ $\widetilde{\mathbf{M}} = \widetilde{\mathbf{L}}\widetilde{\mathbf{U}}$ $\widetilde{b} = \widetilde{\mathbf{M}}^{-1} \widetilde{b}$ solve for \tilde{b} : $\tilde{L}\tilde{U}\tilde{b} = \vec{b}$ At each iteration of GCR we need : $\widetilde{\mathbf{M}}^{-1}\mathbf{M}r^k = \hat{r}^k$ Then solve for \hat{r}^k : $\widetilde{\mathbf{L}}\widetilde{\mathbf{U}}\hat{r}^k = \mathbf{M}r^k$







★ Complete Algorithm Pick M such that : a.) Pick M = M b.) M is easy to factor : M = LU Factor M = LU Calculate new RHS solving for b : LUb = b At the kth step of GCR: Calculate Mr^k solve for p^k : LUp^k = Mr^k

GCR Summary

- Comparison CGR & GE
 Making GCR converge fast can be very problem specific: preconditioning
 - Assume we have a good pre-conditioner (e.g. # iterations k < 10,20)
 Sparse Dense

$$n = 10$$

$$Re(\lambda)$$
Same condition number for both.

Top converges faster - clustered eigenvalues.

- When would one use GCR?
 - Need to have fast convergence rate:
 - 1. cond (M) is not large
 - or eigenvalues of M are clustered in few groups
 - or vegen undes of the uncered in rew groups
 or we have a good "pre-conditioner"
 - And one of the following:
 - 1. need fewer than 16 digits of precision
 - 2. or M is dense
 - 3. or have a fast matrix-vector product algorithm

And what about if you need to re-solve a system with different RHS?



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n = 10



INTRODUCTION TO NUMERICAL SIMULATION



 $\vec{f}_{l} = (0, -10)$

 $\frac{y_2 - y_0}{L_4} \varepsilon_A \left(L_0 - L_A \right) + \frac{y_2 - y_1}{L_B} \varepsilon_B \left(L_0 - L_B \right) + f_{l,y} = 0$


 $I_r - \frac{1}{10}V_r = 0$

 $I_{d} - I_{s}(e^{\frac{V_{d}}{V_{i}}} - 1) = 0$

Circuit Example

10V





Richardson And Linear CONVERGENCE Richardson Iteration Richardson Iteration Definition $x^{k+1} = x^k + f(x^k)$ An iteration stationary point is a solution

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Non-Linear

System of Equations

f(x) = 0

M

 $x^{k+1} = x^{k}$ $\Rightarrow f(x^{k}) = 0$ $\Rightarrow x^{k} = x^{*} \text{ (Solution)}$ ***** Example 1 f(x) = -0.7x + 10Start with $x^{0} = 0$ $x^{1} = x^{0} + f(x^{0}) = 0 + 10 = 10$ $x^{2} = x^{1} + f(x^{1}) = 10 + (-7 + 10) = 13$ $x^{3} = x^{2} + f(x^{2}) = 13 + (-0.7 \cdot 13 + 10) = 14.17$ $x^{5} = 14.25$ $x^{6} = 14.27$ $x^{7} = 14.28$ $x^{8} = 14.28 \longleftarrow \text{Converged!}$



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★ Example 2 f(x) = 2x + 10Start with $x^0 = 0$ $x^1 = x^0 + f(x^0) = 0 + 10 = 10$ $x^2 = x^1 + f(x^1) = 10 + (2 \cdot 10 + 10) = 40$ $x^3 = x^2 + f(x^2) = 40 + (2 \cdot 40 + 10) = 130$ $x^4 = x^3 + f(x^3) = 130 + (2 \cdot 130 + 10) = 400$... No Convergence!!



$$f(x) = x$$

$$f(x) = x$$

$$f(x) = x$$

$$f(x) = b - \mathbf{A}x$$
given
$$x^{k+1} = x^{k} + f(x) = x^{k} + b - \mathbf{A}x^{k}$$

$$x^{*} = x^{*} + f(x^{*}) = x^{*} + b - \mathbf{A}x^{*}$$

$$x^{k+1} - x^{*} = (\mathbf{I} - \mathbf{A})(x^{k} - x^{*})$$

$$x^{k} - x^{*} = (\mathbf{I} - \mathbf{A})(x^{0} - x^{*})$$

$$x^{1} - x^{*} = (\mathbf{I} - \mathbf{A})(x^{0} - x^{*})$$

$$x^{2} - x^{*} = (\mathbf{I} - \mathbf{A})(x^{1} - x^{*}) = (\mathbf{I} - \mathbf{A})^{2}(x^{0} - x^{*})$$

$$\lim_{\substack{k \to \infty \text{ spectral } < 1}} |(\mathbf{I} - \mathbf{A}(\mathbf{A}))| < 1$$

★ Convergence
$$x^{k+1} = x^{k} + f(x^{k})$$

$$x^{*} = x^{*} + f(x^{*})$$

$$f(x) = f(x^{*}) + \frac{df}{dx}(x^{*})(x - x^{*}) + \frac{1}{2}\frac{d^{2}f}{dx^{2}}(x^{*})(x - x^{*})^{2} + \dots$$

$$x^{k+1} - x^{*} = x^{k} - x^{*} + f(x^{*}) + \frac{df}{dx}(x^{*})(x^{k} - x^{*}) + \frac{1}{2}\frac{d^{2}f}{dx^{2}}(x^{*})(x^{k} - x^{*})^{2} + \dots$$

$$|x^{k+1} - x^{*}| = \left|1 + \frac{df}{dx}(x^{*}) + \frac{1}{2}\frac{d^{2}f}{dx^{2}}(x^{*})(x^{k} - x^{*}) + \dots\right| x^{k} - x^{*}|$$

$$|x^{k+1} - x^{*}| = \left|1 + \frac{df}{dx}(x^{*}) + \frac{1}{2}\frac{d^{2}f}{dx^{2}}(x^{*})(x^{k} - x^{*}) + \dots\right| x^{k} - x^{*}|$$
Find a bound: $K^{k} \leq K^{*} = x^{k} - x^{*}|$

$$|x^{k+1} - x^{*}| = K^{*}|x^{k} - x^{*}| \quad \text{for large } k$$
If $K^{*} < 1$ then Richardson converges linearly.
$$|x^{k+1} - x^{*}| \leq \frac{1}{2}|x^{k} - x^{*}| \quad |x^{k} - x^{*}| = 10$$

$$|x^{k+1} - x^{*}| \leq \frac{1}{2}|x^{k} - x^{*}| \quad |x^{k+1} - x^{*}| \leq \frac{1}{2} \cdot 10 = 5$$

$$|x^{k+10} - x^{*}| \leq \frac{(1)^{10} \cdot 10}{x^{*} - x^{*} + f(x^{k})}$$
Exact Solution $x^{*+1} = x^{k} + f(x^{k})$
Exact Solution $x^{*+1} = x^{k} + f(x^{k})$
Computing Differences
$$x^{k+1} - x^{*} = x^{k} - x^{*} + \frac{f(x^{k}) - f(x^{*})}{Neet to Estimate}$$

$$\begin{aligned} x^{k+1} - x^* &= g\left(x^k - x^*\right) \\ x^{k+1} - x^* &= \gamma \left[x^k - x^*\right] \\ \gamma \leq 1 \end{aligned}$$

$$\begin{aligned} \mathcal{K}^* &= \left|1 + \frac{\partial f}{\partial x} \left(x^*\right) + \frac{1}{2} \frac{\partial^2 f}{\partial x^2} \left(x - x^*\right) + \varepsilon\right| \\ \text{Example 1.} \quad f(x) = -0.7x + 10 \\ \mathcal{K}^* &= \left|1 + (-0.7) + 0 + \varepsilon\right| = \left|0.3 + \varepsilon\right| \\ \text{Example 2.} \quad f(x) = 2x + 10 \\ \mathcal{K}^* &= 1 + 2 + 0 + \varepsilon\right| = \left|3 + \varepsilon\right| \\ \text{No convergence} \end{aligned}$$

$$\circ \quad \text{Mean Value Theorem} \\ f(w) - f(v) &= \frac{\delta f(z)}{\delta x} (w - v) \qquad z \in [v, w] \end{aligned}$$

Z

w

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v

• Use Mean Value Theorem
Iteration Equation
$$x^{k+1} = x^k + f(x^k)$$

Exact Solution $x^* = x^* + f(x^*)$
Computing Differences
 $x^{k+1} - x^* = x^k - x^* + f(x^k) - f(x^*) = \left(1 + \frac{\partial f(\tilde{x})}{\partial x}\right)(x^k - x^*)$
 $\overline{x} \in \left|x^k, x^*\right|$
 $|x^{k+1} - x^*| = \left|x^k - x^* + \frac{\partial f}{\partial x}(\tilde{x})(x^k - x^*)\right|$
 $= \left|1 + \frac{\partial f}{\partial x}(\tilde{x})\right|(x^k - x^*)$
 $\leq \gamma |(x^k - x^*)|$
 $\leq \gamma |(x^{k-1} - x^*)|$
 $\leq \gamma r|(x^{k-1} - x^*)|$

• Richardson Theorem
IF
$$\left|1 + \frac{\partial f(\tilde{x})}{\partial x}\right| \le \gamma < 1$$
 for all \tilde{x} s.t. $\left|\tilde{x} - x^*\right| < \delta$
AND $\left|x^0 - x^*\right| < \delta$
THEN $\left|x^{k+1} - x^*\right| \le \gamma \left|x^k - x^*\right|$
OR $\lim_{k \to \infty} \left|x^{k+1} - x^*\right| = \lim_{k \to \infty} x^0 - x^*\right| = 0$
 $\left|1 + \frac{\partial f(\tilde{x})}{\partial x}\right| < 0.1$ $\frac{\partial f}{\partial x} = \frac{x^0 - x}{|x^1 - x||^2}$



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NEWTON'S METHOD Derivation of the Basic Algorithm * 1-D Reminder • Newton Idea Problem : Find x^* such that $f(x^*) = 0$ Approximate f(x) with its Taylor series about x^k : $\hat{f}(x) \cong f(x^k) + \frac{\delta f}{\delta x}(x^k)(x - x^k)$ a straight line Find the solution of the approximation x^{k+1} : $f(x^k) + \frac{\delta f}{\delta x}(x^k)(x^{k+1} - x^k) = 0$ $\Rightarrow x^{k+1} = x^k - \left[\frac{\partial f}{\partial x}(x^k)\right]^{-1} f(x^k)$









• Newton Algorithm $x^0 =$ Initial Guess, k = 0Repeat {

$$\begin{aligned} x^{k+1} &= x^k - \left[\frac{\partial f}{\partial x}(x^k)\right]^{-1} f(x^k) \\ k &= k+1 \\ \} \text{ Until ?} \\ \|x^{k+1} - x^k\| < \text{threshold} \\ \|f(x^{k+1})\| < \text{threshold} \end{aligned}$$

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• Convergence C $\circ f(x)$

Need an "f(x)" check to avoid false convergence.



 $\circ \Delta x$

Need a " Δx " check to avoid false convergence.







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 Taylor Expansion $\frac{\partial f}{\partial x}$ $\frac{\partial f}{\partial x} = |f(x_B) - f(x_A)|$ $\frac{\partial f}{\partial x} = |f(x_B) - f(x_A)|$ XA

Taylor Series Expansion

$$f(x) = f\left(x^{k}\right) + \frac{\partial}{\partial x} f\left(x^{k}\right) \left(x - x^{k}\right) + \frac{1}{2!} \frac{\partial^{2}}{\partial x^{2}} f\left(x^{k}\right) \left(x - x^{k}\right)^{p+1} + \cdots + \frac{1}{n!} \frac{\partial^{n}}{\partial x^{n}} f\left(x^{k}\right) \left(x - x^{k}\right)^{p} + O\left(\left(x - x^{k}\right)^{p+1}\right)$$

$$f(x) = f\left(x^{k}\right) + \frac{\partial}{\partial x} f\left(x^{k}\right) \left(x - x^{k}\right) + \frac{1}{2!} \frac{\partial^{2}}{\partial x^{2}} f\left(x^{k}\right) \left(x - x^{k}\right)^{p+1} + \cdots + \frac{1}{n!} \frac{\partial^{n}}{\partial x^{n}} f\left(x^{k}\right) \left(x - x^{k}\right)^{p} + \frac{1}{(n+1)!} \frac{\partial^{n+1}}{\partial x^{n+1}} f\left(\overline{x}\right) \left(x - x^{k}\right)^{p+1} + \overline{x} \in \left[x, x^{k}\right]$$

Exact Taylor expansion about the iteration x^k evaluated at the solution x^*

$$0 = f\left(x^*\right) = f\left(x^k\right) + \frac{df}{dx}\left(x^k\left)\left(x^* - x^k\right) + \frac{1}{2}\frac{d^2f}{dx^2}\left(\widetilde{x}\right)\left(x^* - x^k\right)^2\right)$$
(1)
some $\underbrace{\widetilde{x} \in \left[x^k, x^*\right]}_{\text{Mean Value Theorem}}$

Approximate Taylor expansion about the solution x^4 evaluated at the next iteration x^{k+1} $(\cdot) = \frac{df}{dt} (\cdot) (\cdot \cdot \cdot \cdot)$

$$0 = f(x^{k}) + \frac{df}{dx} (x^{k}) (x^{k+1} - x^{k})$$
(2)
Subtracting (1) from (2)
$$\frac{df}{dx} (x^{k}) (x^{k+1} - x^{*}) = \frac{d^{2}f}{dx^{2}} (\widetilde{x}) (x^{*} - x^{k})^{2}$$
Dividing through
$$(x^{k+1} - x^{*}) = \left[\frac{df}{dx} (x^{k}) \right]^{-1} \frac{d^{2}f}{dx^{2}} (\widetilde{x}) (x^{*} - x^{k})^{2}$$









Proof. $\begin{vmatrix} x_{1} - x^{*} &| \leq L | (x_{0} - x^{*}) | x_{0} - x^{*} | \\ \Rightarrow &| x_{1} - x^{*} &| \leq L \gamma | x_{0} - x^{*} | \\ \Rightarrow &| x_{2} - x^{*} &| \leq L \gamma | x_{0} - x^{*} | \\ arrow &| x_{2} - x^{*} &| \leq L \gamma | x_{0} - x^{*} | \\ \Rightarrow &| x_{3} - x^{*} &| \leq Y^{2} | x_{1} - x^{*} &| \leq Y^{3} | x_{0} - x^{*} | \\ \Rightarrow &| x_{3} - x^{*} &| \leq Y^{4} | x_{2} - x^{*} &| \leq Y^{7} | x_{0} - x^{*} | \\ \hline &| x^{k+1} - x^{*} &| = \left[\frac{\partial f}{\partial x} (x^{k}) \right]^{-1} \frac{\partial^{2} f}{\partial x^{2}} (\bar{x}) | x^{k} - x^{*} |^{2} \\ k = 0: &| x^{1} - x^{*} &| \leq L | x^{0} - x^{*} | \\ k = 1: &| x^{2} - x^{*} &| \leq L | x^{0} - x^{*} | \\ &\leq L \gamma | x^{0} - x^{*} | | x^{0} - x^{*} | \\ \leq L \gamma | x^{0} - x^{*} | \gamma | x^{0} - x^{*} | \\ \leq Y^{2} | x^{0} - x^{*} | \\ k : &| x^{k+1} - x^{*} &| \leq \gamma^{2(k+1)-1} | x^{0} - x^{*} | \\ \vdots \\ k : &| x^{k+1} - x^{*} &| \leq \gamma^{2(k+1)-1} | x^{0} - x^{*} | \\ \gamma = \left[\frac{\partial f}{\partial x} (x) \right]^{-1} \left[\frac{\partial^{2} f}{\partial x^{2}} (x) \right] x^{0} - x^{*} | < 1 \\ \checkmark Theorem$ If L is bounded $\left[\frac{df}{dx}$ bounded away from zero; $\frac{d^{2} f}{dx^{2}}$ bounded $\right]$ then Newton's method is guaranteed to converge given a "close enough" guess Always converges?

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.













 Jacobian Matrix $J_F(x)\Delta x \approx F(x + \Delta x) - F(x)$ $\begin{vmatrix} \frac{\partial F_1(x)}{\partial x_1} & \cdots & \frac{\partial F_1(x)}{\partial x_N} \\ \vdots & \ddots & \vdots \\ \frac{\partial F_N(x)}{\partial x_N} & \cdots & \frac{\partial F_N(x)}{\partial x_N} \end{vmatrix}$ Δx_1 $J_F(x)\Delta x =$ Δx_N General Setting General setting Problem : Find x^* such that $F(x^*) = 0$, $x^* \in \Re^N$ and $F: \Re^N \to \Re^N$ Approximate F(x) with its Taylor Series about x^k : $\hat{F}(x) \cong F(x^k) + J_F(x^k) (x^{k+1} - x^k)$ Jacobian Matrix Find the solution of the approximation: x^{k+1} $F(x^{k}) + J_{F}(x^{k})(x - x^{k}) = 0 \Longrightarrow x^{k+1} = x^{k} - [J_{F}(x^{k})]^{-1}F(x^{k})$ * Newton Algorithm $x^0 =$ initial guess, k = 0Repeat { Compute $F(x^k)$, $J_F(x^k)$ Solve $J_F(x^k)(x^{k+1} - x^k) = -F(x^k)$ for x^{k+1} k = k + 1} Until $x^{k+1} - x^k$, $f(x^{k+1})$ small enough Nodal Analysis o Strut and Joint $x^* \in \mathbb{R}^2 \text{ and } F: \mathbb{R}^2 \to \mathbb{R}^2$ $F_L \quad J_F(\vec{x}) = \begin{bmatrix} ? & ? \\ ? & ? \end{bmatrix}$ $\frac{x}{l} \varepsilon(l_0 - l) + F_{L_x} = 0$ $\frac{y}{l} \varepsilon(l_0 - l) + F_{L_y} = 0 \quad J_F(\vec{x}) = \begin{bmatrix} * & 0 \\ 0 & 0 \end{bmatrix}$ Finally between the set of the set o Singular Jacobian

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$$\begin{split} J_{F}(\vec{x}) &= \begin{bmatrix} * & 0 \\ 0 & \textcircled{0} \end{bmatrix} \\ &= \begin{bmatrix} \frac{1}{\partial y} \begin{bmatrix} \frac{y}{l} \varepsilon(l_{0} - l) + F_{L_{y}} \end{bmatrix}_{l_{0} = \sqrt{x_{0}^{2} + y_{0}^{2}}} \\ &= \begin{bmatrix} \frac{1}{l} \varepsilon(l_{0} - l) + \frac{y}{l_{0}} \varepsilon(\frac{\partial l}{\partial y}) + \frac{y}{l_{0}} \varepsilon(\frac{\partial l}{\partial y}) + \frac{y}{l_{0}} \varepsilon(\frac{\partial l}{\partial y}) \\ &= \begin{bmatrix} \frac{1}{l} \varepsilon(l_{0} - l) + \frac{y}{l_{0}} \varepsilon(\frac{\partial l}{\partial y}) + \frac{y}{l_{0}} \varepsilon(\frac{\partial l}{\partial y}) + \frac{y}{l_{0}} \varepsilon(\frac{\partial l}{\partial y}) \\ &= \int F_{1}\left(\frac{x + \Delta x}{y + \Delta y}\right) = F_{1}\left(\frac{x}{y}\right) + \frac{\partial F_{1}}{\partial x} \Delta x + \frac{\partial F_{1}}{\partial y} \Delta y \\ &= F_{2}\left(\frac{x + \Delta x}{y + \Delta y}\right) = F_{2}\left(\frac{x}{y}\right) + \frac{\partial F_{2}}{\partial x} \Delta x + \frac{\partial F_{2}}{\partial y} \Delta y \\ &\Rightarrow \vec{F}(\vec{x} + \Delta \vec{x}) \approx \vec{F}(\vec{x}) + \begin{bmatrix} \frac{\partial F_{1}}{\partial x} & \frac{\partial F_{1}}{\partial y} \\ \frac{\partial F_{1}}{\partial x} & \frac{\partial F_{2}}{\partial y} \end{bmatrix} \begin{bmatrix} \Delta x \\ \Delta y \end{bmatrix} \end{split}$$

o Nonlinear Resistor



Computing the Jacobian and the Function

Consider the contribution of one nonlinear resistor Connected between nodes n_1 and n_2

Summing currents at Node n_1 : $F_{n_1}(v) = g(v_{n_1} - v_{n_2}) + \dots$ Summing currents at Node n_2 : $F_{n_2}(v) = -g(v_{n_1} - v_{n_2}) + \dots$

Differentiating at Node n_1 :

$$\frac{\partial F_{n_1}(v)}{\partial v_{n_1}} = \underbrace{\frac{\partial g(v_{n_1} - v_{n_2})}{\partial v_{n_1}}}_{\frac{\partial g}{\partial x}} + \dots \quad \frac{\partial F_{n_1}(v)}{\partial v_{n_2}} = \underbrace{\frac{\partial g(v_{n_1} - v_{n_2})}{\partial v_{n_2}}}_{-\frac{\partial g}{\partial x}} + \dots$$





INTRODUCTION TO NUMERICAL SIMULATION

LECTURE 12.

Multi-Dimensional Newton Method

TODAY'S OUTLINE:

- Multi-Dimensional Newton Method
 - Convergence Analysis
- Damped Newton Schemes
 - Global Convergence
 - Difficulty with Singular Jacobians
- Continuation Schemes (Homotopy Methods)
 - Source/Load Stepping
 - Improving Continuation Efficiency

MULTI-DIMENSIONAL NEWTON METHOD **Convergence Analysis**

* Theorem Statement

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If a.) $J_F^{-1}(x^k) \leq \beta$

(Inverse is bounded)

b.) $|J_F(x) - J_F(y)| \le \ell ||x - y||$ (Derivative is Lipschitz Cont) Then Newton's method converges given a sufficiently close initial guess.





Limiting the changes in X might improve convergence



Global Convergence Newton Algorithm for Solving F(x) = 0 $\vec{x}^0 = \text{initial guess}, k = 0$ Repeat {

Compute $F\left(\vec{\mathbf{x}}^{k}\right) J_{F}\left(\vec{\mathbf{x}}^{k}\right)$ Solve $J_{F}\left(\vec{\mathbf{x}}^{k}\right) \Delta \vec{\mathbf{x}}^{k+1} = -F\left(\vec{\mathbf{x}}^{k}\right)$ for $\Delta \vec{\mathbf{x}}^{k+1}$ $\vec{\mathbf{x}}^{k-k+1} = \vec{\mathbf{x}}^{k} + \text{limited}\left(\Delta \vec{\mathbf{x}}^{k+1}\right)$ k = k+1





* Damped Newton Scheme with Limiting



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General Damping Scheme Solve $J_F(x^k)\Delta x^{k+1} = -F(x^k)$ for Δx^{k+1} $x^{k+1} = x^k + \alpha^k \Delta x^{k+1}$ Key Idea: Line Search







o Example



Nodal Equations with Numerical Values



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Every Step reduces F-Global Convergence!

Example. $\begin{aligned} \left| F\left(\frac{x^{k+1}}{x^k + \alpha^k \Delta x^{k+1}} \right)^2 &\leq 0.9 \left| F\left(x^k \right) \right|_2^2 \\ &\left| F\left(x^{k+1} + \alpha^{k+1} \Delta x^{k+2} \right)_2^2 &\leq 0.9 \left| F\left(x^{k+1} \right)_2^2 &\leq 0.9^2 \left| F\left(x^k \right)_2^2 \\ &\left| F\left(x^{k+100} + \alpha^{k+100} \Delta x^{k+101} \right)_2^2 &\leq 0.9 \left| F\left(x^{k+100} \right)_2^2 &\leq \underbrace{0.9^{101}}_{\to 0} \left| F\left(x^k \right)_2^2 \to 0 \end{aligned} \end{aligned}$

.

Proof.
By definition of the Newton Iteration
$$x^{k+1} = x^{k} - \alpha^{k} \frac{J_{F}(x^{k})^{-1} F(x^{k})}{N_{\text{conton Direction}}}$$
Multidimensional Mean Value Lemma
$$\|F(x) - F(y) - J_{F}(y)(x - y)\| \leq \frac{\ell}{2} \|x - y\|^{2}$$
Combining
$$\|F(x^{k+1}) - F(x^{k}) - J_{F}(x^{k}) \Big[\alpha^{k} J_{F}(x^{k})^{-1} F(x^{k}) \Big] \leq \frac{\ell}{2} \Big[\alpha^{k} J_{F}(x^{k})^{-1} F(x^{k}) \Big]^{2}$$
Combining terms and moving scalars out of norms
$$\|F(x^{k+1}) - (1 - \alpha^{k})F(x^{k})\| \leq (\alpha^{k})^{2} \frac{\ell}{2} \Big] J_{F}(x^{k})^{-1} F(x^{k}) \Big]^{2}$$
Using the Jacobian Bound and splitting the norm
$$\|F(x^{k+1})\| \leq \Big[(1 - \alpha^{k}) F(x^{k}) + (\alpha^{k})^{2} \frac{\beta^{2} \ell}{2} \|F(x^{k})\|^{2} \Big]$$
Yields a quadratic in the damping coefficient.
Simplifying quadratic



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Two Cases:
1.)
$$\left\|F\left(x^{k}\right)\right\| \leq \frac{1}{\beta^{2}t} \operatorname{pick} \alpha^{k} = 1$$
 (Standard Newton)
 $\left\|F\left(x^{k+1}\right)\right\| \leq \left[1-1+1^{2}\frac{\beta^{2}t}{2}\left\|F\left(x^{k}\right)\right]\right] \left\|F\left(x^{k}\right)\right\| \leq \frac{1}{2}\left\|F\left(x^{k}\right)\right\|$
Case 1. $\frac{1}{2C} \geq 1$
 $\frac{1}{2\left[\frac{\beta^{2}t}{2}\left\|F\left(x^{k}\right)\right]\right]} \geq 1$
 $\Rightarrow \left\|F\left(x^{k}\right)\right\| \leq \frac{1}{\beta^{2}t}$ Take the full Newton step.
pick $\alpha^{k} = 1$ (Standard Newton)
2.) $\left\|F\left(x^{k+1}\right)\right\| > \frac{1}{\beta^{2}t}$ pick $\alpha^{k} = \frac{1}{\beta^{2}t}\left\|F\left(x^{k}\right)\right\|$ (Damped Newton)
 $\left\|F\left(x^{k+1}\right)\right\| \leq \left[1 - \frac{1}{\beta^{2}t}\left|F\left(x^{k}\right)\right|\right]\right\|F\left(x^{k}\right) < \left\|F\left(x^{k}\right)\right\|$

Case 2.
$$\frac{1}{2C} < 1$$

 $\beta^{2I}_{2} |F(x^{k})| > \frac{1}{2} \Rightarrow |F(x^{k})| > \frac{1}{\beta^{2}I} + \frac{1}{0} + \frac{1}{2C} + \frac$

If
$$\beta$$
 small ℓ small $\rightarrow \frac{1}{\beta^2 \ell}$ large Always take the
 $1D \int D \int \frac{1}{\beta^2 \ell}$ large $\frac{\partial^2 f}{\partial x^2}$ small
If β large ℓ large $\rightarrow \frac{1}{\beta^2 \ell}$ small
 $1D \int D \int \frac{1}{\beta^2 \ell}$ small
 $\frac{\partial f}{\partial x}$ small $\frac{\partial^2 f}{\partial x^2}$ large

$$\begin{split} & \left|F\left(x^{k+1}\right)\right| \leq \gamma^{k} \left\|F\left(x^{k}\right)\right\| \quad \text{not good enough, need } \gamma \text{ independent from } k \\ & \text{The above result does imply} \\ & \left|F\left(x^{k+1}\right)\right| \leq \left\|F\left(x^{0}\right)\right\| \quad \text{not yet a convergence theorem} \\ & \Rightarrow 1 - \frac{1}{2\beta^{2}\ell} \left\|F\left(x^{k}\right)\right\| \leq 1 - \frac{1}{2\beta^{2}\ell} \left\|F\left(x^{0}\right)\right| \leq \gamma^{0} \\ & \left\|F\left(x^{k+1}\right) \leq 0.9\right| F\left(x^{k}\right) \\ & \left\|F\left(x^{k+2}\right) \leq 0.99\right\| F\left(x^{k+1}\right) \\ & \left\|F\left(x^{k+3}\right) \leq 0.999\right\| F\left(x^{k+2}\right) \\ & \lim_{k \to \infty} \left\|F\left(x^{k}\right)\right\| \to 0? \end{split}$$



CONTINUATION SCHEMES (HOMOTOPY METHODS)







o Heat-conducting bar example









• Template Algorithm Solve $\tilde{F}(x(\lambda), \lambda) = 0$, $x(\lambda, prev) = x(0)$ $\delta \lambda = 0.01$, $\lambda = \delta \lambda$ While $\lambda < 1$ { $\begin{aligned} x^{0}(\lambda) &= x(\lambda_{prev}) \\ \text{Try to solve } \tilde{F}(x(\lambda),\lambda) &= 0 \text{ with Newton} \\ \text{If Newton Converged} \\ &x(\lambda_{prev}) &= x(\lambda), \ \lambda &= \lambda + \delta\lambda, \ \delta\lambda &= 2\delta\lambda \end{aligned}$ Else $\delta \lambda = \frac{1}{2} \delta \lambda, \ \lambda = \lambda_{prev} + \delta \lambda$ }





Source/Load Stepping Does Not Alter Jacobian





* Jacobian Altering Scheme - Description $\widetilde{F}(x(\lambda),\lambda) = \lambda F(x(\lambda)) + (1-\lambda)x(\lambda)$ Observations $\lambda = 0 \quad \widetilde{F}(x(0),0) = x(0) = 0$ $\frac{\partial \widetilde{F}(x(0),0)}{\partial \widetilde{F}(x(0),0)} = \mathbf{1}$ Problem is easy to solve and $\frac{\partial f(x)}{\partial x} = \mathbf{1}$ $\lambda = 1 \quad \widetilde{F}(x(1), 1) = F(x(1))$ Jacobian is definitely nonsingular. $\partial \tilde{F}(x(1),1) = \partial F(x(1))$ Back to the original problem and ∂x *ĉx* original Jacobian

$$R$$

$$V$$

$$F(v) = \frac{v - v_s}{R} + i_d(v) = 0$$
Diode
$$\frac{R}{\lambda}$$

$$V$$

$$V_s$$

$$i_d$$

$$\frac{1}{1 - \lambda}$$

 $v - v_s$

R

slowly

= 0V $\lambda i_d(v)$ = 01 Remove non - linear $1-\lambda$ Remove Resistors components then Insert resistors to then scale them bring them back ground in each mode slowly $R_{\lambda} = \frac{1}{1 - \lambda}$ then remove them slowly







INTRODUCTION TO NUMERICAL SIMULATION

LECTURE 13.

Newton-Method Case Study - Simulating an Image Smother

TODAY'S OUTLINE:

- Image segmentation example description
 Formulation: node-branch or nodal?
- * What Solver for Newton step: sparse LU or GCR?
 - Convergence of Newton GCR
 - Matrix-free idea
- * Continuation?
 - Jacobian altering scheme (for singular Jacobians)

 - Jacobian altering with update improvement
 Arc-Length Continuation (for multiple solutions)

IMAGE EXAMPLE

Simple Smoother

* Circuit Diagram



Array of photo diodes or photo transistors that provide some voltage proportional to the light in that particular spot of the chip





Input Image
 Input Image
 Smoothed Result

- 3



4

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FORMULATION – NODE-BRANCH OR NODAL? The nodal matrix is smaller and diagonally dominant in the linear case. It is also very sparse: 5 non-zeros per row.

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- And do we REALLY need to assemble the Jacobian?

Matrix-free Idea

Consider Applying GCR to the Newton Iterate Equation $J_F(x^k) \Delta x^{k+1} = -F(x^k)$ At each iteration GCR forms a matrix-vector product $J_F(x^k) t' \approx \frac{1}{c} [F(x^k + cx^l) - F(x^k)]$

It is possible to use Newton-GCR without Jacobians!! Need to select a good a

k = k + 1} Until $\left\| \Delta x^{k+1} \right\|, \left\| F(x^{k+1}) \right\|$ small enough



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Convergence of Newton-GCR Basic Algorithm – Solve Accuracy Required After l steps of GCR $J_{F}\left(x^{k}\right)\underbrace{\Delta x^{k+1,l}}_{\substack{\text{Newton} \\ \text{delta from}}} = -F\left(x^{k}\right) + \underbrace{r^{k,l}}_{\substack{\text{GCR} \\ \text{Residual}}}$ Newton delta from / GCRsteps If $a.) \quad \left|J_F^{-1}(x^k)\right| \leq \beta$ (Inverse is bounded) b.) $|J_F(x) - J_F(y)| \le \ell |x - y|$ (Derivative is Lipschitz Cont.) $c.) \quad \left| r^{k,l} \right| \le C \left| F\left(x^k \right)^2 \right|$ (More accurate near convergence) Then The Newton-Iterative Method Converges Quadratically Count the number of iterations $\Rightarrow F(x^k)$ $F(x^k + \varepsilon p^l)$ Number of digits - doubles with every Newton step (quadratic convergence) 16 Number of digits - linear convergence 3 steps GCR - 1 digit = 0.1Accuracy: 8 digits $4 \times (3 \times 8) \times 2 = 192$ Quadratic Newton GCR eval $8 \times 3 \times 2 = 48$ Linear Newton GCR eval



THEOREM	PRACTICALLY
Use as many GCR as needed for quadratic convergence i.e. $ r' \leq C F(x^{4}) ^{2}$ so that only need a few Newton steps	Stop GCR when residual is down by one order of magnitude (might need more Newton steps) What we really want to count is not the number of Newton steps but rather the total number of system evaluations (2 per each GCR)
Newton Quadratic 1 2 4 8 16 32 digits GCR 3 steps per digit 2 evaluations per GCR step	Newton Linear 1 2 3 4 5 6 digits GCR 3 steps per digit 2 evaluations per GCR step
e.g. want 8 digits: $\begin{split} 4_{Newton steps} &\times (3 \times 8)_{GCR steps} \times 2_{eval} \\ &= 192 \text{ evaluations} \\ \text{e.g. want 16 digits:} \\ 5_{Newton steps} &\times (3 \times 16)_{GCR steps} \times 2_{eval} \end{split}$	e.g. want 8 digits: $8_{Newton steps} \times 3_{GCR steps} \times 2_{eval}$ = 48 evaluations e.g. want 16 digits: $16_{Newton steps} \times 3_{GCR steps} \times 2_{eval}$

Proof:
By definition of the Newton-Iterative Method

$$x^{k+1} = x^{k} - \underbrace{J_{F}\left(x^{k}\right)^{-1}\left(F\left(x^{k}\right) + r^{k,l}\right)}_{Approximate Newton Direction}$$
Multidimensional Mean Value Lemma

$$\|F(x) - F(y) - J_{F}(y)(x - y)\| \leq \frac{\ell}{2} \|x - y\|^{2}$$
Combining

$$\left|r\left(x^{k+1}\right) - F\left(x^{k}\right) - J_{F}\left(x^{k}\left[J_{F}\left(x^{k}\right)^{-1}\left(F\left(x^{k}\right) + r^{k,l}\right)\right]\right] \leq \frac{\ell}{2} \left[\left[J_{F}\left(x^{k}\right)^{-1}\left(F\left(x^{k}\right) + r^{k,l}\right)\right]\right]^{2}$$

$$\left\|F\left(x^{k+1}\right) - F\left(x^{k}\right) - F\left(x^{k}\right) + r^{k,l}\right\| \leq \frac{\ell}{2} \left[\left[J_{F}\left(x^{k}\right)^{-1}\left(F\left(x^{k}\right) + r^{k,l}\right)\right]\right]^{2}$$
Using the triangle inequality

$$\left\|F\left(x^{k+1}\right) = \frac{\ell}{2} \left[J_{F}\left(x^{k}\right)^{-1}\left(F\left(x^{k}\right) + r^{k,l}\right)\right]\right\|^{2} + \left|r^{k,l}\right|$$
Using the Jacobian Bound







Source/Load Stepping Does Not Alter Jacobian






λ













Update Improvement



Solve for
$$\delta x$$
: $\left\lfloor \frac{\delta F}{\delta x} \right\rfloor \delta x = -\left[F(x(\lambda_{prev})) - x(\lambda_{prev})\right] \delta \lambda$

 \Rightarrow Initial guess for next Newton :

$$x^{0}(\lambda) = x(\lambda_{prev}) + \delta x$$

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INTRODUCTION TO NUMERICAL SIMULATION

LECTURE 14.

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Methods for Ordinary Differential Equations



Ē **ODE - Backward Euler, Forward Euler, Trape-**

ಲು

zoidal Rule

dt

2

dt





Signal Transmission in an Integrated Circuit ↔ A 2×2 Example

$v_1 \downarrow_{R2} \qquad v_2 \downarrow_{R2} \qquad o$	Constitutive Equations	Conservation Laws
$C_{1} = \langle R_{1} Q \rangle = \langle R_{2} Q \rangle = \langle R_{3} Q \rangle = \langle Q \rangle$	$i_C = C \frac{dv_C}{dt}$	$i_{C_1} + i_{R_1} + i_{R_2} = 0$
Nodal Equations Vields 2×2 System	$i_R = \frac{1}{R} v_R$	$i_{C_2} + i_{R_3} - i_{R_2} = 0$
$\begin{bmatrix} C_1 & 0 \end{bmatrix} \begin{bmatrix} \frac{dv_1}{dt} \end{bmatrix}$	$= -\left[\frac{1}{R_1} + \frac{1}{R_2}\right]$	$-\frac{1}{R_2}$ v_1
$\begin{bmatrix} 0 & C_2 \end{bmatrix} \begin{bmatrix} \frac{dv_2}{dt} \end{bmatrix}$	$-\frac{1}{R_2}$	$\frac{1}{R_3} + \frac{1}{R_2} \left[\begin{array}{c} v_2 \end{array} \right]$
Stamping Procedure for Nodal	Formulation	n in Circuits

$\frac{Capacitor}{q \propto v} \qquad q = 0$ Charge is	$Cv \longrightarrow differentiate$	$d_c = \frac{dq}{dt} = C\frac{dv_c}{dt}$
voltage Conservation Laws:		
$ \begin{aligned} &i_{C_1} + i_{R_1} + i_{R_2} = 0 \\ &i_{C_2} + i_{R_3} - i_{R_2} = 0 \end{aligned} $	$C_{1} = \sum_{R_{1} \in I_{R_{2}}} C_{R_{1}}$	$l \leq \frac{1}{2} C_{j}$

Use Nodal Analysis (substitute branch equations)



Stamping for Capacitors is like Stamping for Resistors



EIGENVALUE ANALYSIS FOR DYNAMICAL SYSTEMS Signal Transmission in an Integrated Circuit A 2×2 Example



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Consider an ODE : $\frac{dx(t)}{dt} = \mathbf{A}x(t), \quad x(0) = x_0$

Eigendecomposition: $\mathbf{A} = \begin{bmatrix} \uparrow & \uparrow & \uparrow \\ E_1 & E_2 & \cdots & E_n \\ \downarrow & \downarrow & \downarrow \end{bmatrix} \begin{bmatrix} \lambda_1 & & \\ \ddots & \\ \lambda_n \end{bmatrix} \begin{bmatrix} \uparrow & \uparrow & \uparrow \\ E_1 & E_2 & \cdots & E_n \\ \downarrow & \downarrow & \downarrow \end{bmatrix}^{-1}$ Change of variables: $\mathbf{E}y(t) = x(t) \Leftrightarrow y(t) = \mathbf{E}^{-1}x(t)$ Substituting : $\frac{d\mathbf{E}y(t)}{dt} = \mathbf{A}\mathbf{E}y(t), \ \mathbf{E}y(0) = x_0$ Multiply by \mathbf{E}^{-1} : $\frac{dy(t)}{dt} = \mathbf{E}^{-1}\mathbf{A}\mathbf{E}y(t) = \begin{bmatrix} \lambda_1 & & \\ & \ddots & \\ & & \lambda_n \end{bmatrix} y(t)$ Decoupled Equations! Decoupling : $\frac{dy_i(t)}{dt} = \lambda_i y_i(t) \Rightarrow y_i(t) = e^{\lambda_i t} y(0)$ Steps for Solving $\frac{dx(t)}{dt} = Ax(t), \quad x(0) = x_0$ 1.) Determine E, λ 2.) Compute $y(0) = \mathbf{E}^{-1} x_0$ 3.) Compute y(t) = $\begin{array}{c|c} \ddots \\ e^{\lambda_n t} \end{array} y(0)$ 4.) $x(t) = \mathbf{E}v(t)$ $y_1(0)$ - λ is real $\lambda < 0$ -> 1





Struts, Joints and Point Mass Example A 2×2 Example

$$y = y_0 + u$$

$$y = y_0 + u$$

$$f_s = EA_c \frac{y - y_0}{y_0} = \frac{EA_c}{y_0} u$$

$$f_s = EA_c \frac{y - y_0}{y_0} = \frac{EA_c}{y_0} u$$

$$f_s + f_m = 0$$

$$f_m = M \frac{d^2u}{dt^2}$$
Define v as velocity (du/dt) to yield a 2 × 2 System
$$\begin{cases} M \frac{d^2u}{dt^2} + \frac{EA_c}{y_0} u(t) = 0 \\ v(t) = \frac{du}{dt} \end{cases} \rightarrow \begin{cases} M \frac{dv}{dt} = -\frac{EA_c}{y_0} u(t) \\ \frac{du}{dt} = v(t) \end{cases}$$

$$\begin{bmatrix} M & 0\\ 0 & 1 \end{bmatrix} \begin{bmatrix} \frac{dv}{dt}\\ \frac{du}{dt} \end{bmatrix} = \begin{bmatrix} 0 & -\frac{EA_c}{y_0}\\ 1 & 0 \end{bmatrix} \begin{bmatrix} v\\ u \end{bmatrix}$$

Let
$$M = 1$$
, $\frac{LA_c}{y_0} = 1$

$$\begin{bmatrix} M & 0 \\ 0 & 1 \\ dt \\ dt \\ dt \end{bmatrix} = \begin{bmatrix} 0 & -\frac{EA_c}{y_0} \\ 1 & 0 \end{bmatrix} \begin{bmatrix} v \\ u \end{bmatrix} \qquad dx = \begin{bmatrix} 0 & -1.0 \\ 1.0 & 0 \\ A \end{bmatrix} x$$

Eigenvalues and Eigenvectors

 $\mathcal{A} = \begin{bmatrix} -1 & -1 \\ i & -i \end{bmatrix} \begin{bmatrix} i & 0 \\ -i \end{bmatrix} \begin{bmatrix} -1 & -1 \\ i & -i \end{bmatrix}^{-1} \qquad y_1(t) = e^s y_1(0) = y_1(0) (\cos t + i \sin t)$ $y_2(t) = e^{-s} y_2(0) = y_2(0) (\cos t - i \sin t)$

eigenvectors eigenvalues

 $\mathbf{A} \underbrace{\vec{E}}_{\text{real complex}} = \lambda \underbrace{\vec{E}}_{\text{complex}}$



Note the system has imaginary eigenvalues

Persistent Oscillation

Velocity, v, peaks when displacement, u, is zero.

Chemical Reaction Example

eigenvectors eigenvalues

✤ A 2×2 Example

Amount of reactant = R, the temperature = T $\frac{dT}{dt} = -T + R$ More reactant causes the temperature to rise, higher temperatures increase heat dissipation causing temperature to fall. $\frac{dR}{dt} = -R + 4T$ Higher temperature raises reaction rates, increased reactant interferes with reaction and slows rate. $\begin{bmatrix} \frac{dT}{dt} \\ \frac{dR}{dt} \end{bmatrix} = \begin{bmatrix} -1 & 1 \\ 4 & -1 \end{bmatrix} \begin{bmatrix} T \\ R \end{bmatrix} \xrightarrow{dt} \frac{dx}{dt} = \begin{bmatrix} -1 & 1 \\ 4 & -1 \end{bmatrix} x$ Eigenvalues and Eigenvectors $A = \begin{bmatrix} 1 & -1 \\ 2 & 2 \end{bmatrix} \begin{bmatrix} 1 & 0 \\ 0 & -3 \end{bmatrix} \begin{bmatrix} 1 & -1 \\ 2 & 2 \end{bmatrix}^{-1} y_1(t) = e' y_1(0)$



FINITE DIFFERENCE METHODS Basic Concepts



SECOND: Represent x(t) using values at t_i







Forward Euler
 Approximation

















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The Forward-Euler is accurate for small timesteps, but goes unstable when the timestep is enlarged

	F.E. explicit	B.E.	Trap
COST	cheap	expensive	expensive
ACCURACY	ok	ok	great
CONVERGENCE	linear	linear	quadratic
OSCILLATION	generates	dissipates	conserves
ADJUST Δt	unstable	ok	ok

* Summary

- Summary
 Convergence
 Did the computed solution approach the exact solution?
 Why did the trap rule approach faster than BE or FE?
 Energy Preservation

 Why did BE produce a decaying oscillation?
 Why did trap rule maintain oscillation amplitude?

 Two timeconstant (stiff) problems

 Why did FE go unstable when the timestep increased?

INTRODUCTION TO NUMERICAL SIMULATION

LECTURE 15.

Multistep Methods I

TODAY'S OUTLINE:

- Multistep Methods
 - e.g. FE, BE, TR
 - Convergence Definitions
 - Local Truncation Error (LTE) and Global Error (GE)
 - Example: LTE and GE for F.E.
 - Minimizing LTE (consistency)
 Minimizing the accumulation of errors (stability)



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Common Algorithms

Multistep Equation:	$\sum_{j=0}^{k} \alpha_j \hat{x}^{l-j} = \Delta t \sum_{j=0}^{k} \beta_j f\left(\hat{x}^{l-j}, u(t_{l-j})\right)$
Forward-Euler Approximation:	$x(t_{l}) \approx x(t_{l-1}) + \Delta t f(x(t_{l-1}), u(t_{l-1}))$
FE Discrete Equation:	$\hat{x}^{l} - \hat{x}^{l-1} = \Delta t f(x(t_{l-1}), u(t_{l-1}))$
Multistep Coefficients:	$k = 1$, $\alpha_0 = 1$, $\alpha_1 = -1$, $\beta_0 = 0$, $\beta_1 = 1$
BE Discrete Equation:	$\hat{x}^{l} - \hat{x}^{l-1} = \Delta t f \left(\hat{x}^{l}, u(t_{l}) \right)$
Multistep Coefficients:	$k = 1, \alpha_0 = 1, \alpha_1 = -1, \beta_0 = 1, \beta_1 = 0$
Trap Discrete Equation:	$\hat{x}^{l} - \hat{x}^{l-1} = \frac{\Delta t}{2} \left[f(\hat{x}^{l}, u(t_{l})) + f(\hat{x}^{l-1}, u(t_{l-1})) \right]$
Multistep Coefficients:	$k = 1, \alpha_0 = 1, \alpha_1 = -1, \beta_0 = \frac{1}{2}, \beta_1 = \frac{1}{2}$







$$\begin{split} \hat{x} &= \lambda x \quad x(0) = x_0 \\ \underline{k} &= 1 \text{ Multistep Method} \\ \hline \alpha_0 \hat{x}^l + \alpha_1 \hat{x}^{l-1} = \Delta t \left(\beta_0 \lambda \hat{x}^l + \beta_1 \lambda \hat{x}^{l-1} \right) \\ & \text{Forward} \quad \hat{x}^l = \hat{x}^{l-1} + \Delta t \cdot \lambda \hat{x}^{l-1} \quad \alpha_0 = 1 \quad \beta_0 = 0 \\ & \text{Euler} \quad \hat{x}^l = \hat{x}^{l-1} + \Delta t \cdot \lambda \hat{x}^{l-1} \quad \alpha_1 = -1 \quad \beta_1 = 1 \\ & \text{When does } x(t) \text{ grow?} \\ & \lambda > 0 \quad x(t) = x_0 e^{\lambda t} \\ & \text{What about F.E.?} \\ & \hat{x}^l = \underbrace{(1 + \Delta t \cdot \lambda)}_{\hat{x}^l (1 + \Delta t \cdot \lambda) > 1} \hat{x}^{l-1} \end{split}$$

Simplified Problems for Analysis

Scalar ODE : $\frac{d}{dt}v(t) = \lambda v(t), v(0) = v_0 \quad \lambda \in C$ Why such a simple Test Problem? • Nonlinear Analysis has many *unrevealing* subtleties • Scalar is equivalent to vector for multistep methods $\frac{d}{dt}x(t) = Ax(t) \xrightarrow{\text{multistep}} \sum_{j=0}^{k} \alpha_j \hat{x}^{l-j} = \Delta t \sum_{j=0}^{k} \beta_j A \hat{x}^{l-j}$ Let $Ey(t) = x(t) \xrightarrow{\text{multistep}} \sum_{j=0}^{k} \alpha_j \hat{y}^{l-j} = \Delta t \sum_{j=0}^{k} \beta_j E^{-1} A E \hat{y}^{l-j}$ • Decoupled Equations $\sum_{j=0}^{k} \alpha_j \hat{y}^{l-j} = \Delta t \sum_{j=0}^{k} \beta_j \left[\lambda_1 \cdot \cdot \cdot \lambda_n \right] \hat{y}^{l-j}$ Scalar Multistep formula: $\sum_{j=0}^{k} \alpha_j \hat{y}^{l-j} = \Delta t \sum_{j=0}^{k} \beta_j \lambda \hat{v}^{l-j}$ Must consider ALL $\lambda \in C$

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Global Error Equation

Multistep formula: Exact solution Amost satisfies Multistep Formula :

Global Error :

 $(t_{l-i}) - \Delta t \sum \beta$ $E^{l} \equiv v(t_{l}) - \hat{v}^{l}$ Difference equation relates LTE to Global error (L.FL) $(\alpha_0 - \lambda \Delta t \beta_0) E^l + (\alpha_1 - \lambda \Delta t \beta_1) E^{l-1} + \dots + (\alpha_k - \lambda \Delta t \beta_k) E^{l-k} = e^l$

 $^{-j} - \Delta t \sum_{i} \beta_{i} \lambda \hat{v}^{l-j} = 0$







$$E^{0} = 0$$

$$E^{1} = e^{1}$$

$$E^{2} = (1 + \Delta t \lambda)E^{1} + e^{2} = (1 + \Delta t \lambda)e^{1} + e^{2}$$

$$E^{3} = (1 + \Delta t \lambda)^{2}e^{1} + (1 + \Delta t \lambda)e^{2} + e^{3}$$

$$E' \approx \int_{\text{roughly}} \int_{t=1}^{t} e^{t}$$

$$E^{1} = e^{1}$$

$$E^{1} = e^{1}$$

$$\begin{split} \text{If } & \left| E^{I} \right| \text{ roughly } \sum \left| e^{I} \right| \leq \sum C(\Delta t)^{2} \\ & \left| E^{\frac{I}{T_{obs}^{I}}} \right| \approx \frac{T}{\Delta t} + \underbrace{C(\Delta t)^{2}}_{\substack{\text{such charged} \\ \text{standards failer}}} \underbrace{\frac{1}{2} \max_{\substack{\text{such charged} \\ \text{standards failer}}} \frac{1}{2} \max_{\substack{\text{such charged} \\ \text{such charged} \\ \text{such charged} }} \frac{1}{2} \max_{\substack{\text{such charged} \\ \text{such charged} \\ \text{su$$

← Global Error Equation Forward-Euler definition $\hat{v}^{l+1} = \hat{v}^l + \Delta r_\lambda \hat{v}^l$ Using the LTE definition $v((l+1)\Delta t) = v(l\Delta t) + \Delta r \lambda v(l\Delta t) + e^l$ Subtracting yields global error equation $E^{l+1} = (l + \Delta r \lambda)E^l + e^l$ Using magnitudes and the bound on e^l $E^{l+1} = |I + \Delta r \lambda]E^l + |e^l| \le (1 + \Delta r |\lambda|)E^l| + C(\Delta t)^2$

♦ A helpful bound on difference equations A lemma bounding difference equation solutions If $|u^{l+1}| \le (1 + ε)|u^l| + b, u^0 = 0, ε > 0$ Then $|u'| \le \frac{e^{2l}}{ε}|b|$

 $u^{1} \leq (1+\varepsilon)(u^{0}) + b = b$ $u^{2} \leq (1+\varepsilon)b + b$ $u^{3} \leq (1+\varepsilon)^{2}b + (1+\varepsilon)b + b$



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- · Forward-Euler is order 1 convergent
- · Bound grows exponentially with time interval
- C related to exact solution's second derivative
- * Exact and Forward-Euler (FE) Plots for Unstable Reactions



* Forward-Euler Errors for Solving Reaction Equation



Exact and forward-Euler (FE) plots for circuit.



Forward-Euler (FE) errors for solving circuit equation.





Local Truncation Error:
$$\sum_{j=0}^{k} \alpha_{j} v(t_{l-j}) - \Delta t \sum_{j=0}^{k} \beta_{j} \frac{d}{dt} v(t_{l-j}) = e^{t}$$
If $v(t) = t^{p} \Rightarrow \frac{d}{dt} v(t) = pt^{p-1}$

$$\sum_{j=0}^{k} \alpha_{j} \frac{((k-j)\Delta t)^{p}}{v(t_{k-j})} - \Delta t \sum_{j=0}^{k} \beta_{j} \frac{p((k-j)\Delta t)^{p-1}}{dt} = e^{k}$$

$$\Rightarrow (\Delta t)^{p} \left[\sum_{j=0}^{k} \alpha_{j} (l-j)^{p} - \sum_{j=0}^{k} \beta_{j} p(l-j)^{p-1} \right] = e^{k}$$
If $\sum_{j=0}^{k} \alpha_{j} (k-j)^{p} - \sum_{j=0}^{k} \beta_{j} p(k-j)^{p-1} = 0$ then $e^{k} = 0$ for $v(t) = t^{p}$
As any smooth $v(t)$ has a locally accurate Taylor series in t :
If $\sum_{j=0}^{k} \alpha_{j} v(t_{l-j}) - \sum_{j=0}^{k} \beta_{j} p(k-j)^{p-1} = 0$ for all $p \leq p_{0}$
Then $\sum_{j=0}^{k} \alpha_{j} v(t_{l-j}) - \sum_{j=0}^{k} \beta_{j} \frac{d}{dt} (v_{l-j}) = e^{l} = \underbrace{C(\Delta t)^{p_{0}+1}}_{c(l-1)\Delta t}$

$$x(t) = t \quad \frac{dx(t)}{dt} = 1$$





 $-\sum_{k=0}^{n} \beta_{i} p(k-j)^{p-1} = 0$ Exactness Constraints : $\sum \alpha_i (k-j)^p$ j=0

For k = 2, yields a 5×6 system of equations for Coefficients



Forward-Euler $\alpha_0 = 1$, $\alpha_1 = -1$, $\alpha_2 = 0$, $\beta_0 = 0$, $\beta_1 = 1$, $\beta_2 = 0$, FE satisfies p = 0 and p = 1 but not $p = 2 \Rightarrow LTE = C(\Delta t)^2$ Backward-Euler $\alpha_0 = 1$, $\alpha_1 = -1$, $\alpha_2 = 0$, $\beta_0 = 1$, $\beta_1 = 0$, $\beta_2 = 0$, BE satisfies p = 0 and p = 1 but not $p = 2 \Rightarrow LTE = C(\Delta t)^2$ Trapezoidal Rule $\alpha_0 = 1$, $\alpha_1 = -1$, $\alpha_2 = 0$, $\beta_0 = 0.5$, $\beta_1 = 0.5$, $\beta_2 = 0$, Trapezoidal satisfies p = 0, 1, or 2 but not $p = 3 \Rightarrow LTE = C(\Delta t)^3$ First introduce a normalization, for example $\alpha_0 = 1$ Solve for the 2-step method with lowest LTE $\alpha_0 = 1, \alpha_1 = 0, \alpha_2 = -1, \beta_0 = 1/3, \beta_1 = 4/3, \beta_2 = 1/3$ Satisfies all five exactness constraints \Rightarrow LTE = $C(\Delta t)^5$ Solve for the 2-step explicit method with lowest LTE $\alpha_0 = 1, \alpha_1 = 4, \alpha_2 = -5, \beta_0 = 0, \beta_1 = 4, \beta_2 = 2$ Satisfies all five exactness constraints \Rightarrow LTE = $C(\Delta t)^4$









For a "good" method

$$E^{I} \cong k \sum_{i=1}^{I} e^{i} \quad I \propto \frac{T}{\Delta t}$$
FELTE Δt^{2}

$$|E^{I}| \leq kI \max_{i} |e^{i}|$$



Stability Definition Multistep Method Difference Equation $\begin{array}{l} (\alpha_0 - \lambda \Delta t \beta_0) E^{l} + (\alpha_1 - \lambda \Delta t \beta_1) E^{l-1} + \dots + (\alpha_k - \lambda \Delta t \beta_k) E^{l-k} = e^{l} \\ \text{Definition: A multistep method is stable if as } \Delta t \to 0 \\ \max |E^{l}| \leq C(T) \frac{T}{2} \max |e^{l}| \end{array}$

$$\max_{I \in \left[0, \frac{T}{\Delta I}\right]} \left| \frac{E^{i}}{\sum_{i \in \left[0, \frac{T}{\Delta I}\right]}} \right| \leq \frac{C(T)}{\sum_{i \in \left[0, \frac{T}{\Delta I}\right]}} \frac{1}{\sum_{i \in \left[0, \frac{T}{\Delta I}\right]}} \left| \frac{e^{i}}{\sum_{i \in \left[0, \frac{T}{\Delta I}\right]}} \right|$$

Stability means: Global error is bounded by a constant times the sum of the L.T.E.'s

Solving Difference Equations

* Convolution Consider a general kth order difference equation $a_0 x^l + a_1 x^{l-1} + \dots + a_k x^{l-k} = u^l$ Which must have k initial conditions $\begin{aligned} x^0 &= x_0, \ x^{-1} = x_1, \ \cdots, \ x^{-k} = x_k \\ \text{As is clear when the equation is in update form} \\ x^1 &= -\frac{1}{a_0} \Big(a_1 x^0 + \cdots + a_k x^{-k+1} - u^1 \Big) \end{aligned}$ It can be shown that the solution of a difference equation is simply a convolution sum: x can be related to u by $x^{l} = \sum_{j=0}^{l} h^{l-j} u^{j}$ Calculating h If $a_0 z^k + a_1 z^{k-1} + \dots + a_k = 0$ has distinct roots $\zeta_1, \zeta_2, \dots, \zeta_k$ Then $x^{\ell} = \sum_{j=0}^{k} h^{\ell j} u^{j}$ where $h^{\ell} = \sum_{j=1}^{k} \gamma_{j} (\zeta_{j})^{\ell}$ To understand how h is derived, first a simple case Suppose $x^{\ell} = \zeta x^{\ell-1} + u^{\ell}$ and $x^{-1} = 0$ $x^{l} = \varsigma x^{0} + u^{1} = u^{1}, \quad x^{2} = \varsigma x^{1} + u^{2} = \varsigma u^{1} + u^{2}$ $x_{j=0}^{l} = \sum_{j=0}^{l} \varsigma_{j=0}^{l-j} u^{j}$ time index time index * Three Important Observations If $|\varsigma_i| < 1$ for all *i*, then $|x^i| \le C \max |u^{ij}|$ where *C* does not depend on *l* If $|\varsigma_i| > 1$ for any *i*, then there exists a bounded u^j such that $|x^l| \to \infty$

If $|\varsigma_i| \le 1$ for all *i*, and if $|\varsigma_i| = 1$, ς_i is distinct then $|x^i| \le CI \max |u^j|$

$$x^{l} = \varsigma x^{l-1} + u^{l} \qquad x^{0} = 0$$

$$x^{1} = \varsigma x^{0} + u^{1}$$

$$x^{2} = \varsigma x^{1} + u^{2} = \varsigma u^{1} + u^{2}$$

$$x^{l} = \varsigma^{l-1}u^{1} + \varsigma^{l-2}u^{2} + \varsigma^{l-3}u^{3} + \dots u^{l}$$

• Convolution Sum
• When roots are not distinct
• When roots are not distinct
•
$$h^{I} = \sum_{q=1}^{Q} \sum_{m=0}^{M_{q}-1} \gamma_{q,m}(l)^{m} (\varsigma_{q})^{l} \qquad x^{I} = \sum_{j=0}^{I} h^{I-j} u^{j}$$

Roots of $a_{0}z^{k} + a_{1}z^{k-1} + \dots + a_{k} = 0$
• Bounding Terms
 $x^{I} = \sum_{q=1}^{Q} \sum_{m=0}^{M_{q}-1} \left[\sum_{j=0}^{I} \gamma_{q,m}(l-j)^{m} (\varsigma_{q})^{l-j} u^{j} \right]$
If $|\varsigma_{i}| < 1$, then $|R_{q,m}| \le C \max_{i} |u^{i}|$
Independent of l
If $|\varsigma_{q}| < (1+\varepsilon)$, then $|R_{q,0}| \le C \frac{e^{t^{2}}}{\varepsilon} \max_{j} |u^{j}|$
Lemma: bounds distinct roots

Stability Theorem

•

Theorem: A multistep method is stable if and only if Roots of $\alpha_0 z^k + \alpha_1 z^{k-1} + \ldots + \alpha_k = 0$ either: 1. Have magnitude less than one 2. Have magnitude equal to one and are distinct Note: from the previous slides it is easy to see that a multistep method is stable if the same conditions hold for: $(\alpha_0 - \lambda \Delta t \beta_0) z^k + (\alpha_1 - \lambda \Delta t \beta_1) z^{k-1} + \dots + (\alpha_k - \lambda \Delta t \beta_k) = 0$ But the theorem says: forget about the B's ...





Dahlquist's First Stability Barrier

For a stable, explicit k-steps method, the maximum degree of a polynomial t^p that can be integrated exactly is less than or equal to k: $p_0 \le k$ (note there are 2k coefficients).

For a stable, implicit k-steps method, (with $2k + 1 \operatorname{coeff}$) $p_0 \le k + 2$ if k is even $p_0 \le k + 1$ if k is odd STABLE: EXPLICIT IMPLICIT k = 1 step $p_0 \le 1$ (FE)k = 2 steps $p_0 \le 2$ k = 3 steps $p_0 \le 3$ $p_0 \le 2 \text{ (TR)}$ $p_0 \le 4$ $p_0 \le 4$

2. Global Condition: One step errors grow slowly (stability)

INTRODUCTION TO NUMERICAL SIMULATION

LECTURE 16.

Multistep Methods II: Small and Large Time-Steps Stability

TODAY'S OUTLINE:

SMALL Timestep Issues for Multistep Methods
 Stability

- Difference Equations
 Stability & Consistency implies Convergence
- Stability & Consistency implies Converget
 LARGE Timestep Issues
- Absolute Stability for two time-scale examples.
 - Oscillators



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We made the LTE so small, how come the Global error is so large?

Stability Definition

Multistep Method Difference Equation $(\alpha_0 - \lambda \Delta t \beta_0) E^l + (\alpha_1 - \lambda \Delta t \beta_1) E^{l-1} + \dots + (\alpha_k - \lambda \Delta t \beta_k) E^{l-k} = e^l$ Definition: A multistep method is stable if as $\Delta t \to 0$ $| \Delta t | = -t = -t = T$

$$\max_{l \in \left[0, \frac{T}{\Delta t}\right]} E^{l} \leq \underbrace{C(T)}_{\text{interval}} \frac{I}{\Delta t} \max_{l \in \left[0, \frac{T}{\Delta t}\right]} e^{l}$$

Stability means: Global Error is bounded by a constant times the sum of the LTE's

Solving Difference Equations

x

♦ Convolution Consider a general kth order difference equation a₀x^l + a₁x^{l-1} + ··· + a_kx^{l-k} = u^l Which must have k initial conditions x⁰ = x₀, x⁻¹ = x₁, ··· x^{-k} = x_k as is clear when the equation is in update form x^l = -1/a₀ (a₁x⁰ + ··· + a_kx^{-k+1} - u^l)

it can be shown that the solution of a difference equation is simply a convolution sum:

can be related to
$$u$$
 by $x^{l} = \sum_{i=0}^{l} h^{l-j} u^{j}$

$$BE \qquad \psi(t) = \lambda v$$

$$\psi' = \psi(t) = \lambda v$$

$$\psi' = \psi^{t-1} - \Delta t \lambda \psi' = 0$$

$$v(t\Delta t) - v((t-1)\Delta t) - \Delta t \lambda v(t\Delta t) = e^{t}$$

$$(1 - \Delta t \lambda) \quad E^{t} = -E^{t-1} = e^{t}$$

$$E^{1} = \frac{1}{1 - \Delta t \lambda} E^{0} + \frac{1}{1 - \Delta t \lambda} e^{1}$$

$$E^{2} = \frac{1}{1 - \Delta t \lambda} E^{1} + \frac{1}{1 - \Delta t \lambda} e^{2} = \left(\frac{1}{1 - \Delta t \lambda}\right)^{2} e^{1} + \frac{1}{1 - \Delta t \lambda} e^{2}$$

$$E^{t} = \left(\frac{1}{1 - \Delta t \lambda}\right)^{t} e^{1} + \left(\frac{1}{1 - \Delta t \lambda}\right)^{t-1} e^{2} + \dots + \frac{1}{1 - \Delta t \lambda} e^{t}$$

$$E^{t} = \sum_{j=1}^{t} \left(\frac{1}{1 - \Delta t \lambda}\right)^{t-j} \left(\frac{1}{1 - \Delta t \lambda}\right) e^{j}$$

← Calculating *h* If $a_0 z^k + a_1 z^{k-1} + \dots + a_k = 0$ has distinct roots $\zeta_1, \zeta_2, \dots, \zeta_k$ Then $x^l = \sum_{j=0}^l h^{l j} u^j$ where $h^l = \sum_{j=1}^k \gamma_j (\zeta_j)^j$ To understand how *h* is derived, first a simple case

To understand how *h* is derived, first a simple case Suppose $x^{l} = \varsigma x^{l-1} + u^{l}$ and $x^{-1} = 0$ $x^{1} = \varsigma y^{0} + u^{1} = u^{1}$, $x^{2} = \varsigma x^{1} + u^{2} = \varsigma u^{1} + u^{2}$

 $x_{1}^{l} = \sum_{j=0}^{l} \zeta_{i}^{l-j} u^{j} \underbrace{\text{time index}}_{\text{index}}$

♦ Three Important Observations If $|\varsigma_i| < 1$ for all *i*, then $|x^{l}| \le C \max_j |u^{j}|$ where *C* does not depend on *l* If $|\varsigma_i| > 1$ for all *i*, then there exists a bounded u^{j} such that $|x^{l}| \to \infty$

 $\frac{\text{If } |\varsigma_i| \le 1 \text{ for all } i, \text{ and } i |\varsigma_i| = 1, \varsigma_i \text{ is distinct then } |x'| \le Cl \max_{i < j} |u'|}{|u_j|}$

$$x^{l} = \lambda x^{l-1} + u^{l}$$

$$x^{-1} = 0$$

$$x^{0} = u^{0}$$

$$x^{1} = \lambda x^{0} + u^{1} = \lambda u^{0} + u^{1}$$

$$x^{2} = \lambda x^{1} + u^{2} = \lambda^{2} u^{0} + \lambda u^{1} + u^{2}$$

$$x^{l} = \sum_{i=0}^{l} \lambda^{l-i} u^{i}$$

Convolution Sum
 When roots are not distinct

o Bounding Terms

$$x^{l} = \sum_{q=1}^{O} \sum_{m=0}^{M_{q}-1} \left[\sum_{j=0}^{l} \gamma_{q,m} (l-j)^{m} (\varsigma_{q})^{j-j} u^{j} \right]$$
If $|\varsigma_{1}| < 1$, then $|R_{q,m}| \le C \max_{R_{q,m}} |u^{j}|$
If $|\varsigma_{q}| < (1 + \varepsilon)$, then $|R_{q,0}| \le C \frac{\varepsilon^{2l}}{\varepsilon} \max_{j} |u^{j}|$
Lemma: bounds distinct roots

Stability Theorem

 $\begin{array}{ll} \hline \text{Theorem: A multistep method is stable if and only if}\\ \text{Roots of } & \alpha_0 z^k + \alpha_1 z^{k-1} + \cdots + \alpha_k = 0 \text{ either:} \\ \bullet & \text{Have magnitude equal to one and are distinct}\\ \text{Note: from the previous slides it is easy to see that a multistep method is stable if the same conditions hold for:}\\ & (\alpha_0 - \lambda \Delta \beta_0) z^k + (\alpha_1 - \lambda \Delta \beta_1) z^{k-1} + \cdots + (\alpha_k - \lambda \Delta \beta_k) = 0\\ \text{But the theorem says: forget about the betas...}\\ \bullet & \text{''Proof''}\\ & \text{Given the Multistep Method Difference Equation}\\ & (\alpha_0 - \lambda \Delta \beta_0) E^l + (\alpha_1 - \lambda \Delta \beta_1) E^{l-1} + \cdots + (\alpha_k - \lambda \Delta \beta_k) E^{l-k} = e^l\\ \text{If, as } \Delta t \to 0, \text{ roots of } (\alpha_0 - \lambda \Delta \beta_0) z^{k} + \ldots + (\alpha_k - \lambda \Delta \beta_k) = 0\\ & \circ & \text{Less than one in magnitude or}\\ & \circ & \text{Are distinct and bounded by } 1 + \kappa \Delta t, \kappa > 0\\ \text{Then from the aside on difference equations}\\ & \\ & max_k \\ & I_k \in \left[0, \frac{T}{\Delta t}\right] E^l \left| \leq C \frac{e^{\kappa \Delta t}}{\Delta t} \max_{l \in \left[0, \frac{T}{\Delta t}\right]} \left| e^l \right| \leq \frac{Ce^{\kappa T}}{C(T)} \frac{T}{\Delta t} \max_{l \in \left[0, \frac{T}{\Delta t}\right]} \right| e^l \end{aligned}$

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Best explicit 2-step method $\alpha_0 = 1$, $\alpha_1 = 4$, $\alpha_2 = -5$, $\beta_0 = 0$, $\beta_1 = 4$, $\beta_2 = 2$





Dahlquist's First Stability Barrier

For a stable, explicit k-steps method, the maximum degree of a polynomial t^{ρ} tat can be integrated exactly is less than or equal to k: $p_0 \leq k$ (note there are 2k coefficients).

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For a stable, implicit k-steps method, (with 2k+1 coeff) $p_0 \le k+2$ if k is even $p_0 \le k+1$ if k is odd

STABLE:	EXPLICIT:	IMPLICIT:
k = 1 step	$p_0 \le 1 (FE)$	$p_0 \leq 2 (\text{TR})$
k = 2 steps	$p_0 \leq 2$	$p_0 \leq 4$
k = 3 steps	$p_0 \leq 3$	$p_0 \leq 4$

Conditions for Convergence, Stability, and Consistency

 Local Condition: One step errors are small (consistency) Exactness Constraints up to *p₀* (*p₀* must be > 0) ⇒ max <sub>*l*∈[0, ^T/_M] |*E^l*| ≤ C₁(Δ*t*)^{*p*₀+1} for Δ*l* < Δ*t*₀
 Global Condition: One step errors grow slowly (stability) roots of ^k_{j=0} α_j z^{k-j} = 0 Inside the unit circle and distinct ⇒ max <sub>*l*∈[0, ^T/_M] |*E^l*| ≤ C₂ ^T/<sub>Δ*t*</sup> <sub>*l*∈[0, ^T/_M]} |*e^l*|
</sub></sub></sub></sub>

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 $i_{C1} \vee i_{R1} \vee i_{R2}$



















Stability Definitions Region of Absolute Stability for a Multistep method:

Values of $\lambda \Delta t$ where roots of $\sum_{j=0}^{k} (\alpha_j - \lambda \Delta t \beta_j) z^{k-j} = 0$ are inside the unit circle.

A-Stable: A-stable: A method is A-stable if its region of absolute stability includes the entire left-half of the complex plane Dahlquist's Second Stability Barrier: There are no A-stable multistep methods of convergence order greater than 2, and the trap rule is the most accurate.







Why does FE result grow, BE result decay and the Trap rule preserve oscillations FE Large Timestep Oscillator Example







Trap Large Timestep Oscillator Example



- Large Timestep Issues
 o Two Time-Constant Stable Problem (Circuit)
 - FE: stability, not accuracy, limited timestep size
 BE: was A-stable, any timestep could be used.

 - Trap Rule most accurate A-stable m-step method
 - o Oscillator Problem
 - · Forward-Euler generated an unstable difference equation regardless of timestep size.


- Backward-Euler generated a stable (decaying) difference equation
- regardless of timestep size.Trapezoidal rule mapped the imaginary axis to the unit circle, regardless of timestep size:

 - Decaying ODE are mapped to stable difference equations
 Unstable ODE are mapped to unstable difference equations
 Osciallating ODE are mapped to oscillating difference equations

- SUMMARY ◆ Small Timestep Issues for Multistep Methods Local truncation error and exactness
- O Local truncation error and exactness
 D lifference equation stability
 Stability + consistency implies convergence
 Large Timestep Issues
 Absolute stability for two time-scale examples
 Oscillators

INTRODUCTION TO NUMERICAL SIMULATION

LECTURE 17.

Methods for Computing Periodic Steady-State

TODAY'S OUTLINE:

- * Periodic Steady-State Problems
- Application examples and simple cases
- Review Full Time Integration Methodss
 Finite-Difference Methods for Periodic Steady-State
- Shooting Methods
 - State Transition Function
 - Sensitivity Matrix
- Matrix Free Approach
- Spectral Methods (Harmonic Balance)
 - Galerkin Methods
 - Collocation Methods









 $\frac{dx(t)}{dt} = F(x(t)) + u(t)$ dt • Then if *u* is periodic with period *T* and $x(t_0 + T) = x(t)$ for some t_0 $\Rightarrow x(t+T) = x(t)$ for all $t > t_0$



ω 17 ODE Τ. Periodic Steady State Analysis









 $x(t) = ke^{i\omega t}$ for linear systems

Simple Example [RLC Filter, Spring+Mass+Dashpot]







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REVIEW FULL TIME INTEGRATION METHODS Computing Steady State

 $\begin{aligned} \hat{x}^{i} &- \hat{x}^{i-1} = F(\hat{x}^{i}) + u(i\Delta t) \quad \text{Solve for } \hat{x}^{i} \text{ using Newton} \\ \Delta t & \alpha_{0} = 1 \quad \alpha_{1} = -1 \quad \beta_{0} = 1 \end{aligned}$

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 $\frac{dx(t)}{dt} = F$ $\frac{x(t)}{state}$ $x(0) = x_0$ $+ \underbrace{u(t)}_{input}$ dt Initial Condition • Backward Euler Equation for timester l $\hat{x}^{l} - \hat{x}^{l-1} = \Delta t \left[F(\hat{x}^{l}) + u(l\Delta t) \right]$ How do we solve the backward-Euler Equation?













LU or GCR?

In GCR: at each step need matrix vector product

$$\begin{bmatrix} \alpha_0 I - \Delta t \beta_0 \frac{\partial f}{\partial x} (\hat{x}^{l,j}, u(t_l)) \end{bmatrix} r^k \cong$$

$$\cong \alpha_0 r^k - \frac{\Delta t \beta_0}{\epsilon} \left[f (\hat{x}^{l,j} + \epsilon r^k, u(t_l)) - f (\hat{x}^{l,j}, u(t_l)) \right]$$





FINITE-DIFFERENCE METHODS FOR PERIODIC STEADY-STATE
Boundary-Value Problem
Basic Formulation



N Differential Equations : $\frac{d}{dt}x_i(t) = F_i(x(t))$ N Periodicity Constraints : $x_i(T) = x_i(0)$



← Finite Difference Methods ○ Linear Example Problem $\frac{dx(t)}{dt} = Ax(t) + \underbrace{u(t)}_{input} \quad t \in [0, T] \quad \underbrace{x(T) = x(t)}_{periodicity}$ constraint Discretize e.g. with Backward-Euler $\hat{x}^1 = x(0) + \Delta t \cdot (A\hat{x}^1 + u(\Delta t))$ $\hat{x}^2 = \hat{x}^1 + \Delta t \cdot (A\hat{x}^2 + u(2\Delta t))$ \vdots $\hat{x}^L = \hat{x}^{L-1} + \Delta t \cdot (A\hat{x}^L + u(L\Delta t))$ Periodicity implies $\hat{x}^L = \hat{x}^0$





 $\begin{aligned} \frac{d\kappa(t)}{dt} &= F(x(t)) + \underbrace{u(t)}_{input} \quad t \in [0,T] \quad \underbrace{x(T) = x(t)}_{periodicity} \\ \text{Discretize e.g. with Backward-Euler} \\ H_{FD} \begin{bmatrix} \begin{bmatrix} \hat{x}^1 \\ \hat{x}^2 \\ \vdots \\ \hat{x}^L \end{bmatrix} = \begin{bmatrix} \hat{x}^1 - \hat{x}^L - \Delta t \begin{bmatrix} F(\hat{x}^1) + u(\Delta t) \end{bmatrix} \\ \hat{x}^L - \hat{x}^L - \Delta t \begin{bmatrix} F(\hat{x}^2) + u(2\Delta t) \end{bmatrix} = 0 \\ \vdots \\ \hat{x}^L - \hat{x}^L - 1 - \Delta t \begin{bmatrix} F(\hat{x}^L) + u(L\Delta t) \end{bmatrix} \\ \text{Solve the huge system using Newton's Method} \end{aligned}$



-NL-





 $x(t) = \mathbf{A}e^{\lambda t}$
$$\begin{split} x_{t_0} &= \mathbf{A} e^{\lambda t_0} \to \mathbf{A} = x_{t_0} e^{-\lambda t_0} \\ x(t) &= x_{t_0} e^{\lambda(t-t_0)} \qquad \qquad x(t_1) = x_{t_0} e^{\lambda(t_1-t_0)} \end{split}$$







 $\begin{aligned} & \bigstar \quad \text{State Transition Function Example} \\ & \frac{dx(t)}{dt} = \lambda x(t) \\ & \Phi \Big(x_{t_0}, t_0, t_1 \Big) = e^{\lambda (t_1 - t_0)} x_{t_0} \end{aligned}$



★ Abstract Formulation Solve $\frac{x(t_0 + T)}{\Phi(x_{t_0, t_0, t_0} + T) = x_{t_0}} = x_{t_0}$ for a t_0 of your choice e.g. pick $t_0 = 0$ $H(x_0) = \Phi(x_0, 0, T) - x_0 = 0$

> Evaluate $H(x^k)$ Integrate $\frac{d}{dt}x(t) = F(x(t)) + u(t)$ $x(0) = x^k$ Over the interval [0,T]

 $\begin{aligned} x_0 & \text{unknown} \\ H(x_0) &= 0 \\ \underline{\phi(x_0, 0, T) - x_0}_{H(x_0)} &= 0 \\ \frac{\partial H}{\partial x_0} & \begin{bmatrix} \frac{\partial \phi}{\partial x_0} \\ & -\mathbf{I} \end{bmatrix} \Delta x_0^{k+1} = -\left[\phi(x_0^k, 0, T) - x_0^k\right] \\ & \text{UNTIL} \left[\Delta x_0^{k+1}\right], \left[H(x_0^k) \text{ small}\right] \end{aligned}$

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Solve $H(x_0) = 0$ for x_0 using Newton's method $J_H(x_0^k)(x_0^{k+1} - x_0^k) = -H(x_0^k)$

where $J_H(x_0^k) = \frac{\partial \Phi(x_0^k, 0, T)}{\partial x_0} - I$ $k \leftarrow 0$ Note this is not so different to guess anymore; $x^k \leftarrow x_0$ it is just one distribution of temperatures. REPEAT $\left[\frac{\partial \phi}{\partial x_0}\right]_{x_0^k}$ $-\mathbf{I} \left[\Delta x_0^{k+1} = -\left[\phi(x_0^k, 0, T) - x_0^k \right] \right]$ $x_0^{k+1} \leftarrow x_0^k + \Delta x_0^{k+1}$ $k \leftarrow k+1$ UNTIL $\left|\Delta x_0^k\right|, \left|\phi\left(x_0^{k-1}, 0, T\right) - x_0^{k-1}\right|$ small STEP 1 STEP 2 STEP k $x_1(t)$ $x_1(t)$ $x_1(t)$ $\phi(x_0^1,0,T) \neq x_0^1$ $\phi(x_0^2, 0, T) \neq x_0^2$ $\phi(x_0^k, 0, T)$ $x_0^{1/2}$ $= \frac{x_0^k}{T_t^0}$ T t Not Converged Not Converged Converged



















 Observations on Sensitivity Matrix Newton at each timestep uses same matrices

$$\frac{\partial \Phi(x,0,T)}{\partial x} \approx \prod_{l=1}^{L} \left(\underbrace{I - \Delta t \frac{\partial F(\tilde{x}^{l})}{\partial x}}_{Timestep} \right)^{-1}$$
Timestep
Newton
Jacobian
Formula simplifies in the linear case
$$\frac{\partial \Phi(x,0,T)}{\partial x} \approx (I - \Delta t A)^{-L}$$

$$\begin{aligned} \frac{dx}{dt} &= \mathbf{A}x(t) + u(t) & F'(x(t)) = \mathbf{A}x(t) \\ \mathbf{x}(t) &= \underbrace{e^{\mathbf{A}'}x(\mathbf{0})}_{\downarrow} & T = \Delta t \cdot t \\ \begin{bmatrix} \widehat{c}\Phi\\\widehat{c}x \end{bmatrix} \Delta x(\mathbf{0}) &\cong \phi(x(\mathbf{0}) + \Delta x(\mathbf{0}), \mathbf{0}, T) - \phi(x(\mathbf{0}), \mathbf{0}, T) \\ &\cong e^{\mathbf{A}\Delta t}[x(\mathbf{0}) + \Delta x] - e^{\mathbf{A}\Delta t}x(\mathbf{0}) \\ &= [e^{-\mathbf{A}\Delta t}]^{-t} \Delta x(\mathbf{0}) \\ &= [e^{-\mathbf{A}\Delta t}]^{-t} & \text{First Order Expansion (B.E.)} \end{aligned}$$

Matrix-Free Approach Basic Setup

Basic Schup REPEAT $\begin{bmatrix}
\frac{\partial \phi}{\partial x_0} |_{x_0^k} - \mathbf{I} \\
\frac{\int \Delta x_0^{k+1}}{x_0^{k+1} \leftarrow x_0^k} + \Delta x_0^{k+1} = -\left[\phi(x_0^k, 0, T) - x_0^k\right] \\
\frac{\chi_0^{k+1}}{k} \leftarrow k + 1 \\
\text{UNTIL} |\Delta x_0^k|, ||\phi(x_0^{k-1}, 0, T) - x_0^{k-1}|| \text{ small}$

Start with $\frac{dx(t)}{dt} = F(x(t)) + u(t)$ $H(x_0) = \Phi(x_0, 0, T) - x_0 = 0$ Solve $H(x_0) = 0$ for x_0 using Newton's method $J_H(x_0^k) x_0^{k+1} - x_0^k) = -H(x_0^k)$ where $J_H(x_0^k) = \frac{\partial \Phi(x_0^k, 0, T)}{\partial x_0} - I$











F.D. represents one period using *L* points

Maybe we can use other ways to represent the shape

 $Periodic \rightarrow use \ Fourier \ Coefficients$



Works well (need few coefficients) if function smooth harmonics Let's solve for x_l

SPECTRAL METHODS (HARMONIC BALANCE) Fourier Representation Truncation Approximation

• Periodic Function – Fourier Series

$$x(t) = \sum_{l=-\infty}^{\infty} X_l e^{-i2\pi l \frac{l}{T}}$$
• Approximate a function with truncated series

$$x(t) \approx \sum_{l=-L}^{L} X_l e^{-i2\pi l \frac{l}{T}}$$











★ Residual c Plug representation in differential equation $\frac{R(\vec{X}, t)}{\text{Residual}} = \frac{d}{dt} \left(\sum_{l=-L}^{L} X_{l} e^{-i2\pi l t} \right) - F\left(\sum_{l=-L}^{L} X_{l} e^{-i2\pi l t} \right) - u(t)$

























INTRODUCTION TO NUMERICAL SIMULATION

LECTURE 18.

Finite Difference Methods for Boundary Value Problems

TODAY'S OUTLINE:

- Informal Finite Difference Methods
- Heat Conducting Bar
- Comparing FEM and FD in 1-D
- Convergence Analysis for Heat Equation
 - Local Truncation Error and Consistency
- Global Error and Stability
 Formal Generalization to any PDE
 - Consistency + Stability yields Convergence

INFORMAL FINITE DIFFERENCE METHODS

- Heat Conducting Bar
 - Cut the bar into short sections
 Assign each cut a temperature





 $\frac{\partial}{\partial x} \kappa \frac{\partial T(x)}{\partial x} = -h_s \Rightarrow -\frac{\partial^2 u(x)}{\partial x^2} = f(x)$ $-u_{xx}(x) = f(x)$

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Errors accumulate but they need to decrease back to zero on the other side thanks to boundary conditions (B.C.) (*This is probably an easier problem*!)













Backward Euler $\hat{x}^1 - \hat{x}^0 = \Delta t \lambda \left[\lambda \hat{x}^1 + u(t_1) \right] \Longrightarrow \hat{x}^1 - x(0) = \Delta t \lambda \left[\lambda \hat{x}^1 + u(t_1) \right]$ $\int 1 - \lambda \Delta t$ $\left\lceil u(t_1) + x(0) \right\rceil$ \hat{x}^1 \hat{x}^2 -1 $1 - \lambda \Delta t$ $u(t_2)$ -1 $1 - \lambda \Delta t \int \hat{x}^{t}$ $u(t_1)$ Lower Triangular input $\hat{u} = \begin{pmatrix} \hat{u}_1 \\ \hat{u}_2 \\ \vdots \\ \hat{u}_{n-1} \end{pmatrix},$ (2 -1 $f(x_1)$ -1 2 -1 $f(x_2)$ $\mathbf{A} = \frac{1}{\Delta x^2}$ ··. ··. ··. -1 2 -1 f = $f(x_{n-1})$ -1 2 û, $f(x_n)$ (Symmetric) $\mathbf{A} \in \mathfrak{R}^{n \times n} \quad \underline{\hat{u}}, \underline{f} \in \mathfrak{R}^{n}$ Is **A** nonsingular? For any $\underline{v} = \{v_1, v_2, \dots, v_n\}^T$ $\underline{v}^T \mathbf{A} \underline{v} = \frac{1}{\Delta x^2} \left(v_1^2 + \sum_{i=2}^n (v_i - v_{i-1})^2 + v_n^2\right)$ Hence $\underline{v}^T \mathbf{A} \underline{v} > 0$, for any $\underline{v} \neq 0$ (A is SPD) $\mathbf{A}\underline{\hat{u}} = f$: $\underline{\hat{u}}$ exists and is unique



COMPARING FEM AND FD IN 1-D

Residual Equation
Partial Differential Equation Form

$$-\frac{\partial^{2} u}{\partial x^{2}} = f \quad u(0) = 0 \quad u(1) = 0$$
Step 1: Choose Basis Functions to represent the solution

$$u(x) \approx u_{h}(x) = \sum_{i=1}^{n} \omega_{i} \frac{\varphi_{i}(x)}{p_{insterments}}$$
Step 2: Generate equations for the basis functions weights setting residual
orthogonal to some test functions

$$\int_{0}^{1} \frac{\varphi_{i}(x)R(x)dt = 0 \quad R(x) = \sum_{i=1}^{n} \omega_{i} \frac{d^{2}\varphi_{i}(x)}{dx^{2}} + f(x)$$
FEM – Basis Weights
• FEM – Basis Weights
• Galerkin Scheme
Force the residual to be "orthogonal" to the basis functions

$$\int_{0}^{1} \varphi_{i}(x)R(x)dt = 0$$
Generates *n* equations in *n* unknowns

$$\int_{0}^{1} \varphi_{i}(x)\left[\sum_{i=1}^{n} \omega_{i} \frac{d^{2}\varphi_{i}(x)}{dx^{2}} + f(x)\right]dx = 0 \quad l \in \{1, ..., n\}$$
• Linear System

$$\sum_{i=1}^{n} \omega_{i} \int_{0}^{1} \frac{d\varphi_{i}}{dx} \frac{d\varphi_{i}}{dx} dx = \int_{0}^{1} \frac{\varphi_{i}(x)dx}{F_{i}}$$

$$= \int_{1}^{1} \frac{1}{2x} \begin{bmatrix} 2 & -1 \\ -1 & 2 & \ddots \\ & \ddots & \ddots & -1 \\ -1 & 2 \end{bmatrix} \begin{bmatrix} \omega_{i} \\ \omega_{i} \\ \omega_{i} \\ \omega_{i} \end{bmatrix} = \begin{bmatrix} F_{i} \\ F_{i} \\ \vdots \\ F_{n} \end{bmatrix}$$

* Comparing - FD & FEM (hat basis) - 1D problem











0.05

.

0.2 0.3

04 05 X

0.1



389

11

0.9

0.7 0.8

0.6

Use: Taylor Series

$$\frac{u(x_{i+1}) - 2u(x_{i}) + u(x_{i-1})}{\Delta x^{2}} \cong \frac{\partial^{2} u}{\partial x^{2}}$$

$$u(x_{i+1}) = u(x_{i}) + \Delta x \cdot u^{i}(x_{i}) + \frac{\Delta x^{2}}{2} u^{*}(x_{i}) + \frac{\Delta x^{3}}{3!} u^{**}(x_{i}) + \frac{\Delta x^{4}}{4!} u^{(4)}(x) \Big|_{x \in [x_{i-1}, x_{i+1}]}$$

$$-2u(x_{i}) = -2u(x_{i})$$

$$u(x_{i-1}) = u(x_{i}) - \Delta x \cdot u^{i}(x_{i}) + \frac{\Delta x^{2}}{2} u^{*}(x_{i}) - \frac{\Delta x^{3}}{3!} u^{**}(x_{i}) + \frac{\Delta x^{4}}{4!} u^{(4)}(x) \Big|_{x \in [x_{i-1}, x_{i+1}]}$$

$$u(x_{i+1}) - 2u(x_{i}) + u(x_{i-1}) = \Delta x^{2} u^{*}(x_{i}) + \frac{\Delta x^{4}}{12} u^{(4)}(x)$$

$$u(x_{i+1}) - 2u(x_{i}) + u(x_{i-1}) = \Delta x^{2} u^{*}(x_{i}) + \frac{\Delta x^{4}}{12} u^{(4)}(x)$$

$$u(x_{i+1}) - 2u(x_{i}) + u(x_{i-1}) = \Delta x^{2} u^{*}(x_{i}) + \frac{\Delta x^{2}}{12} u^{(4)}(x)$$

$$u(x_{i+1}) - 2u(x_{i}) + u(x_{i-1}) = \Delta x^{2} u^{*}(x_{i}) + \frac{\Delta x^{2}}{12} u^{(4)}(x)$$

$$u(x_{i+1}) - 2u(x_{i}) + u(x_{i-1}) = \Delta x^{2} u^{*}(x_{i-1}) + \frac{\Delta x^{2}}{12} u^{(4)}(x)$$
Substitute back and get τ_{i} L.T.E.

$$-\left[u^{*}(x_{i}) - \frac{\Delta x^{2}}{12} u^{(4)}(x) + \frac{x_{i}(x_{i-1}, x_{i-1})}{12}\right] - f_{i} = \tau_{i} \Rightarrow \tau_{i} = \frac{\Delta x^{2}}{12} u^{(4)}(x) + \frac{x_{i}(x_{i-1}, x_{i-1})}{12}$$
Since $u^{*} = f$

✤ Globa Discretization Error Analysis – Error Equation

Discretization Error Analysis – Error Equation Let $e_j = u(x_j) \cdot \hat{u}_j$ be the discretization error. $-\frac{u(x_{j+1}) - 2u(x_j) + u(x_{j-1})}{\Delta x^2} = f(x_j) + \tau_j$ $-\frac{\hat{u}_{j+1} - 2\hat{u}_j + \hat{u}_{j-1}}{\Delta x^2} = f(x_j)$ Subtracting $-\frac{e_{j+1} - 2e_j + e_{j+1}}{\Delta x^2} = \tau_j$ $1 \le j \le n$ and $e_0 = e_{n+1} = 0$





Ar





 $\|\hat{u}\|_{\infty} = \max_{i} \hat{u}_{i} | = \max_{i} \sum_{j} \alpha_{ij} f_{j} | \leq \max_{i} \left(\sum_{j} \alpha_{ij} \right) \max_{i} |f_{i}| \leq \frac{1}{8} \|\underline{f}\|_{\infty}$ • Convergence Using the discrete stability estimate on $A \underline{e} = \underline{\tau}$ $\left\|e\right\|_{\infty} \leq \frac{1}{8}\left|\tau\right|_{c}$ or $\max_{1 \le i \le n} u(x_i) - \hat{u}_i | \le \frac{\Delta x^2}{96} \max_{0 \le i \le 1} u^{(4)}(x)$ A- priori Error Estimate $Ae = \tau$ $e = \mathbf{A}^{-1} \mathbf{\tau} \qquad \text{Local} \qquad \|\mathbf{\tau}\|_{e} \leq \frac{\Delta x^{2}}{12} u^{(4)}(x)$ Global Error bounded $\|e\|_{\infty} = \max_{i \in [1,n]} |e_i| = \|\mathbf{A}^{-1}\boldsymbol{\tau}\|_{\infty} \leq \underbrace{\|\mathbf{A}^{-1}\|_{\infty}}_{\text{Stability } 1..\text{T.E.}}$ $\left\|\mathbf{A}^{-1}\right\|_{\infty} = \max \frac{\left\|\mathbf{A}^{-1}x\right\|_{\infty}}{\|x\|_{\infty}} = \max_{\|x\|_{\infty} = 0} \left\|\mathbf{A}^{-1}x\right\|_{\infty} = \max_{\text{all rows} i} \sum_{j} |\alpha_{ij}| \qquad \max \text{ row sum}$ $\begin{bmatrix} \pm 1 \\ \pm 1 \\ \vdots \\ \pm 1 \\ \vdots \\ \pm 1 \end{bmatrix}_{H_{v}=1} \quad \mathbf{A} = \frac{1}{\Delta r^{2}} \begin{bmatrix} 2 & -1 & \\ -1 & 2 & \ddots & \\ & \ddots & \ddots & -1 \\ & & -1 & 2 \end{bmatrix}$

We know \mathbf{A}^{-1} has all positive entries

Can look at max column sum since A, A⁻¹ symmetric

What x should we pick to maximize $\|\mathbf{A}^{-1}x\|_{\mathcal{L}}$? $x = \begin{bmatrix} 1 \\ 1 \end{bmatrix}$



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Summary

 For a simple model problem we can produce numerical approximations of

arbitrary accuracy.

- An **a-priori error estimate** gives the asymptotic dependence of the solution error on the discretization size Δx .

FORMAL GENERALIZATION TO ANY PDE

Definitions

Consider a linear elliptic differential equation $\mathcal{L}u = f$ and a difference scheme $\hat{\mathcal{L}}\hat{u} = \hat{f}$

e.g. Different PDE

$$\mathcal{L}u = f - \frac{\partial^2 u}{\partial x^2} + \alpha u(x) = f(x) \quad \text{Helmholtz}$$

$$\hat{\mathcal{L}}\hat{u} = \hat{f} - \frac{\hat{u}_{j-1} - 2\hat{u}_j + \hat{u}_{j+1}}{\Delta x^2} + \alpha \hat{u}_j = f_j \quad \text{F.D. eq on F.D. solution } \hat{u}_j$$

```
    Consistency + Stability yields Convergence

    Consistency

              The difference scheme is consistent with the differential equation if:
              For all smooth functions v
              (\hat{\mathcal{L}}_{\underline{\mathcal{V}}} - \hat{f})_i \to 0, for j = 1, ..., n when \Delta x \to 0.
              (\hat{\mathcal{L}}_{\underline{V}} - \hat{\underline{f}})_{j} \to 0, for j = 1, ..., n when \Delta x \to 0.
              (\hat{\mathcal{L}}_{\underline{\nu}} - \hat{f})_i = \mathcal{O}(\Delta x^p) for all j \implies p is order of accuracy
                <u>Consistency</u> LTE \tau \to 0 if \Delta x \to 0
                    \hat{\mathcal{L}}u - \hat{f} = \tau Exact solution u into F.D. equation
                 -\frac{u(x_{j-1}) - 2u(x_j) + u(x_{j+1})}{\Delta x^2} + \alpha u(x_j) - f(x_j) = \tau_j
In this case note \tau_j = \frac{1}{12} \Delta x^2 u^{(4)}(x)|_{x \in [-,]}
                              As for Poisson Eq., p = 2
     o Truncation Error
              (\hat{\mathcal{L}}\underline{u} - \hat{f})_i = \tau_i, for j = 1, ..., n or, \hat{\mathcal{L}}\underline{u} - \hat{f} = \underline{\tau}.
              The truncation error results from inserting the exact solution into the
              difference scheme.
              Consistency \Rightarrow \|\mathbf{\tau}\|_{\infty} = \mathcal{O}(\Delta x^{p})

    Error Equation

              OriginalScheme \hat{\mathcal{L}}\hat{\underline{u}} = \hat{f}
              Consistency \hat{\mathcal{L}}u = \hat{f} + \tau
              The error e = u - \hat{u} satisfies \hat{\mathcal{L}}e = \tau
                      \hat{f}e = \tau
                                                    e = u - \hat{u}
                   -\frac{u(x_{j-1})-2u(x_j)+u(x_{j+1})}{\Delta x^2}+\alpha u(x_j)-f(x_j)=\tau_j Exact Solution into F.D. Equation
                   \frac{-\frac{\hat{u}_{j-1} - 2\hat{u}_j + \hat{u}_{j+1}}{\Delta x^2} + \alpha \hat{u}_j - f_j = 0}{-\frac{e_{j-1} - 2e_j + e_{j+1}}{\Delta x^2} + \alpha e_j = \tau_j}
                                                                                                         F.D. Solution û
                                                                                                       into F.D. Equation
                                                                                                            Subtract
```

• Stability

$$\begin{aligned} \mathbf{Matrix Norm} \quad \|\mathbf{M}\|_{\infty} = \sup_{v \in \mathbb{R}^{d}} \frac{\|\mathbf{M}v\|_{\infty}}{|\mathbf{y}|_{\infty}} \\
\text{The difference scheme is stable if } \quad \|\hat{\mathcal{L}}^{-1}\|_{\infty} \leq C \quad (\text{independent of } \Delta \mathbf{x}) \\
\|\hat{\mathcal{L}}_{\text{Helm}} \hat{u} = \left(\alpha \begin{bmatrix} 1 & & \\ 1 & \ddots & \\ & & 1 \end{bmatrix} + \frac{1}{(\Delta \mathbf{x})^{2}} \begin{bmatrix} 2 & -1 & & \\ -1 & 2 & \ddots & \\ & \ddots & \ddots & -1 \\ & & -1 & 2 \end{bmatrix} \begin{bmatrix} \hat{u}_{1} \\ \hat{u}_{1} \\ \vdots \\ \hat{u}_{n} \end{bmatrix} \\
\end{aligned}$$

$$\begin{aligned} \hat{f}^{e} = \tau \\ e = \left[\hat{f}^{-1} \right]_{\Sigma} \\
\hat{f} = \begin{bmatrix} 2 + * & -1 & & \\ -1 & 2 + * & \ddots & \\ & \ddots & \ddots & -1 \\ & & -1 & 2 + * \end{bmatrix} \| \mathbf{A}^{-1} \|_{\infty} \leq \frac{1}{8} \\
\| \mathbf{M} \|_{\infty} = \sup_{\|v\|_{\infty} = 1} \| \mathbf{M} v \|_{\infty} = \sup_{\|v\|_{\infty} = 1} \left(\max_{j \in \mathbb{N}} \sum_{j=1}^{n} m_{ij} v_{j} \right) \\
& = \max_{i} \left(\sup_{\|v\|_{\infty} = 1} \left| \sum_{j=1}^{n} m_{ij} v_{j} \right| \right) \\
& = \max_{i} \left(\sup_{\|v\|_{\infty} = 1} \left| \sum_{j=1}^{n} m_{ij} v_{j} \right| \right) \\
& = \max_{i} \left(\sup_{\|v\|_{\infty} = 1} \left| \sum_{j=1}^{n} m_{ij} v_{j} \right| \right) \\
& = \max_{i} \left(\sup_{\|v\|_{\infty} = 1} \left| \sum_{j=1}^{n} m_{ij} v_{j} \right| \right) \\
& \text{Convergence} \\
& \text{Error Equation} \quad e = \hat{\mathcal{E}}^{-1} \tau \\
& \text{Taking Norms} \quad \| \mathcal{E} \|_{\infty} = \| \hat{\mathcal{E}}^{-1} \mathbf{x} \|_{\infty} \leq \| \hat{\mathcal{E}}^{-1} \|_{\infty} \| \mathbf{x} \|_{\infty} \leq \| \hat{\mathcal{E}}^{-1} \|_{\infty} C \Delta x^{\mu} = C_{1} \Delta x^{\mu} \end{aligned}$$



INTRODUCTION TO NUMERICAL SIMULATION

LECTURE 19.

Boundary Value Problems - Solving 3-D Finite-Difference Problems

TODAY'S OUTLINE:

- Finite Difference Matrices in 1-D, 2-D, and 3-D
- Gaussian Elimination Cost:
- Bandlimited GE
- Sparse GE
- Krylov Method Cost
 - Counting iterations: Communication Lower Bound

FINITE DIFFERENCE MATRICES IN 1-D, 2-D, AND 3-D Structural Analysis of Automobiles



- Equations
 - Force-displacement relationships for mechanical elements (beams, plates, shells) and sum of forces = 0.
- Partial Differential Equations of Continuum Mechanics
- Drag Force Analysis of Aircraft



- Equations
 - Navier-Stokes Partial Differential Equations.



Equations

 The Poisson Partial Differential Equation
 FD Matrix Properties
 2-D Discretized Problem

Discretized Poisson




















o 3-D Discretized Problem















GAUSSIAN ELIMINATION COST
* Bandlimited GE
• Dense GE
- Picture

$$\begin{bmatrix} A_{11} & A_{12} & A_{13} & A_{14} \\ 0 & A_{22} & A_{23} & A_{24} \\ 0 & 0 & A_{33} & A_{34} \\ 0 & 0 & 0 & A_{43} \end{bmatrix}$$
- Algorithm
For *i* = 1 to *n* - 1 {
For *j* = *i* + 1 to *n* {
For *i* = 1 to *n* - 1 {
For *i* = *i* to *n* - 1 {
For *i* = *i* to *n* + 1 {
For *i* = *i* to *n* + 1 {
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Complexity of Sparse GE

	matrix size	dense block			
1 - D	n = m	d = 1	da	nain aina	
2 - D	$n = m^2$	d = m	do	main size =	m
3 – D	$n = m^3$	$d = m^2$			
1-D O()	$a) + O(d^{\hat{e}}) = O(m)$	← 100 pt grid		O(100) ops	1 µsec
2-D O()	$n) + O(d^2) = O(m^3)$	$(\leftarrow 100 \times 100 \text{ grid})$		$O(10^{6})$ ops	10 msec
3-D O()	$a) + O(d^{\tilde{e}}) = O(m^6)$	$(\rightarrow 100 \times 100 \times 100)$	grid	$O(10^{12})$ ops	2h 46m
e.g. on a For 3-D	1 GFlops comput problems – can w	ter using $m = 100$ and we do any better with	nd a co GCR	constant $c = 10$	

KRYLOV METHO The Generaliant O Algorit	DD COST zed Conjugate Resid hm Complexity	dual A	Algorithm			
\widetilde{n}	Compl $\leftarrow r^{k-1} - \checkmark^{k-1} (\mathbf{A} p)$	ute A	r^{k-1} Span proc	rse Mat luct cos	rix-vector sts $O(n)$	
P_{k-1} P_{k-1}	$\leftarrow \frac{\widetilde{p}_{k-1}}{\ \mathbf{A}\widetilde{p}_{k-1}\ }$,) (inne	r prod	uct : $O(n$)
\mathcal{Y}_{k-1}	$\leftarrow \begin{pmatrix} r^{k-1} \end{pmatrix}^T \mathbf{A} p_{k-1} \\ \leftarrow \begin{pmatrix} r^{k-1} \end{pmatrix}^T \mathbf{A} p_{k-1} \end{pmatrix}$		inne	r prod	uct : $O(n$)
$x \leftarrow r^k \leftarrow$	$-x + y_{k-1}p_{k-1}$ $-r^{k-1} - y_{k-1}\mathbf{A}p_{k-1}$	-1	mult	tiplica	tions: O	(n)
o Compl	ithm is O(kn) where exity of sparse GCF matrix size we	k = 1 R (genorst call	number of iteration neral worst case ite ase iterations ($k = i$	rations n))	
1 2 -	$D \qquad n = m$ $D \qquad n = m^2$		$k = m$ $k = m^2$	de	omain size =	= <i>m</i>
3 – 1-D	$D \qquad n = m^3 \\ O(kn) = O(m^2)$	←	$k = m^3$ 100 pt grid	0	(10^4) ops	0.1 ms
2-D 3-D e.g.	$O(kn) = O(m^*)$ $O(kn) = O(m^6)$ on a 1 GFlops com	← ← puter	100×100 grid $100 \times 100 \times 100$ g using $m = 100$ and	orid O a cons	$(10^{\circ}) \text{ ops}$ $(10^{12}) \text{ ops}$ tant $c = 10$	l sec 2h 46m
But	how many iteration	s k de	bes it really take?			







Communication Lower Bound



 $\mathbf{A}^{k} r^{0}$ is nonzero in m^{th} entry after k = m iterations

Need at least *m* iterations for $(x^{k+1})_m = \left(x^0 + \sum_{j=1}^k \alpha_j r^j\right)_m \neq 0$

o 2-D Case



	$k = \frac{\log \left(\frac{1}{\sqrt{2}}\right)}{\log \left(\frac{1}{\sqrt{2}}\right)}$	$\frac{g^{\frac{\gamma}{2}}}{\kappa-1} \underset{\kappa+1}{\cong} O(m)$			
¢	GCR achie Complexity of spa matrix s	eves communication arse GCR (no precon- ize iterations ($k =$	lower bound O(n nd) for 3D FD me n)	m)! ethods	
	$\begin{array}{ll} 1-D & n=m\\ 2-D & n=m \end{array}$	k = m $k = m$	domain size	c = m	
	$\begin{array}{ccc} 3-D & n=m\\ 1-D & C\\ 2-D & C\\ 3-D & C\\ e.g. \text{ on a J GFlo}\\ But how many i \end{array}$	$\begin{array}{l} 3 \qquad k = m \\ O(kn) = O(m^2) \qquad \leftarrow \\ O(kn) = O(m^3) \qquad \leftarrow \\ O(kn) = O(m^4) \qquad \leftarrow \\ ps \text{ computer using } n \\ terations k \text{ does it respectively} \end{array}$	100 pt grid 100×100 grid $100 \times 100 \times 100$ n = 100 and a con- cally take?	$O(10^4)$ ops $O(10^6)$ ops grid $O(10^8)$ ops istant $c = 10$	0.1 ms 10 ms 1 sec
*	Work for Banded Dimension 1 2 3 GCR faster than Could be fas Must defeat	Gaussian Eliminatio Banded GE O(m) $O(m^4)$ $O(m^7)$ banded in GE in 2 ter, 3-D matrix only the communication	on, Relaxation and Sparse GE O(m) $O(m^3)$ $O(m^6)$ and 3 dimensions m^3 nonzeros. lower bound!	d GCR GCR $O(m^2)$ $O(m^3)$ $O(m^4)$	
	$m = 100 \ 3-D$	$SPG - O(10^{12}) o$	ps ← 10,000 sec		
		GCR – O(10 ⁸) o	ps ← 1 sec		
	Dimension 1 2 3	Banded GE 1 µs 1 sec 11days 6h	Sparse GE 1 µs 10 ms 2h 46m	GCR (no precon 0.1 ms 10 ms 1 sec	ditioner)
	e.g. on a 1 GFlo GCR faster t Could be fas Must defeat	ps computer using <i>n</i> han banded GE in 2 ter, 3-D matrix only the communication	n = 100, and cons and 3 dimension m^3 nonzeros. lower bound!	tant c = 10	
	Dimension 1 2 3	Banded GE 10 μs 2h 46m 317.000 years	Sparse GE 10 µs 10 sec 317 years	GCR (no precon- 10 ms 10 sec 2h 46m	ditioner)
	e.g. on a 1 GFlo	ps computer using h	n = 1000, and cor	stant c = 10	

Dimension	Banded GE	Sparse GE	GCR	GCR + preconditioner				
1	O(m)	O(m)	$O(m^2)$	O(km)				
2	$O(m^4)$	$O(m^3)$	$O(m^3)$	$O(km^2)$				
3	$O(m^7)$	$O(m^6)$	$O(m^4)$	$O(km^3)$				
GCR+good	preconditioner (i	.e. $k = 5 - 10$) is	always the	best choice for 2D & 3D				
problems!								
Dimension	Banded GE	Sparse GE	GCR	GCR + preconditioner				
1	10 µs	10 µs	10 ms	0.1 ms				
2	2h 46m	10 sec	10 sec	0.1 sec				
3	317,000 years	317 years	2h 46m	1m 20s				
e.g. on a 1 GFlops computer using $\underline{m} = 1000$, $k=10$, and constant $c=10$ Cube Example, Interior Problem – Memory Usage								
1	Cube examp	le, Interior	Problem	- Memory usage				
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INTRODUCTION TO NUMERICAL SIMULATION

LECTURE 21.

Finite Element Methods & GCR Preconditioners

TODAY'S OUTLINE:

- * Finite Element Methods
 - The FEM Basis Functions
 - The Galerkin scheme
 - The FEM linear system
 - An FEM Example
- Energy minimization view of FEM
- GCR Preconditioners for FD and FEM
 - **Diagonal Preconditioners**
 - Block Diagonal Preconditioners _
 - Incomplete Factorization Preconditioners
 - -Communication Improving Preconditioners: Gauss-Seidel
 - Examples of GCR + Preconditioners for 3D FD

Finite Element Methods





Question: What is the temperature distribution along the bar





$u(x) = \sum_{j=1}^{n} w_j \varphi_j(x)$ u(0) = 0 u(1) = 0

What are good basis functions?

- 1.) $\varphi(x) = a_0 + a_1 x \rightarrow \text{No, second derivative is always zero}$
- 2.) $\varphi(x) = a_0 + a_1 x + a_2 x^2 \rightarrow \text{Could work}$
- 3.) I can try to build into the basis function the boundary conditions e.g. $\varphi_i(x) = \sin 2\pi i x$ [and second derivative is nonzero] Good choice. This is the same used by Spectral Methods. Fourier Series.

$$\varphi_1(x)^{\dagger}$$
 $\varphi_2(x)^{\dagger}$ $\varphi_2(x)^{\dagger}$ $\varphi_2(x)^{\dagger}$

4.) Piecewise Constant. Good. This leads to Finite Difference Methods

$$u(x) = \begin{pmatrix} u(x) \\ h & 2h & 3h & 4h \end{pmatrix}$$



Step 2: Generate equations for the basis functions weights setting residual orthogonal to some test functions

conditioners

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Assume the initial condition is satisifed $\sum_{i=1}^{n} \omega_i \varphi_i(0) = x_0$ Define the Residual Define the Residual $\begin{aligned} R(t) &= \sum_{i=1}^{n} \omega_i \frac{d\varphi_i(t)}{dt} - A \sum_{i=1}^{n} \omega_i \varphi_i(t) - u(t) \end{aligned}$ Select weights to "minimize" the residual Force the residual to be zero at *n* test points $R(t_i) &= \sum_{i=1}^{n} \omega_i \frac{d\varphi_i(t_i)}{dt} - A \sum_{i=1}^{n} \omega_i \varphi_i(t_i) - u(t_i) = 0 \qquad l \in \{1, ..., n\}$ Generates *n* equations in *n* unknowns

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THE GALERKIN SCHEME

Force the residual to be "orthogonal" to the basis functions

 $\int \varphi_I(x) R(x) dt = 0$









THE FEM LINEAR SYSTEM













 Properties of A_h (2 -1) $\begin{pmatrix} 2 & -1 \\ -1 & 2 & -1 \\ & -1 & 2 & -2 \\ & & \ddots \end{pmatrix}$ $\underline{A}_{h}^{1} = \frac{1}{h}$ -1 2 -1 -1 2 \underline{A}_{h}^{1} is SPD [Symmetric Positive Definite]; and diagonally dominant; and sparse; and tridiagonal. ✤ "Load" Vector Elements: F_h $\begin{array}{c} \varphi_{i} \\ \hline T_{h}^{i} + T_{h}^{i+1} \\ \hline T_{h}^{i} - T_{h}^{i+1} \end{array} \qquad F_{hi} = \int_{0}^{1} f \varphi_{i} dx \\ F_{hi} = \int_{T_{k}^{i}}^{1} f \varphi_{i} dx + \int_{T_{k}^{i+1}}^{1} f \varphi_{i} dx, \quad i = 1, \dots, n; \end{array}$

COMPARING FEM AND FD IN 1-D

 Residual Equation Partial Differential Equation Form Partial Direction $x_{1}^{2} = f(u(0) = 0)$ $-\frac{\partial^{2} u}{\partial x^{2}} = f(u(0) = 0)$ Step 1: Choose Basis Functions to represent the solution $(x_{1}^{2}) = \sum_{i=1}^{n} \alpha_{i} = \omega_{i}(x)$

$$u(x) \approx u_h(x) = \sum_{i=1}^{\infty} \omega_i \underbrace{\varphi_i(x)}_{\text{Basis functions}}$$

Step 2: Generate equations for the basis functions weights setting residual orthogonal to some test functions

$$\int_{0}^{1} \frac{\phi_{1}(x)}{\tau_{\text{rest}}} R(x) dt = 0 \quad R(x) = \sum_{i=1}^{n} \omega_{i} \frac{d^{2} \varphi_{i}(x)}{dx^{2}} + f(x)$$

Functions

FEM – Basis Weights

o Galerkin Scheme Force the residual to be "orthogonal" to the basis functions

$$\int_{0}^{1} \varphi_{I}(x)R(x)dt = 0$$
Generates *n* equations in *n* unknowns
$$\int_{0}^{1} \varphi_{I}(x) \left[\sum_{i=1}^{n} \omega_{i} \frac{d^{2} \varphi_{I}(x)}{dx^{2}} + f(x) \right] dx = 0 \quad I \in \{1, ..., n\}$$
• Linear System
$$\sum_{i=1}^{n} \omega_{i} \int_{0}^{1} \frac{d\varphi_{I}}{dx} \frac{d\varphi_{I}}{dx} dx = \int_{0}^{1} \varphi_{I}f(x)dx$$

$$\frac{1}{\Delta x} \begin{bmatrix} 2 & -1 \\ -1 & 2 & \ddots \\ & \ddots & \ddots & -1 \\ -1 & 2 \end{bmatrix} \begin{bmatrix} \omega_{1} \\ \omega_{2} \\ \vdots \\ \omega_{n} \end{bmatrix} = \begin{bmatrix} F_{1} \\ F_{2} \\ \vdots \\ F_{n} \end{bmatrix}$$
• Comparing - FD & FEM (hat basis) - 1D problem
$$FEM = \frac{1}{\Delta x} \begin{bmatrix} 2 & -1 \\ -1 & 2 & \ddots \\ & \ddots & \ddots & -1 \\ -1 & 2 \end{bmatrix} \begin{bmatrix} \omega_{1} \\ \omega_{2} \\ \vdots \\ \omega_{n} \end{bmatrix} = \begin{bmatrix} \int_{0}^{1} \varphi_{I}(x)f(x)dx \\ \int_{0}^{1} \varphi_{2}(x)f(x)dx \\ \vdots \\ \int_{0}^{1} \varphi_{n}(x)f(x)dx \\ \vdots \\ \int_{0}^{1} \varphi_{n}(x)f(x)dx \end{bmatrix}$$
FD
$$\frac{1}{\Delta x} \begin{bmatrix} 2 & -1 \\ -1 & 2 & \ddots \\ & \ddots & \ddots & -1 \\ -1 & 2 \end{bmatrix} \begin{bmatrix} \hat{u}_{1} \\ \hat{u}_{2} \\ \vdots \\ \hat{u}_{n} \end{bmatrix} = \begin{bmatrix} f(x_{1}) \\ f(x_{2}) \\ \vdots \\ f(x_{n}) \end{bmatrix}$$



F.D.





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solution vector $\begin{bmatrix} \hat{u}_1 \\ \hat{u}_2 \end{bmatrix}$

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Top View





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* Review of FEM
      Introduce basis representation u(x) \approx u_h(x) = \sum_{i=1}^{n} \omega_i \phi_i(x)
                                                                                        Basis
                                                                                      Functions
       \Rightarrow u_h(x) is a weighted sum of basis functions
      The basis functions define a space
       X_h = \langle v \in X_h | v = \sum \beta_i \varphi_i \text{ for some } \beta_i s \rangle
                                  i=1
      Example
                                                             Piecewise linear Space
              "Hat" basis functions
                    \phi_1 = \phi_5 = \phi_5
      Partial Differential Equation Form
      -\frac{\partial^2 u}{\partial x^2} = f \quad u(0) = 0 \quad u(1) = 0
"Nearly" Equivalent Weak Form
           \int_{\Omega} \frac{\partial u}{\partial x} \frac{\partial v}{\partial x} dx = \int_{\Omega} f v dx \quad \text{for all } v
                                Ω
                a(u,v)
                                   l(v)
       Introduced an abstract notation for the equation u must satisfy
            a(u,v) = l(v) for all v
       a(u,u) defines a norm a(u,u) = \|u\|^2
       u is restricted to be 0 at 0 and 1!!!
       Key Idea:
       Using the norm properties, it is possible to show
            If a(u_h, \varphi_i) = l(\varphi_i) for all \varphi_i \in \{\varphi_1, \varphi_2, \dots, \varphi_n\}
                 Note: this is basically Galerkin
           Then \underbrace{\|u - u_h\|}_{\text{Solution}} = \min_{w_k \in X_h} \underbrace{\|u - w_h\|}_{\text{Projection}}
                          Error
                                                    Error
            Or in other words: given some basics, Galerkin finds the best coefficients
```



 Error Analysis – Energy Norm In words: even if you knew u, you could not find a w_h in X_h more accurate than u_h in the energy norm. FEM Convergence Analysis Choosing the basis functions space X_h specifies the error: $|||e|||=|||u-u_h|||$ U. Question: How does that error change when I change basis? For piecewise linear: $|||u - u_h||| = O\left(\frac{1}{n}\right)$ ENERGY MINIMIZATION VIEW OF FEM Problem of Interest Helmholtz Equation in 1D Boundary Value Problem (BVP) - Strong Form $-u''(x) + \alpha u(x) = f(x) \qquad \alpha \ge 0$ $x \in (0,1)$ u(0) = u(1) = 0Describes many physical phenomena (e.g.): - Temperature distribution in a bar - Deformation of an elastic bar - Deformation of a string under tension Solution Properties - The solution u(x) always exists u(x) is always smoother than the data f(x)given f(x) the solution u(x) is **unique**



In words: Over all functions w in X, u that satisfies $-u_{xx} + \alpha u = f$ in Ω u(0) = u(1) = 0makes J(w) as small as possible o Proof Let w = u + vThen $J\left(\frac{\underset{u^{\chi}}{\underline{u}}}{\underset{u^{\chi}}{\underline{v}}}\right) = \frac{1}{2}\int_{0}^{1} (u+v)_{x} (u+v)_{x} dx + \frac{\alpha}{2}\int_{0}^{1} (u+v)(u+v) dx - \int_{0}^{1} f(u+v) dx$ $J(u+v) = \frac{1}{2} \int_{0}^{1} (u_x u_x + \alpha u u) dx - \int_{0}^{1} f u dx \quad J(u)$ $+ \int_{0}^{2} (u_x v_x + \alpha u v) dx - \int_{0}^{0} f v dx \quad \delta J_v(u)$ first variation $+\frac{1}{2}\int_{0}^{1} (v_{x}v_{x} + \alpha vv) dx \qquad > 0 \text{ for } v \neq 0$ $\delta J_{\text{Poisson}}(w) = \int_{0}^{1} u_x v_x dx - \int_{0}^{1} f v dx$ If $v_x = \varphi_i$ then Galerkin $\delta J_v(u) = \int (u_x v_x + \alpha u v) dx - \int f v dx$ $= \overline{v(0)} u_{x}(0) - \overline{v(1)} u_{x}(1) - \int_{0}^{1} u_{xy} v dx + \alpha \int_{0}^{1} u v dx - \int_{0}^{1} f v dx$ $J(\underbrace{u+v}_{u}) = J(u) + \frac{1}{2} \int_{0}^{1} (v_{v}v_{x} + \alpha vv) dx, \quad \forall v \in X$ $\frac{>0 \text{ unless } v=0}{J(w) > J(u), \quad \forall w \in X, w \neq u}$ 1 \Rightarrow u is **THE**minimizer of J(w)



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$$u_{h}(x) = \sum w_{i}\varphi_{i}(x)$$

$$-u_{xx} + \alpha u - f(x) = 0$$

$$-\frac{\partial^{2}}{\partial x^{2}} (\sum w_{i}\varphi_{i}(x)) + \alpha \sum w_{i}\varphi_{i}(x) - f(x) \cong 0$$

$$R(x)$$

$$\int_{0}^{1} \varphi_{i}(x) \sum_{i=1}^{n} w_{i} \frac{\partial^{2}}{\partial x^{2}} \varphi_{i}(x) dx + \int_{0}^{1} \varphi_{i}(x) \alpha \sum_{i=1}^{n} w_{i}\varphi_{i}(x) dx - \int_{0}^{1} \varphi_{i}(x) f(x) = 0$$

$$J^{R}(\underline{w} \in \Re^{n}) \equiv J\left(\sum_{j=1}^{n} w_{j}\varphi_{j}\right)$$

$$= -\frac{1}{2} \underline{w}^{T} \underline{A}_{h} \underline{w} - \underline{w}^{T} \underline{F}_{h}$$

$$\underline{F}_{h} \in \Re^{n} : F_{hi} = \int_{0}^{1} f\varphi_{i} dx$$

$$\underline{A}_{h} \in \Re^{n \times n} : A_{hij} = \int_{0}^{1} \left(\frac{d\varphi_{j}}{dx} \frac{d\varphi_{j}}{dx} + \alpha \varphi_{i}\varphi_{j}\right) dx$$



$$J^{R}(u_{h} + v) = \frac{1}{2}(u_{h} + v)^{T} A_{h}(u_{h} + v) - (u_{h} + v)^{T} F_{h}$$

$$= \frac{1}{2}u_{h}^{T} A_{h}u_{h} - u_{h}^{T} E_{h} + \frac{1}{2}v^{T} A_{h}u_{h} + \frac{1}{2}u_{h}^{T} A_{h}v - v^{T} E_{h} + \frac{1}{2}v^{T} A_{h}v$$

$$J^{R}(u_{h} + v) = J(u)$$

$$+ \frac{(A_{h}u_{h} - E_{h})^{T}}{v^{R}(u_{h})} SPD$$

$$+ \frac{1}{2}v^{T} A_{h}v$$

$$SPD$$
If (and only if)
$$\delta J^{R}_{v}(u_{h}) = 0, \quad \forall v \in \Re^{n}$$

$$\int (\nabla J^{R}(u_{h}) = A_{h}u_{h} - E_{h} = 0$$
then
$$J(w = u_{h} + v) > J(u_{h}), \quad \forall v \neq 0$$
Poisson $\rightarrow \min J(u)$

$$u = \sum w_{i}\varphi_{i}$$
Find $u_{h} \in \Re^{n}$ such that
$$A_{h}u_{h} = E_{h} \implies u_{h}(x) = \sum_{j=1}^{n} u_{hj}\varphi_{j}(x)$$
SPD \Rightarrow existence and uniqueness
$$Firror Analysis - Energy Norm
$$J(u_{h}) < J(w_{h}), \quad \forall w_{h} \in X_{h}, \quad w_{h} \neq u_{h}$$
If $e = u - u_{h}$

$$\|u - u_{h}\|$$
where the energy norm is defined as:
$$\|e\| = \sqrt{\int_{0}^{1} (e_{x}e_{x} + \alpha ee) dx$$$$



GCR Preconditioners



DIAGONAL PRECONDITIONERS Let $\mathbf{A} = \mathbf{D} + \mathbf{A}_{nd}$ Apply GCR to $(\mathbf{D}^{-1}\mathbf{A})\mathbf{\vec{x}} = (\mathbf{I} + \mathbf{D}^{-1}\mathbf{A}_{nd})\mathbf{\vec{x}} = \mathbf{D}^{-1}\mathbf{\vec{b}}$ The inverse of a diagonal is cheap to compute Usually improves convergence Krylov Methods - Convergence Analysis GCR Optimality Property: $\left\|\vec{r}^{k}\right\| = \left\|\varphi_{k}(\mathbf{M})\vec{r}^{0}\right\| \le \left\|\widetilde{\varphi}_{k}(\mathbf{M})\vec{r}^{0}\right\| \le \left\|\widetilde{\varphi}_{k}(\mathbf{M})\right\| \left\|\vec{r}^{0}\right\|$ $\widetilde{\wp}_{k}$ is any k^{th} order poly such that $\widetilde{\wp}_{k}(0) = \mathbf{I}$ Therefore Any polynomial which satisfies the constraint can be used to get an upper bound on $\left\| \stackrel{\|'}{_{\vec{F}}^{0}} \right\| \leq \left\| \widetilde{\varphi}_{k} \left(\mathbf{M} \right) \right\|$ Low Order Poly lustered Eigenvalu



BLOCK DIAGONAL PRECONDITIONERS



Tridiagonal Matrices factor quickly



Throw away fill-ins produced by other fill-ins Throw away fill-ins produced by fill-ins of other fill-ins, etc.



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 ${\cal L}_{ij}$

















 Faster Converging GCR GCR already achieves Communication Lower bound <u>Preconditioner must accelerate communication</u> Multiplying by Preconditioner must move values by more than one grid point.

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$$\underbrace{(\mathbf{L} + \mathbf{D})^{-1} \mathbf{A}}_{\mathbf{M}} x = (\mathbf{L} + \mathbf{D})^{-1} b$$

 $(\mathbf{U} + \mathbf{D})^{-1} \mathbf{D} (\mathbf{L} + \mathbf{D})^{-1} \mathbf{A} x = (\mathbf{U} + \mathbf{D})^{-1} \mathbf{D} (\mathbf{L} + \mathbf{D})^{-1} b$



$$\begin{array}{c} u_{0} \ \hat{u}_{1}^{(\mathrm{new})} & \hat{e}^{_{n^{-}}} \\ & \hat{u}_{1}^{(\mathrm{new})} & \hat{u}_{2}^{(\mathrm{new})} & \hat{e}^{_{n^{-}}} \\ & & \hat{u}_{n-1}^{(\mathrm{new})} & \hat{u}_{n}^{(\mathrm{new})} & \hat{e}^{_{n^{-}}} \\ & \hat{u}_{n-2}^{(\mathrm{new})} & \hat{u}_{n-1}^{(\mathrm{new})} & \hat{e}^{_{n^{-}}} \\ & \hat{u}_{n-2}^{(\mathrm{new})} & \hat{u}_{n-1}^{(\mathrm{new})} & \hat{e}^{_{n^{-}}} \end{array}$$

Derivation of the SGS Iteration Equations Forward Sweep (half step): $(\mathbf{D} + \mathbf{L})x^{k+\frac{1}{2}} + \mathbf{U}x^k = b$ Backward Sweep (half step): $(\mathbf{D} + \mathbf{U})x^k + \mathbf{L}x^{k+\frac{1}{2}} = b$ $\Rightarrow x^{k+1} = (\mathbf{D} + \mathbf{U})^{-1} \mathbf{L} (\mathbf{D} + \mathbf{L})^{-1} \mathbf{U} x^k + (\mathbf{D} + \mathbf{U})^{-1} b \dots$ - $(\mathbf{D} + \mathbf{U})^{-1} \mathbf{L} (\mathbf{D} + \mathbf{L})^{-1} b$ $\Rightarrow x^{k+1} = x^k - (\mathbf{D} + \mathbf{U})^{-1} \mathbf{D} (\mathbf{D} + \mathbf{L})^{-1} \mathbf{A} x^k + (\mathbf{D} + \mathbf{U})^{-1} \mathbf{D} (\mathbf{D} + \mathbf{L})^{-1} b$



EXAMPLES OF GCR + PRECONDITIONERS FOR 3D FD

	compter	matrix size	dense b	lock			
	1 - D	n = m	k = 5 -	-10	domoin sins		
	2 - D	$n = m^2$	k = 5 -	-10	domain size	c = m	
	3 - D	$n = m^3$	k = 5 -	-10			
	1-D	O(kn) = 10 O(kn)	(<i>m</i>) ←	100 pt	grid	$O(10^3)$ ops	1 µs
	2-D	O(kn) = 10 O(kn)	$(m^2) \leftarrow$	100×1	00 grid	$O(10^5)$ ops	0.1 ms
	3-D	O(kn) = 10 O	$(m^3) \leftarrow$	100×1	00×100 grid	$O(10^7)$ ops	10 ms
	e.g. on	a 1 GFlops co	mputer us	ing m =	100 and a cor	stant $c = 10$	
÷	Krylov -	Work for Bar	nded Gaus	sian Elin	nination, Rela	ixation, and GC	R
	Dimen	sion Bande	d GE	Sparse G	E GCR	GCR + pre	conditioner
	1	<i>O</i> ()	m)	O(m)	$O(m^2)$	O(k	m)
	2	O(n	n^4)	$O(m^3)$	$O(m^3)$	O(k	m^2)
	3	O(n	n^7)	$O(m^6)$	$O(m^4)$	O(k	m^3)
	GCR+	good precondi	tioner is a	lways the	best choice	for 2D & 3D pro	oblems!

Dimension	Banded GE	Sparse GE	GCR	GCR + preconditioner
1	1 µs	1 µs	0.1 ms	1 µs
2	1 sec	10 ms	10 ms	0.1 sec
3	11d 6h	2h 46m	1 sec	10 ms
e.g. on a 1 G	Flops computer	using $\underline{m} = 1000$), <i>k</i> =10, and	constant $c=10$
GCR+good J	preconditioner is	s always the bes	st choice for	2D & 3D problems!
Dimension	Banded GE	Sparse GE	GCR	GCR + preconditioner
1	10 µs	10 µs	10 ms	0.1 ms

1	10 µs	10 µs	10 ms	0.1 ms
2	2h 46m	10 sec	10 sec	0.1 sec
3	317,000 years	317 years	2h 46m	1m 20s
e.g. on a	1 GFlops computer i	using $\underline{m} = 100$	0, k=10, and con	nstant c=10
GCR+goo	od preconditioner (i.	e. $k = 5 - 10$) is	always the bes	t choice for 2D & 3D
problems	!			

* Cube Example, Interior Problem - Memory Usage



INTRODUCTION TO NUMERICAL SIMULATION

Lecture 22.

PDE - BEM Integral Equation Method I

TODAY'S OUTLINE:

- Poisson and Laplace Equations: examples
 Exterior Versus Interior Problems
- Green's Function
- Basis Functions
- * Collocation Method
- * Galerkin Method

OVERVIEW OF COURSE



MOR: $\nabla^{2}\Psi + \kappa^{2}\Psi = -f \xrightarrow{\text{discretize}} \frac{dx}{dt} = \mathbf{A}x(t) + bu(t) \xrightarrow{\text{MOR}} \frac{d\hat{x}}{dt} = \hat{\mathbf{A}}\hat{x}(t) + \hat{b}\hat{u}(t)$ ODE



3.21PDE - BEM Integral Equation Method H



 Electrostatic Analysis The Laplace Partial Differential Equation. Electrostatic Analysis $\Psi(x) \leftrightarrow$ electrostatic potential $\nabla \Psi(x) \propto$ electric field



Drag Force Analysis of Aircraft





 $\frac{Potential Flow}{\nabla \Psi = air velocity}$

EXTERIOR VERSUS INTERIOR PROBLEMS



Finite Difference/Finite Element Method Meshing Problem



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Exterior Problem in Electrostatics





- High viscosity

- Size of structure







Capacitance on a Microprocessor Signal Line



What is common about these problems

- Exterior Problems

 - Drag Force in MEMs device fluid (air) creates drag.
 - Coupling in a Package Field sin exterior create coupling
- Capacitance of a Signal Line Fields in exterior.
- Quantities of Interest are on the Surface
 - MEMs device just want surface traction force
 - Package just want coupling between conductors
- Signal Line just want surface charge
 Exterior Problem is Linear and Space-Invariant

 - MEMs Exterior Stokes Flow equation (linear)
 - Package Maxwell's equation in free space (linear)
 - Signal Line Laplace's equation in free space (linear)
- But problems are geometrically very complex!!

Exterior Problems

* Why not use Finite-Difference or FEM methods? 2D Heat Flow Example



GREEN'S FUNCTION

Laplace's Equation

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 Green's Function Definition: Green's Function G(x,x') is the solution at test location x produced by a unit point source at location x' e.g. for Laplace operator: $\nabla^2 G(x, x') = \delta(x - x')$ o 2–D
$$\begin{split} \Psi(x,y) &= \log \sqrt{(x-x_0)^2 + (y-y_0)^2} & G(x,x') \\ \frac{\partial^2 \Psi}{\partial x^2} + \frac{\partial^2 \Psi}{\partial y^2} &= 0 \quad \text{for all } (x,y) \neq (x_0,y_0) \\ G(x,x') &= \frac{1}{2\pi} \log |x-x'| & \delta(x-x') \\ x' &= \frac{1}{\sqrt{(x-x_0)^2 + (y-y_0)^2 + (z-z_0)^2}} \\ \frac{\partial^2 \Psi}{\partial x^2} + \frac{\partial^2 \Psi}{\partial y^2} + \frac{\partial^2 \Psi}{\partial z^2} &= 0 \quad \text{for all } (x,y,z) \neq (x_0,y_0,z_0) \\ G(x,x') &= -\frac{1}{4\pi |x-x'|} \end{split}$$

Proof: Just differentiate and see!



















How do we solve the integral equation?











Basis Functions for a Circle – unwrapped circle geometry



 $\mathcal{Q} = \int_{l} \sigma(x) dx = \int_{i=1}^{n} \omega_{i} \phi_{i}(x)$ $\Psi(x) = \oint_{l} G(x, x') \sum_{i=1}^{n} \omega_{i} \phi_{i}(x') dx'$

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 $\Psi(x)_{\text{boundary}} = \int_{\text{boundary}} G(x, x') \sigma(x') dx'$ $= \int_{\substack{\text{approximate} \\ \text{boundary}}} G(x, x') \sum_{j=1}^{n} \omega_{j} \varphi_{j}(x) dx'$ $=\sum_{j=1}^{n}\omega_{j}\int_{\text{boundary}}G(x,x')\varphi_{j}(x')dx'$ $=\sum_{j=1}^{n}\omega_{j}\int_{\text{line }l_{j}}G(x,x')dx'$ How do we determine the ω_i 's? $\nabla^2 \Psi = 0 + \underbrace{\Psi(x)}_{\text{surface}}$ Basis Functions $\sigma(x) \cong \sum_{j=1}^{n} w_j \varphi_j(x)$ Integral Equation $\Psi(x)_{surface} = \int_{surface} G(x, x')\sigma(x')dS'$ $\begin{array}{lll} \text{Substitute} & \Psi(\mathbf{x}) |_{surface} & \cong \int \limits_{surface} G(x,x') \sum_{j=1}^n w_j \varphi_j(x) dS' \end{array}$ $\cong \sum_{j=1}^{n} w_{j} \int_{surface} G(x, x') \varphi_{j}(x) dS'$ If basis functions $\underbrace{\Psi(\mathbf{x})_{xurface}}_{\text{given}} \cong \sum_{j=1}^{n} w_j \int_{I_j} G(x, x') dS'$ constant unknowns $\varphi_{j}(x') = \begin{cases} 1 & \text{if } x' \in I_{j} \\ 0 & \text{if } x' \notin I_{j} \end{cases}$ We are going to try and find the weights w_j

Q1. How do I find $w_i^{(2)}$ Q2. Given w_i , what is the charge on the surface? $\sigma(x) = \sum_{j=1}^{w} w_j \varphi_j(x)$ Q3. Given w_i , what is the potential anywhere outside? $\varphi(x) = \sum_{j=1}^{w} w_j \int G(x, x') dS'$ Can evaluate this anywhere outside if we evaluate on the surface we obtain the given B.C.

Residuals

$$R(x) \equiv \Psi(x)_{boundary} - \int_{\substack{approx\\surface}} G(x, x') \sum_{i=1}^{n} \omega_i \varphi_i(x') dS$$

Residual Minimization We will pick the ω/s to make R(x) small. General Approach : Pick a set of test functions φ₁, φ₂,..., φ_n and force R(x) to be orthogonal to the set ∫φ_i(x)R(x)dS = 0 for all i.

$$\begin{split} \phi_i &= ?\\ & \int \phi_i(x) R(x) = 0 \to R(x_{l_i}) = 0 \qquad i = 1, 2, \dots, N\\ & \int \delta(x - x_{l_i}) R(x) = R(x_{l_i}) \end{split}$$

Residual Minimization Using Test Functions

$$\int \phi_l(x)R(x)dS = \int \phi_l(x)\Psi(x)dS - \int \int_{approx} \int_{approx} \phi_l(x)G(x,x')\sum_{j=1}^n \phi_j\phi_j(x')dS'dS = 0$$

We will generate the different methods by choosing the $\phi_1, \phi_2, ..., \phi_n$
Collocation : $\phi_l(x) = \delta(x - x_{l_1})$ (point - matching)







































3.22PDE Τ. **BEM Integral Equation Method II**



BASIS FUNCTION APPROACH Piecewise Constant Basis Function









- Discretize Unknowns σ(x) → N unknowns using basis functions
 Discretize Equations → Collocation or Galerkin
- Residual Definition

$$R(x) \equiv \Psi(x) - \int_{\substack{approx\\surface}} G(x, x') \sum_{i=1}^{n} \omega_i \varphi_i(x') dS'$$

 Residual Minimization We will pick the ω, 's to make R(x) small. General Approach : Pick a set of test functions φ₁, φ₂,..., φ_n and

force R(x) to be orthogonal to the set $\int \phi_i(x) R(x) dS = 0$ for all *i*. • Residual Minimization Using Test Functions

$$\int \phi_{i}(x)R(x)dS = \int \phi_{i}(x)\Psi(x)dS - \int \int_{\substack{approx \\ surface}} \phi_{i}(x)G(x,x')\sum_{j=1}^{n} \omega_{j}\phi_{j}(x')dS'dS = 0$$
We will generate the different methods by choosing the $\phi_{1}, \phi_{2}, ..., \phi_{n}$

Collocation: $\phi_i(x) = \delta(x - x_i)$ (point - matching)



Galerkin Method: $\phi_i(x) = \phi_i(x)$ (basis = test)

COLLOCATION METHOD * Basis Function Approach

Collocation : $\phi_i(x) = \delta(x_i)$ (point - matching)

$$\begin{split} \int &\delta(x - x_{t_i}) R(x) dS = R(x_{t_i}) = \Psi(x_{t_i}) - \int_{\substack{approx\\surface}} G(x_{t_i}, x') \sum_{j=1}^{n} \omega_j \varphi_j(x') dS' = 0 \\ \Rightarrow \Psi(x_{t_i}) = \sum_{k=1}^{n} \omega_j \int_{a} G(x_{t_i}, x') \varphi_j(x') dS' \end{split}$$

$$\begin{bmatrix} A_{1,1} & \cdots & A_{1,n} \\ \vdots & \ddots & \vdots \\ A_{n,1} & \cdots & A_{n,n} \end{bmatrix} \begin{bmatrix} \omega_1 \\ \vdots \\ \psi_n \end{bmatrix} = \begin{bmatrix} \Psi(\mathbf{x}_{t_1}) \\ \vdots \\ \Psi(\mathbf{x}_{t_n}) \end{bmatrix}$$



Centroid Collocation















If G(x, x') = G(x', x) then $A_{i,j} = A_{i,j}$





3D PANEL INTEGRATION Collocation Approach – Calculating Matrix Elements



Example: Thin Metal Strip Piecewise Constant **Basis Function** $\varphi_{j}(x') = \begin{cases} 1 & x' \in panel \ j \\ 0 & x' \notin panel \ j \end{cases}$ Unknowns: w/ charges on each panel Assume $\Psi(x) = 1$ given on the metal strip WI w; 11 Known Potential Wie Toeplitz Matrix Collocation $\begin{aligned} &A_{1, -6 \atop \text{nor source}} = \int_{panel -6}^{1} \frac{1}{|x_{c_1} - x'|} dS' = \frac{\text{Area Panel } 6}{|x_{c_1} - x_{c_6}|} = \frac{1}{5} \\ &A_{2, -6 \atop \text{nor source}} = \int_{panel -6}^{1} \frac{1}{|x_{c_2} - x'|} dS' = \frac{\text{Area Panel } 6}{||x_{c_1} - x_{c_2}||} = \frac{1}{4} \end{aligned}$ $A_{6,1} = \int_{ponell} \frac{1}{\|x_{c_6} - x_{c_1}^{e}\|} dS' = \frac{\text{Area Panel I}}{|x_{c_6} - x_{c_1}|\|} = \frac{1}{5}$ $A_{6,1} = A_{1,6}$ Why? Collocation in general gives $A_{6,1} \neq A_{1,6}$ But if Area_{panel 1} = Area_{panel 6}, then $A_{6,1} = A_{1,6}$ Self-term What about A1,1? $A_{1,1} = \int_{pamelt} \frac{1}{\|x_{c_1} - x'\|} dS' = \frac{\text{Area Panel I}}{\|x_{c_1} - x_{c_1}\|} = \frac{1}{0} \sum_{N_0}$



* Basis Function Approach - Calculating "Self-Term"











★ The Generalized Conjugate Residual Algorithm –
o The kth step of GCR
compute Ar^{k-1}

$$p_{k-1} \leftarrow r^{k-1} - \sum_{j=0}^{k-2} (Ap_j)^T (Ar^{k-1})p_j$$

A - Orthonormalize
search direction
 $p_{k-1} \leftarrow \frac{\tilde{p}_{k-1}}{\|A\bar{p}_{k-1}\|}$
A - Orthonormalize
search direction
 $p_{k-1} \leftarrow r^{k-1} T Ap_{k-1}$
 $p_{k-1} \leftarrow r^{k-1} T Ap_{k-1}$
complexity of symmetric GCR (e.g. for Galerkin)
complexity of symmetric GCR (e.g. for Galerkin)
complexity of symmetric GCR (e.g. for Galerkin)
 $p_{k-1} \leftarrow r^{k-1} - \sum_{j=0}^{k-2} (Ap_j)^T (Ar^{k-1})p_j$
 $\tilde{p}_{k-1} \leftarrow r^{k-1} - \sum_{j=0}^{k-2} (Ap_j)^T (Ar^{k-1})p_j$
 $\tilde{p}_{k-1} \leftarrow r^{k-1} - \sum_{j=0}^{k-2} (Ap_j)^T (Ar^{k-1})p_j$
 $p_{k-1} \leftarrow \frac{\tilde{p}_{k-1}}{\|A\bar{p}_{k-1}\|}$
 $y_{k-1} \leftarrow (r^{k-1})^T Ap_{k-1}$
 $p_{k-1} \leftarrow \frac{\tilde{p}_{k-1}}{\|A\bar{p}_{k-1}\|}$
 $p_{k-1} \leftarrow r^{k-1} - p_{k-1}Ap_{k-1}$
 $p_{k-1} \leftarrow p_{k-1}Ap_{k-1}Ap_{k-1}$
 $p_{k-1} \leftarrow p_{k-1}Ap_{k$

Number of nonzeros = $7 \cdot 10^6 \rightarrow 7$ million O(n)

3D Integral Equation Method

Number of unknowns = 10^6 Number of nonzeros = $10^{12} \rightarrow 8,000$ gigabytes $O(n^2)$





Cube Example, Interior Problem – Memory Usage

		Solv interi	ing a cube or proble			m	
	F.D.			I.E. (BEM)			
	80's	2005	<i>O</i> ()	80's	2005	<i>O</i> ()	
# panels on edge m	100	800 × 8		100	800 × 8		
# unknowns n	۱M	× 512	<i>O</i> (<i>m</i> ³)	60,000	× 64	0(6m ²	
# nonzero entries	7M	× 512	$O(7n) = O(7m^3)$	3.6 G	×4096	$O(n^2) = O(36m^2)$	
memory	56 MB	× 512 30 GB	O(56n)= $O(56m^3)$	*	×*************************************	0(8n ²) 0(290n	
				0.5 GB	×64 30 GB	O(8000 =	
					0	(48000n	
			GCR + I	Fast Me	thods		



 $\nabla \Psi = 0 \implies \Psi(x) = \int G(x, x) G(x) dx \implies Ax = 0$ (1) Choice of basis functions
(2) Choice of evaluation – Collocation or Galerkin

INTRODUCTION TO NUMERICAL SIMULATION

LECTURE 24.A.

PDE – BEM Integral Equation Method III Fast Methods for Integral Equation Solvers

TODAY'S OUTLINE:

- Fast Matrix-Vector Product
 - Multipole Algorithm (for Laplace Green Function)
 Precorrected FFT (Green Function Independent)

FAST MATRIX-VECTOR PRODUCT











Solvers





source

cluster

Key Point: If test is much further away (r larger) then the clusters can be larger (R larger by same amount)

Error $\leq k$

10% error bound









• Approximate potential at point *i*: $v_i(r_i, \phi_i, \theta_i) \approx \sum_{j=0}^{order} \sum_{k=-j}^{j} L_j^k Y_j^k(\phi_i, \theta_i) r_i^{j}$



- Complete calculation consists of:
 - 1. Build multipoles (Upward Pass).
 - 2. Build locals (Downward Pass).
 - 3. Evaluation local expansions and nearby charge potential (Evaluation Pass).



Potential Distribution



- Potential given by $\psi(x) = -\frac{x_3}{2||x||^3}$
- Charge given by $\sigma(x) = -\frac{3}{8\pi}x_3$

Discretization Convergence



· Higher-order multipole expansions needed for higher accuracy

Two Sphere Example

Potential Distribution



- Potential on each sphere : $\psi(x) = -\frac{x_3}{2\|x\|^3}$
- · Does not correspond to a simple physical problem.





Precorrected FFT (Green Function Independent)

P-FFT Matrix Vector Product

- Problem: Solve iteratively $\mathbf{A}\vec{w} = \Psi$
- At each iteration evaluate matrix-vector products Aw using the following steps:
 Grid Generation

Jild			





Inhomogeneity Problem







 Computational Complexity of PFFT++ is nearly O(N) PFFT++ is available at www.rle.mit.edu/cpg







INTRODUCTION TO NUMERICAL SIMULATION

LECTURE 24.B.

Model-Order Reduction





3.24Model Order Reduction H



electroplated windings [Daniel96]

frequency

3

Example: Micro-Inductor in a DC/DC Power Converter



How is the frequency dependency of the power loss in the inductor affecting the dynamics of the power converter and its overall efficiency?

- * Motivation: Analysis Produces Impedance vs. Frequency Curves
- How are parasitic and the resonances of the power PCB, package, IC interconnects distribution grid affecting the impedance? <u>Analysis tools</u> can produce for instance impedance vs. frequency curves.

Example: Power Grid used Feeding Circuit Blocks



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* Example: Accelerometer and RF Resonator

What is the Drag force on the fingers of a resonator or accelerometer?
How does it affect the quality factor?



Pictures generated by FastStokes (Thanks to Xin Wang)
 Example: Micro-Inductor and Resonator in a Wireless Transceiver







What are the forces applied on the mirror, how do they affect the dynamic response of the mirror?

 Example: Micromirror Switch in a Dense Wavelength Division Multiplexing Optical Communication System








FROM PDE TO A STATE SPACE MODEL

- Compare PDE Solvers and Model Order Reduction
 - PDE Solvers:
 - Accurate
 - Relatively fast (minutes to hours)
 - Challenges:
 - very large matrices: sparse (FD or FEM), or dense (BEM)
 - Application: verification and characterization of component properties
 - Model Order Reduction
 - Preserve PDE solver accuracy
 - Model construction relatively fast (same as PDE solvers minutes to hours)
 - Challenges: same as PDE solvers but only in model construction
 - Only capture input/output behavior (don't show field distributions)
 - But model evaluation in msec (e.g. get dynamical response from any input)
 - Application: analysis of functionality and interaction with other components
- * Model Order Reduction State of the Art

Model Order Reduction for simple linear systems is well understood (e.g. interconnect, heat diffusion)

$$\frac{dx}{dt} = A x(t) + B u(t)$$
$$y(t) = C x(t)$$

Not many techniques yet for NON-LINEAR systems

$$\frac{dx}{dt} = F[x(t)] + B u(t)$$
$$y(t) = C x(t)$$

Heat Conducting Bar





Temperature Differential Equation $2^{2\pi}$

$$\underbrace{\underbrace{\gamma}_{\substack{\text{specific}\\\text{heat}}} \frac{\partial T(x,t)}{\partial t} - \underbrace{\kappa}_{\substack{\text{thermal}\\\text{conductivity}}} \frac{\partial^2 T(x,t)}{\partial x^2} = h(x)\underbrace{u(t)}_{\substack{\text{specific}\\\text{input}}}$$

Spatial Discretization (except at end)



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$$\begin{split} \gamma \frac{d\hat{T}_i}{dt} &- \frac{\kappa}{\left(\Delta x\right)^2} \left(\hat{T}_{i+1} - 2\hat{T}_i + \hat{T}_{i-1} \right) = h(x_i) u(t) \quad i \in [1, \dots, N-1] \\ & \gamma \frac{d\hat{T}_i}{dt} - \frac{\kappa}{\left(\Delta x\right)^2} \left(\hat{T}_N - \hat{T}_{N-1} \right) = h(x_N) u(t) \\ & T_{end} = \hat{T}_N \end{split}$$

1 × 1
 Small amount of information to give the user
 Does not give access to interior information
 But yet accounts for it and produces correct I/O behavior
 Note: if compone gives me a different u input I do not need to

Note: if someone gives me a different u input I do not need to rediscretize the bar and solve.





Dynamic Linear Case – State-Space Description
 Original Dynamical System - Single Input/Output



Reduced Dynamical System $dx_{i}(t)$

$$\frac{dx_r(t)}{dt} = \underbrace{A_r}_{qxq} x(t) + \underbrace{b_r}_{qel} \underbrace{u(t)}_{scolar} \quad \underbrace{v_r(t)}_{scolar} = \underbrace{c_r}_{scol}^T x_r(t)$$

 $q \ll N$, but input/output behavior preserved







LECTURE 25.

Model-Order Reduction II

TODAY'S OUTLINE:

- * Problem Setup: from PDE to large ODE
- Reduction via eigenmode truncation method
- Reduction via transfer function fitting
 - Point Matching
 - Least Square
 - Quasi-convex Optimization Method
- Reduction via Projection Framework

FROM PDE TO LARGE ODE

From 3D geometry to small state space systems (MOR)





LECTURE 25. Model-Order Reduction II



- eliminate it!

 $y(t) = c^T x(t) = c^T \mathbf{E} w(t) = \left(\mathbf{E}^T c\right)^T w(t)$

Solving Decoupled Equations

 $w_i(t) = \int e^{\lambda_i(t-\tau)} \widetilde{b}_i u(\tau) d\tau$ Assuming Zero Initial Conditions

Output Equation

 $y(t) = \sum_{i=1}^{N} \widetilde{c}_{i} w_{i}(t)$

$\frac{d}{dt}w_1 = \lambda_1 w_1 + \widetilde{b}_1 u(t)$	$y(t) = \sum \widetilde{c}_i w_i(t)$
:	$\widetilde{c}_i = 0$ <i>i</i> th node unobservable!
$\frac{d}{dt}w_N = \lambda_N w_N + \widetilde{b}_N u(t)$	$\widetilde{b}_i = 0$ <i>i</i> th node uncontrollable!

Dynamic Linear Case – Reduced Models via Mode Truncation



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LECTURE 25. Model-Order Reduction II

 $y(t) = \sum_{i=1}^{q} \widetilde{c}_{i} w_{i}(t)$ Output Equation Certain modes are not affected by the input $\widetilde{b}_{k+1}, \ldots, \widetilde{b}_N$ are all small Certain modes do not affect the output $\widetilde{c}_{k+1}, \ldots, \widetilde{c}_N$ are all small

 Keep least negative eigenvalues (slowest modes) - Look at response to a constant input

$$w_{i}(t) = \int_{0}^{t} e^{\lambda_{i}(t-\tau)} \widetilde{b}_{i}u(\tau) d\tau = \underbrace{\frac{1}{\lambda_{i}} \left(\widetilde{b}_{i}u - \widetilde{b}_{i}ue^{\lambda_{i}t} \right)}_{\text{Smallif}}$$



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Bilateral Laplace Transform : $X(s) = \int_{-\infty}^{\infty} X(t)e^{-st} dt$ Key Transform Property : $sX(s) = \int_{-\infty}^{\infty} \frac{dx(t)}{dt}e^{-st} dt$

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LECTURE 25. Model-Order Reduction II





* Model Order Reduction via Eignmode Analysis Pole - Residue Form Pole - Zero Form (SISO)

$$\begin{aligned} H(s) &= \sum_{i=1}^{n} \frac{\widetilde{c}_{i}\widetilde{b}_{i}}{s - \lambda_{i}} \\ h(t) &= \sum_{i=1}^{n} \widetilde{c}_{i}\widetilde{b}_{i}e^{\lambda_{i}t} \end{aligned} \qquad H(s) &= \prod_{i=1}^{n-1} (s - \zeta_{i}) \\ &\prod_{i=1}^{n} (s - \lambda_{i}) \end{aligned}$$

Ideas for reducing order:

- Drop terms with small residues $c_i b_i$
- Drop terms with large negative Re (λ_i) ("fast" modes)
 Remove pole/zero near-cancellations
- Cluster poles that are "together"



Counting Degrees of Freedom

Dynamical System $\frac{d\hat{x}}{dt} = \hat{\mathbf{A}}\hat{x}(t) + \hat{b}u(t)$ $\hat{y}(t) = \hat{c}^T \hat{x}(t)$

$$\frac{\text{Transfer Function}}{H(s)} = \frac{r_1}{(s-\lambda_1)} + \dots + \frac{r_n}{(s-\lambda_n)}$$

$$= k \frac{(s-\zeta_1)\dots(s-\zeta_{n-1})}{(s-\lambda_1)\dots(s-\lambda_n)}$$
Degrees of Freedom
$$2q$$
coefficients

$$1 + a_1 s + \dots + a_n s^n$$

Fully Invertible Change of Coordinates

 $b_0 + b_1 s + \dots + b_{n-1} s^n$

Reduced Model Transfer Function dî an an

$$\overline{dt} = \mathbf{A}\hat{\mathbf{x}}(t) + bu(t) \implies H(s) = \hat{c}^T (sI - \hat{\mathbf{A}})^{-1} \hat{b}$$
$$\hat{\mathbf{y}}(t) = \hat{c}^T \hat{\mathbf{x}}(t)$$

Apply any invertible change of coordinates to the state $\hat{x}(t) = U\tilde{x}(t)$

Degrees of Freedom

 q^2+2q

coefficients

2q

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$$I \xrightarrow{dt} dt = U^{-1} \hat{\mathbf{A}} U \tilde{\mathbf{x}}(t) + U^{-1} \hat{b} u(t) \Rightarrow \tilde{H}(s) = \hat{c}^T (sI - \hat{\mathbf{A}})^{-1} \hat{b} = H(s)$$
$$v(t) = \hat{c}^T U \tilde{\mathbf{x}}(t)$$

Many Dynamical System have the same transfer function!!!

$$\widetilde{H}(s) = \mathbf{C}^{T} \mathbf{U} \left(s\mathbf{I} - \mathbf{U}^{-1} \mathbf{A} \mathbf{U} \right)^{-1} \mathbf{U}^{-1} b = \mathbf{C}^{T} \mathbf{U} \mathbf{U}^{-1} \left(s\mathbf{I} - \mathbf{A} \right)^{-1} \mathbf{U} \mathbf{U}^{-1} b = \mathbf{C}^{T} \left(s\mathbf{I} - \mathbf{A} \right)^{-1} b \qquad \mathbf{U} = \begin{bmatrix} \uparrow & \uparrow & \uparrow \\ u_{1} & u_{2} & \cdots & u_{q} \\ \downarrow & \downarrow & \downarrow \end{bmatrix}$$

o Model Order Reduction via Rational Transfer Function Fitting **Original System Transfer Function:**

$$H(s) = \frac{b_0 + b_1 s + \dots + b_{N-1} s^{N-1}}{1 + a_1 s + \dots + a_N s^N}$$
 rational function

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LECTURE 25. Model-Order Reduction II

Model Reduction =Find a low order $(q \le N)$ rational function matching $\hat{H}(s) = \frac{\hat{b}_0 + \hat{b}_1 s + \dots + \hat{b}_{q-1} s^{q-1}}{1 + \hat{a}_1 s + \dots + \hat{a}_q s^q}$ reduced order rational function $(sI - \mathbf{A})\mathbf{X}(s) = b\mathbf{U}(s)$ $\mathbf{X}(s) = (sI - \mathbf{A})^{-1} b \mathbf{U}(s)$ $\mathbf{Y}(s) = \underbrace{\mathbf{C}^{T}(sI - \mathbf{A})^{-1}b}_{\mathbf{H}(s)} \mathbf{U}(s)$

Point Matching

 Rational Transfer Function Fitting: via Point Matching



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Lecture 25.

Model-Order Reduction II

o Point Matching can be ill-conditioned



Also... missing data can cause severe accuracy problems

 $s_1H(s_1)a_1+\ldots+s_1^qH(s_1)a_q^r-(b_0+b_1s_1+\ldots+b_{q-1}s^{q-1})=0$



Hard to Solve Systems – Fitting Example
 Polynomial Interpolation
 Table of Data



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INTRODUCTION TO NUMERICAL SIMULATION

Lecture 25.











LECTURE 25. Model-Order Reduction II o Relaxation of the H-inf norm MOR setup [Sou, Megretski, Daniel] minimize $H(z) - \frac{p(z)}{q(z)} - \frac{r(z^{-1})}{q(z^{-1})}$ Anti-stable term subject to $\frac{\deg(q) = m, \quad \deg(p) \le m,}{\deg(r) < m}$ Stability: q(z) Schur polynomial (roots inside unit circle) Passivity, and possibly other constraints Benefit: Relaxation equivalent to a quasi-convex program Drawback: May obtain suboptimal solutions How bad is our relaxation? 0 THEOREM: Let (q^*, p^*, r^*) = arg min $H(z) - \frac{p(z)}{q(z)} - \frac{r(z^{-1})}{q(z^{-1})}$ Such that $\deg(q) = m$, q(z) - q(z) and q(z)Such that $\deg(q) = m$, q(z) is Schur polynomial Then $\left| H(z) - \frac{p(z)^*}{q(z)^*} \right|_{\infty} < m\sigma_{m+1}^{(\mu)}(H) - m^{+1^{th}}$ Hankel singular value \circ Equivalent Quasi-Convex Setup $\underset{a,b,c}{\text{minimize}} \quad \left\| H(e^{j\omega}) - \frac{b(e^{j\omega}) + jc(e^{j\omega})}{a(e^{j\omega})} \right\| \leftarrow \text{quasi-convex function}$ $\deg(a) = m, \ \deg(b) \le m$ $\deg(c) \leq m$, subject to Stability: $a(e^{j\omega}) > 0, \forall \omega \in [0, \pi]$ convex set Passivity: $b(e^{j\omega}) > 0, \forall \omega \in [0, \pi]$

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LECTURE 25. Model-Order Reduction II

This is a quasi-convex program, because

 $a(e^{j\omega}) = 2\cos(m\omega) + 2\cos((m-1)\omega)a_{m-1} + \dots + a_0 > 0$

defines an intersection of halfspaces

 $\begin{array}{l} \text{This is a } quasi-convex \text{ program, because} \\ 2\cos\left(\left[m\omega_{0}\quad(m-1)\omega_{0}\quad\dots\quad0\right]\right)\left[1\quad a_{m-1}\quad\dots\quad a_{0}\right]>0\\ 2\cos\left(\left[m\omega_{1}\quad(m-1)\omega_{1}\quad\dots\quad0\right]\right)\left[1\quad a_{m-1}\quad\dots\quad a_{0}\right]>0\\ 2\cos\left(\left[m\omega_{2}\quad(m-1)\omega_{2}\quad\dots\quad0\right]\right]\left[1\quad a_{m-1}\quad\dots\quad a_{0}\right]>0\\ 2\cos\left(\left[m\omega_{2}\quad(m-1)\omega_{2}\quad\dots\quad0\right]\right]\left[1\quad a_{m-1}\quad\dots\quad a_{0}\right]>0\\ \text{defines an intersection of halfspaces} \end{array}\right.$ o Solving the Quasi-Convex Program $\begin{array}{c} \min_{a,b,c} \left\|H\left(e^{j\omega}\right)-\frac{b(e^{j\omega})+jc\left(e^{j\omega}\right)}{a(e^{j\omega})}\right\| &\longleftarrow \text{quasi-convex set}\\ &\text{deg}\left(a\right)=m, \quad \text{deg}\left(b\right)\leq m\\ &\text{deg}\left(c\right)\leq m,\\ \text{subject to}\\ &\text{Stability: } a\left(e^{j\omega}\right)>0, \forall \omega\in[0,\pi]\\ &\text{Passivity: } b\left(e^{j\omega}\right)>0, \forall \omega\in[0,\pi] \end{array}$

Standard problem. Use for example by the ellipsoid algorithm

LECTURE 25.

Model-Order Reduction II

• Example 2: RF Inductor with Substrate (from field solver)

- RF inductor with substrate effect captured by layered Green's function [Hu Dac 05]
- · System matrices are frequency dependent
- · Full model has infinite order





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• Example 4: Model of graphic card package (from measurement)

- Industry example of a multi-port device (390 frequency samples)
- 12th order SISO reduced models are constructed
- Bounded realness constraint is imposed
- Frequency weight is employed



Example 5: Large IC power distribution grid (from field solver)

- Power distribution grid (dimension size = 7mm, wire width = $2 \mu m$)
- Blue: full model (order 2046)

• Red: QCO 40th order reduced model (positive real)





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Model-Order Reduction II



 Notes: now few variables (q << N) in the state, but still thousands of equations (N)

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LECTURE 25.

Model-Order Reduction II Approaches for picking V and U o Use Eigenvectors of the system matrix (modal analysis) Eigenmode Analysis $\mathbf{A} = \mathbf{A}^{T}$ $\mathbf{A} = \mathbf{E} \mathbf{A} \mathbf{E}^{-1} = \begin{bmatrix} \uparrow & \uparrow & \uparrow \\ E_{1} & E_{2} & \dots & E_{N} \end{bmatrix}^{L_{1}} \quad \lambda_{2} \\ \downarrow & \downarrow & \downarrow & \downarrow \end{bmatrix} \begin{bmatrix} \uparrow & \uparrow & \uparrow \\ E_{1} & E_{2} & \dots & E_{N} \end{bmatrix}^{L_{1}} \quad \lambda_{2} \\ \vdots \\ \chi(t) = \mathbf{E} \tilde{\chi} = \begin{bmatrix} \uparrow & \uparrow & \uparrow & \uparrow \\ E_{1} & E_{2} & \dots & E_{q} & E_{q+1} & \dots & E_{N} \\ \downarrow & \downarrow & \downarrow & \downarrow & \downarrow \end{bmatrix}$ Invertible change of coordinates $\mathbf{x} = \begin{bmatrix} \uparrow & \uparrow & \uparrow \\ E_{1} & E_{2} & \dots & E_{q} \\ \downarrow & \downarrow & \downarrow \\ \hline \text{pick some of these} \end{bmatrix} \tilde{\chi} \quad \text{Non-invertible change of coordinates}$ $\mathbf{x} = \begin{bmatrix} \uparrow & \uparrow & \uparrow \\ E_{1} & E_{2} & \dots & E_{q} \\ \downarrow & \downarrow & \downarrow \\ \hline \mathbf{v}_{q} \text{ eigenvectors} \\ \mathbf{v}^{T} = ? \quad \mathbf{v}^{T} \mathbf{V} = I \qquad \mathbf{v}_{q} = \begin{bmatrix} \uparrow & \uparrow & \uparrow & \uparrow \\ E_{1} & E_{2} & E_{3} & \dots & E_{q} \\ \downarrow & \downarrow & \downarrow & \downarrow \\ \downarrow & \downarrow & \downarrow \\ \end{bmatrix}$



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INTRODUCTION TO NUMERICAL SIMULATION

LECTURE 25. Model-Order Reduction II





 Use Singular Vectors of System Grammians Product (Truncated Balance Realizations)

o Use Krylov Subspace Vectors (Moment Matching)



LECTURE 26.

Model-Order Reduction III

TODAY'S OUTLINE:

- Truncated Balance Realizations
 - * Krylov Subspace Moment Matching
 - Preserving passivity
 - Need for orthogonalization (Arnoldi process)
 - Overall computational complexity
- * Comparing Truncated Balance Realizations and Krylov Subspace Moment Matching

TRUNCATED BALANCE REALIZATIONS

- * Approaches for picking V and U
 - Use Eigenvectors of the system matrix (modal analysis)
 - o Use Frequency Domain Data
 - Compute $x(s_1), x(s_2), ..., x(s_k)$
 - Use the SVD to pick $q \le k$ important vectors
 - o Use Time Series Data

 - Computer x(t₁), x(t₂), ..., x(t_k)
 Use the SVD to pick q ≤ k important vectors
 Use Singular Vectors of System Grammians Product (Truncated Balance
 - Realizations) o Use Krylov Subspace Vectors (Moment Matching)
- Observability Gramian



Energy of the output y(t) starting from state x with no input:

$$\|y(t)\|_{x}^{2} = \int_{0}^{\infty} y(t)^{T} y(t) dt = \int_{0}^{\infty} \left(\mathbf{C} e^{\mathbf{A} t} x \right)^{T} \mathbf{C} e^{\mathbf{A} t} x dt = x^{T} \left[\int_{0}^{\infty} e^{\mathbf{A}^{T} t} \mathbf{C}^{T} \mathbf{C} e^{\mathbf{A} t} dt \right] x$$

$$\frac{W_{O}}{Observabily Gramian}$$

Note: it is also the solution of $\mathbf{A}^T W_O + W_O \mathbf{A} = -\mathbf{C}^T \mathbf{C}$ Note: If $x = x_i$ the *i*th eigenvector of W_O :











 Controllability Gramian Minimum amount input energy required to drive the system to a specific state x:





Native Controllability/Observability MOR

 Suppose I could compute a basis for the strongly observable and/or strongly controllable spaces. Projection-based MOR can give a reduced model that deletes weakly observable and/or weakly controllable modes.

- o Problem: What if the same mode is strongly controllable, but weakly observable?
- Are the eignvalues of the respective Gramians even unique? Changing Coordinate System
- Consider an invertible change of coordinates: x(t) = Ux(t)
 We know that the input/output relationship will be unchanged.
 But what about the Gramians, and their cigenvalues? $\widetilde{W}_{O} \leftarrow U^{T} W_{O} U \quad \widetilde{W}_{C} \leftarrow U^{-1} W_{C} U^{-T}$
- Gramians and their eigenvalues change! Hence the relative degrees of observability and controllability are properties of the coordinate system
- o A bad choice of coordinates will lead to bad reduced models if we look at controllability and observability separately.What coordinate system should we use then?

$$\begin{split} W_{O} &= \int_{0}^{\infty} e^{A^{T}} \mathbf{C}^{T} \mathbf{C}^{e^{A^{T}}} dt \qquad \mathbf{x} = \mathbf{U} \mathbf{\tilde{x}} \\ \begin{cases} \frac{d\mathbf{\tilde{x}}}{dt} = \mathbf{U}^{-1} \mathbf{A} \mathbf{U} \mathbf{\tilde{x}}(t) + \mathbf{U}^{-1} b u(t) \\ \frac{d\mathbf{\tilde{x}}}{dt} = \mathbf{C} \mathbf{U} \mathbf{\tilde{x}}(t) \end{cases} \\ & \mathbf{\tilde{W}}_{O} &= \int_{0}^{\infty} e^{\left[\mathbf{U}^{-1} \mathbf{A} \mathbf{U}\right]^{T}} (\mathbf{C} \mathbf{U})^{T} (\mathbf{C} \mathbf{U}) e^{\mathbf{U}^{-1} \mathbf{A} \mathbf{U} t} dt \\ &= \int_{0}^{\infty} \mathbf{U}^{T} e^{A^{T}} \mathbf{U}^{-T} \mathbf{U}^{T} \mathbf{C}^{T} \mathbf{C} \mathbf{U} \mathbf{U}^{-1} e^{A^{T}} \mathbf{U} dt \\ &= \mathbf{U}^{T} W_{O} \mathbf{U} \end{cases} \\ & \mathbf{W}_{O} = \mathbf{X} \mathbf{A} \mathbf{X}^{T} \qquad \mathbf{U} \text{ invertible} \\ & \mathbf{\tilde{W}}_{O} = \underbrace{\mathbf{U}^{T} \mathbf{X} \mathbf{A} \mathbf{X}^{T} \mathbf{U} \\ & \text{ these are not the new \\ eigenvectors in general \end{cases} \qquad \mathbf{X}^{T} \mathbf{U} \neq (\mathbf{U}^{T} \mathbf{X})^{-1} \\ & \mathbf{U}^{T} \mathbf{X} \mathbf{X}^{T} \mathbf{U} = \mathbf{U}^{T} \mathbf{U} \neq I \quad = I \text{ only if } \mathbf{U} \text{ is orthonormal} \end{split}$$

Balancing Fortunately the eigenvalues of the product of the Gramians (Hankel singular values) do not change when changing coordinates:

Diagonal matrix with eigenvalues of the product

Subscription individual with eigenvalues of the product

$$W_{c}W_{O} = S\Sigma^{2}S^{-1}$$
The eigenvectors change

$$U^{-1}W_{c}U^{-T}$$

$$U^{T}W_{O}U = U^{-1}W_{c}W_{O}U = (U^{-1}S)\Sigma^{2}(U^{-1}S)^{-1}$$
And since W_{c} and W_{c} are symmetric a change of
coordinate matrix U can be found that diagonalize both
 $\Sigma\Sigma = \Sigma^{2}$. In Balanced Coordinates the Gramians are equal and diagonal
Selection of Vectors for the Columns of the Reduced Order Projection Matrix.
In balanced coordinates it is easy to select the best vectors for the reduced
model: we want the subspace of vectors that are at the same time most
controllable and observable:
 $U^{-1}W_{c}U^{-T}U^{T}W_{O}U = \Sigma^{2}$
Simply pick the eigenvecors
the diagonal (Hankel singular values)
In other words the ones corresponding to the largest eigenvalues of the
controllability and observability Gramians product.
> Truncated Balance Realization Summary
• The good news: $|H(j\omega) - H_{q}(j\omega)|_{\infty} \leq \sum (\Sigma_{q+1,q+1} + ... + \Sigma_{N,N})$

- Can do even a bit better with the optimal Hankel Reduction The bad news: $A^T W_0 + W_0 A = -C^T C$ It is expensive: Need to compute the Gramians (solve Lyapunov equation)
 Need to compute eigenvalues of the product: O(N³) The bottom line: - If the size of your system allows $O(N^3)$ computation, Truncation Balance Realization or Hankel Reduction are a much better choice than the any other reduction method. - But if you cannot afford $O(N^3)$ computation (e.g. dense matrix with N > 5000) then PRIMA or PVL or Quasi-Convex-Optimization are better choices. * Approaches for picking V and Uo Use Eigenvectors of the system matrix (modal analysis) o Use Frequency Domain Data - Compute $x(s_1), x(s_2), ..., x(s_k)$ Use the SVD to pick q < k important vectors Use Time Series Data - Compute $x(t_1), x(t_2), ..., x(t_k)$ Use the SVD to pick q < k important vectors o Use Singular Vectors of System Grammians Product (Truncated Balance Realizations) o Use Krylov Su ace Vectors (Moment Matching) Original System Transfer Function Moments $H(s) = c^{T} (sI - \mathbf{A})^{-1} b$ $\begin{aligned} F = -c^{T} \underbrace{(I - s\mathbf{A}^{-1})^{-1}}_{\text{Taylor expand}} \mathbf{A}^{-1} b \\ &= -c^{T} \underbrace{(I - s\mathbf{A}^{-1})^{-1}}_{\text{with respect to } s} \mathbf{A}^{-1} b \\ &= 1 + x + x^{2} + x^{3} + \dots \\ & \left(I - \underbrace{s\mathbf{A}^{-1}}_{\text{small}}\right)^{-1} = I + s\mathbf{A}^{-1} + s^{2}\mathbf{A}^{-2} + s^{3}\mathbf{A}^{-3} + \dots \end{aligned}$ $=\sum_{k=0}^{\infty}c^{T}A^{-(k+1)}bs^{k}$ ► ↑ _____

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- We even have bounds for the error

$$H_{r}(s) = \frac{b_{0}^{r} + b_{1}^{r}s + \dots + b_{q-1}^{r}s^{q-1}}{1 + a_{1}^{r}s + \dots + a_{q}^{r}s^{q}} = m_{0} + m_{1}s + \dots + m_{2q-1}s + \dots$$

$$c^{T} \underline{\Lambda^{-1}b} \quad c^{T} \underline{\Lambda^{-2}b} \quad c^{T} \underline{\Lambda^{-3}b} \quad \dots \quad c^{T} \underline{\Lambda^{-4}b}$$
eigenvectors
$$A^{-1}b = A^{-1}[\alpha_{1}\vec{e}_{1} + \dots + \alpha_{N}\lambda_{N}^{1}\vec{e}_{N}]$$

$$= \alpha_{1}\lambda_{1}^{-1}\vec{e}_{1} + \dots + \alpha_{N}\lambda_{N}^{1}\vec{e}_{N}$$

$$A^{-4}b = \alpha_{1}\lambda_{1}^{-4}\vec{e}_{1} + \dots + \alpha_{N}\lambda_{N}^{1}\vec{e}_{N}$$

$$d_{q-1} = \begin{bmatrix} m_{q} \\ m_{q+1} \\ \vdots \\ m_{2q-1} \end{bmatrix}$$

$$M_{n-1} \quad \dots \quad m_{2q-3} \quad m_{2q-2} \end{bmatrix} \begin{bmatrix} a_{q} \\ a_{q-1} \\ \vdots \\ m_{2q-1} \end{bmatrix} = \begin{bmatrix} m_{q} \\ m_{q+1} \\ \vdots \\ m_{2q-1} \end{bmatrix}$$

$$A \text{ Canonical Form for Model Order Reduction}$$

$$sx = Ax + \overline{b} u$$

$$y = c^{T}x$$

$$E = A^{-1} \quad \text{Assuming } A \text{ is non-singular we can east the dynamical linear system into a canonical form for model Order Reduction for model order reduction for model order reduction for model order reduction simple for educational purposes the notation simple for education and purposes the nota$$











> Combine Point and Moment Matching: Multipoint Moment Matching

- · Multipole expansion points give larger band
- Moment (derivates) matching gives more accurate behavior in between expansion points



KRYLOV SUBSPACE MOMENT MATCHING

Preserving Passivity

- Interconnected Systems
 - In reality, reduced models are only useful when connected together with models of other components in a composite simulation
 - Consider a state-space model connected to external circuitry (possibly with feedback)



 Can we assure that the simulation of the composite system will be wellbehaved? At least preclude non-physical behavior of the reduced model?
 Need to Preserve Passivity for Models of Passive Interconnect





Note: passive! Designers will connect models use them in ODE time domain simulators. If the models are not passive they can generate energy and the simulation may explode!!

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Interconnecting Passive Systems

The interconnection of stable models is not necessarily stable
 But the interconnection of passive models is a passive model:

but the interconnection of passive models is a passive model:



Sufficient Conditions for Passivity
 sEx = x + Bu
 y = Cx
 Sufficient conditions for passivity
 1. C = B^T
 2. x^TEx ≤ 0, for all x
 i.e. E is negative semidefinite
 Note that these are NOT necessary conditions (common misconceptions)
 Example Finite Difference on Poisson Equation



We already know the Finite Difference matrices is positive semidefinite. Hence $E=-A^{-1}$ is negative semidefinite.









Vectors will quickly line up with dominant eigenspace!

$$\begin{split} \mathbf{E} &= \mathbf{V} \mathbf{A} \mathbf{U}^T \qquad \mathbf{U} = \begin{bmatrix} \mathbf{b}, \mathbf{E} \mathbf{b}, \mathbf{E}^2 \mathbf{b}, \mathbf{E}^3 \mathbf{b}, \dots \end{bmatrix} \\ & b = \alpha_1 \nu_1 + \alpha_2 \nu_2 + \ldots + \alpha_N \nu_N \\ & \mathbf{E} \mathbf{b} = \alpha_1 \mathbf{E} \nu_1 + \alpha_2 \mathbf{E} \nu_2 + \ldots + \alpha_N \mathbf{E} \nu_N \\ &= \alpha_1 \lambda_1 \nu_1 + \alpha_2 \lambda_2 \nu_2 + \ldots + \alpha_N \lambda_N \nu_N \\ & \mathbf{E}^2 \mathbf{b} = \alpha_1 \lambda_1^2 \nu_1 + \alpha_2 \lambda_2^2 \nu_2 + \ldots + \alpha_N \lambda_N^2 \nu_N \\ & \mathbf{E}^{k-1} \mathbf{b} = \alpha_1 \lambda_1^{k-1} \nu_1 + \alpha_2 \lambda_2^{k-1} \nu_2 + \ldots + \alpha_N \lambda_N^{k-1} \nu_N \cong \alpha_1 \lambda_1^{k-1} \nu_1 \\ & \mathbf{V} \text{ very large} \\ & \text{ They all become linearly dependent and parallel to the eigenvector with the largest eigenvalue} \end{split}$$

Overall Computational Complexity

$$\begin{split} \vec{u}_{1} &= b / \left\| b \right\| & \text{Normalize first vector} \\ \text{For } i &= 1 \text{ to } k \\ \vec{u}_{i+1} &= E \vec{u}_{i} & \text{Generates new Krylov} \\ \text{For } j &= 1 \text{ to } i & O(22222222) \\ \vec{u}_{i+1} &\leftarrow \vec{u}_{i+1} - \left(\vec{u}_{i+1}^{T} \vec{u}_{j} \right) \vec{u}_{j} \text{ Orthogonalize new vector} \\ \vec{u}_{i+1} &\leftarrow \frac{1}{\left\| \vec{u}_{i+1} \right\|} & \vec{u}_{i+1} \\ \vec{u}_{i+1} &\leftarrow O(n) \\ \end{array}$$

Generating Vectors for the Krylov Subspace
 Most of the computation cost is spent in calculating:

$$\vec{u}_{i+1} = E_h \vec{u}_h = (A - s_h I)^{-1} \vec{u}_i$$

$$\vec{u}_{i+1} = (\mathbf{A} - s_h I)^{-i} \vec{u}_i \qquad \Longrightarrow (\mathbf{A} - s_h I) \vec{u}_{i+1} = \vec{u}_i$$

- $(A s_h I)^{-1} \vec{u}_{i+1} = \vec{u}_i$
- If we have a good preconditioner and a fast matrix vector product each new vector is calculated O(n)

O(n)

- The total complexity for calculating the projection matrix U_q is O(qn)

COMPARING TRUNCATED BALANCE REALIZATIONS AND KRYLOV SUBSPACE MOMENT MATCHING

 $\begin{aligned} \mathbf{U}_{q}^{T}\mathbf{E}\mathbf{U}_{q}\frac{d\hat{\mathbf{x}}}{dt} &= \hat{\mathbf{x}}(t) + \mathbf{U}^{T}b\hat{\mathbf{x}}(t) \quad \text{Reduced Model} \\ \hat{\mathbf{y}}(t) &= \mathbf{C}^{T}\mathbf{U}_{q}\hat{\mathbf{x}}(t) \end{aligned}$



MOR technique	Computational Complexit
TBR ('81) Hankel ('84) (have error bounds)	Bottleneck: SVD O(N^3) e.g. 10months, 80GB, for $N=100,000$
Moment matching ('97) (no error bounds)	Bottleneck: matrix-vector prod $O(qN^2)$ e.g. 7days, 80GB, for <i>N</i> =100,000 <i>q</i> =10
Moment Matching + pFFT matrix-vector	O(<i>qN</i> log <i>N</i>) e.g. 8hours, 0.3GB for <i>N</i> =100,000 <i>q</i> =1
QuasiConvex Optimization + pFFT field solver	$O(mN \log N)$ Same as above, but have error bound!

Model Order Reduction Computational Complexity (time and memory)