

Lecture 17: Repeated Games 2

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Agenda:

- Repeated games
- Imperfect monitoring
 - price trigger strategies

Recall: v^* set of feasible and strictly individually rational pay off set.

Theorem 1 (Nash folk theorem:) $\forall v \in V^*; \exists \underline{\delta} < 1$ s.t. $\forall \delta \geq \underline{\delta}$ there is a NE of the $G^\infty(\delta)$ with payoffs v .

Proof: Strategy constructed \Rightarrow grim trigger with punishment by minmax.

If i plays the strategy \Rightarrow he gets v_i

If i deviates from the strategy in some period t ; then denoting $\bar{v}_i = \max_{\mathbf{a}} g_i(\mathbf{a})$, the most that I could get. $(1 - \delta)[v_i + \delta v_i + \dots + \delta^{t-1}v_i + \delta^t \bar{v}_i + \delta^{t+1} \underline{v}_i + \delta^{t+2} \underline{v}_i + \dots]$

following the suggested strategy will be optimal if,

$$\begin{aligned} v_i &\geq (1 - \delta)v_i + \delta^t(1 - \delta)\bar{v}_i + \delta^{t+1}\underline{v}_i \\ &= v_i - \delta^t[v_i + (1 - \delta)\bar{v}_i - \delta\underline{v}_i + (\delta v_i - \delta v_i)] \end{aligned}$$

so if we can show that the expression in the bracket, is non negative then we are done. So we need to show that:

$$\begin{aligned} &\Rightarrow (1 - \delta)(v_i - \bar{v}_i) + \delta(v_i - \underline{v}_i) \geq 0 \\ &\Rightarrow \frac{\delta}{(1 - \delta)}(v_i - \underline{v}_i) \geq \bar{v}_i - v_i \end{aligned}$$

simply pick: $\underline{\delta} = \max_i \frac{v_i - \bar{v}_i}{2v_i - \bar{v}_i - \underline{v}_i}$. Check that $\underline{\delta} < 1$. (Since $\underline{v}_i < v_i$). \square

Remark: Nash folk theorem says that essentially anything goes as a NE when players are patient enough.

\Rightarrow strategies involve unrelenting punishment which may be very costly for the punisher to carry out \Rightarrow non-credible threat.

Example:

	L (q)	R ($1 - q$)
U	6, 6	0, -100
D	7, 1	0, -100

The unique NE: (D, L) . On the other hand by minmax: $\underline{v}_1 = 0, \underline{v}_2 = 1$

\Rightarrow strategies constructed in the proof require play 2 to play R in every period following a deviation!

Perfect folk theorems:

1. Friedman: Any feasible payoff $> e_i$, where $e = [e_i]$ $i = 1, \dots, I$, the pay off vector for some static equilibrium, can be obtained as an SPE of $G^\infty(\delta)$

Proof: grim trigger with punishment by NE in phase 2 (subgame perfect)

2. What happens if static NE payoffs $>$ minmax payoff?

- In PD minmax = NE payoff

Theorem 2 *Perfect folk theorem: (Fudenberg & Maskin) Assume that $\dim(v^*) = I$ (non empty interior) then for any $v \in V^*$; $\exists \underline{\delta} < 1$; s.t. $\forall \delta \geq \underline{\delta}$ there is an SPE of $G^\infty(\delta)$ with payoffs v .*

Proof Sketch: Fix $v \in V^*$. Construct an SPE that achieves it.

Again assume that \exists some action profile $\mathbf{a} = (a_1, \dots, a_I)$ s.t. $g_i(\mathbf{a}) = v_i$

Key to the proof: Find strategies that allows us to "reward" agents $j \neq i$ in the event i deviates and has to be minmaxed.

Steps:

- choose $\hat{v} \in \text{Int}(V^*)$ s.t. $\underline{v}_i < \hat{v}_i < v_i \forall i$
- choose $\epsilon > 0$ s.t. the vector

$$\hat{v}(\epsilon) = (\hat{v}_1 + \epsilon, \hat{v}_2 + \epsilon, \dots, \hat{v}_i, \hat{v}_{i+1} + \epsilon, \dots, \hat{v}_I + \epsilon)$$

- let \mathbf{a}^i be the profile with $g(\mathbf{a}^i) = \hat{v}(\epsilon)$
- choose T (punishment length) s.t. $\forall i$;

$$\max_a g_i(a) + T\underline{v}_i < \min_a g_i(a) + T\hat{v}_i$$

- let $m_i = (m_i^i, m_{-i}^i)$ be the profile that minmaxes i

$$g_i(m_i^i, m_{-i}^i) = \underline{v}_i$$

Now consider the following strategy for $i = 1, \dots, I$

- 1) play a_i ($g(\mathbf{a}) = v$) so long as no one deviates. If j alone deviates, go to 2_j
- 2 _{j}) play m_j^i for T periods then go to 3_j , If no one deviates. If during 2_j , k deviates, restart phase 2_k
- 3 _{j}) play \mathbf{a}^j ($g(\mathbf{a}^j) = \hat{v}^j(\epsilon)$) so long as no one deviates. If k deviates, go to 2_k

□

remarks:

1. Note that strategies involve punishments & rewards.
 2. Check if is an SPE
 3. $\dim v^* = I \Rightarrow$ ensures that each player i can be *singled out* for a punishment for a deviation.
- Folk theorem for finitely repeated game:
 - if the stage game has a unique NE, then unique perfect equilibrium is to play the static NashE in every period of the subgame.
 - with multiple stage NE \Rightarrow can use rewards & punishments
 - Equilibrium concept do very little to pin down play by patient player?...?

Repeated games with imperfect public monitoring

Before each player observed the actions of others at the end of each period.

\Rightarrow instead, assume players observe only an imperfect signal of stage game actions.

Motivational Example Green & Porter (84) Cournot competition with noisy demand. Firms set output q_1^t, \dots, q_I^t privately. Level of demand stochastic. Each firm's pay off depends on his own output and on the publicly *observed market price*. Firms do not observe each other's output levels.

Low market price $\rightarrow \begin{cases} \text{high produced output} \\ \text{low demand} \end{cases}$

General Model:

- A_1, \dots, A_I finite action sets: $A = A_1 \times \dots \times A_I$ Each action profile induces a probability distribution over $y \in Y$
- y_i rv $\in Y$: finite set.
 $\pi(\cdot, \mathbf{a})$: pmf on y parametrized by the action profile \mathbf{a} .
- Player i 's realized payoff $r_i(a_i, y)$
- Expected stage payoff.

$$g_i(\mathbf{a}) = \sum_{y \in Y} \pi(y, \mathbf{a}) r_i(a_i, y)$$

- Strategy $\{s_i^t\}$

$$s_i^t : h^t \rightarrow A_i$$

$$h^t = (y^0, \dots, y^{t-1})$$

- Player i 's average discounted payoff.

$$(1 - \delta) \sum_{t=0}^{\infty} \delta^t g_i^t(\mathbf{a})$$

example: Green & Porter's Noisy Cournot Model: Consider the following trigger strategy:

1. Play q_1, \dots, q_I (desirable outcome) If $p^t \leq f$, go to phase 2
2. Play q_1^c, \dots, q_I^c (static NE) for T periods then return to 1